

Universal gates from braiding and fusing anyons on quantum hardware

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Chiu Fan Bowen Lo^{1✉}, Anasuya Lyons¹, Dan Gresh², Michael Mills², Peter E. Siegfried², Maxwell D. Urmey², Nathanan Tantivasadakarn³, Henrik Dreyer⁴, Ashvin Vishwanath¹, Ruben Verresen^{5✉} & Mohsin Iqbal^{4✉}

A quantum computer requires the ability to store and manipulate information globally to protect against local noise. Topologically ordered phases^{1,2} offer two routes: encoding information in the ground-state subspace³ or in anyonic excitations^{1,4,5}. The toric code¹ exemplifies the first approach but does not intrinsically support a universal gate set. The latter—topological quantum computation—implements gates by braiding non-Abelian anyons⁶ around each other. However, the simplest non-Abelian generalizations of the toric code cannot achieve universality by braiding alone^{7–9}. Here we demonstrate that anyon fusion, used as a computational primitive, renders these minimally non-Abelian topologically ordered states universal. We prepare a 54-qubit ground state of the quantum double of S_3 , the smallest non-Abelian group, on the H2 processor of Quantinuum. We encode logical information in the global fusion space of non-Abelian anyons, and by combining braiding with fusion, we realize a universal topological gate set and read-out, which we demonstrate by topologically preparing a magic state. This demonstrates that the S_3 topologically ordered state is scalably preparable, yet rich enough to support a universal gate set. More broadly, this work opens up new pathways for harnessing the intrinsic properties of quantum matter to manipulate quantum information.

The appeal of topologically ordered phases of matter lies in their ability to encode information in global properties insensitive to local perturbations^{1,2}. The toric code¹ exemplifies this idea: an emergent $\mathbb{Z}_2$ gauge theory¹⁰ whose anyonic quasiparticles exhibit non-trivial exchange statistics^{11,12}, forming the foundation of the paradigmatic surface code^{13,14}, where quantum information is stored nonlocally in the degenerate ground-state subspace³. A qualitatively different approach is known as topological quantum computation, in which information is directly encoded in the internal degrees of freedom of non-Abelian anyons^{1,4}, which is minimally realized in quantum double models based on finite non-Abelian groups¹.

Until recently, the implementation of topological quantum computation was out of experimental reach. However, the discovery that, for solvable groups, constant-depth adaptive circuits can prepare non-Abelian quantum double models^{15–19} as well as their anyons^{19–21} has prompted renewed theoretical interest in their utility^{22–26}. These measurement-based approaches substantially lower the experimental barriers to realizing these phases on quantum hardware, with a recent experimental realization of D_4 topological order²⁷.

Despite increased experimental accessibility to non-Abelian defects and topological orders^{27–33}, a fundamental limitation remains. Although topological orders based on finite groups are desirable because of being minimal non-Abelian generalizations of the toric code, their braiding operations do not provide a universal set of quantum gates^{7–9}. This

stands in contrast to more exotic topological phases such as Fibonacci topological order, in which braid-only universality can be achieved⁴, albeit at the cost of significantly greater complexity in their realization^{29,31}, in part because of conjectured no-gos for scalable preparation protocols¹⁵. Moreover, their irrational quantum dimensions complicate the search for leakage-free gates^{34–36}, in contrast to the integer quantum dimensions of anyons in quantum doubles. Moreover, although a finite decoding threshold has recently been proven for quantum doubles with solvable groups by leveraging the similarity with the toric code³⁷, no similar results exist for more exotic topological order.

Even in idealized settings for topological quantum computation, however, measurements of anyon type are required to extract information encoded in non-Abelian anyons^{1,38}. This observation raises a natural question: given measurement has already been proven instrumental for preparing and manipulating non-Abelian topological phases, can it also be harnessed to enhance their computational power? Theoretical work has suggested that topological orders associated with cyclic anyon fusion rules³⁹, of the form $a \times \bar{a} = a + \dots$, become universal when anyon fusion measurement is incorporated into the gate set^{8,9,40–44}. Although this cyclic criterion excludes the previously realized D_4 topological order, it applies to the quantum double of the smallest non-Abelian group, S_3 (refs. 41,45–47).

This work presents, to our knowledge, the first experimental realizations of a universal topological gate set based on the braiding and

¹Department of Physics, Harvard University, Cambridge, MA, USA. ²Quantinuum, Broomfield, CO, USA. ³C. N. Yang Institute for Theoretical Physics, Stony Brook University, Stony Brook, NY, USA.

⁴Quantinuum, Munich, Germany. ⁵Pritzker School of Molecular Engineering, University of Chicago, Chicago, IL, USA. ✉e-mail: chiufanbowenlo@g.harvard.edu; verresen@uchicago.edu; mohsin.iqbal@quantinuum.com

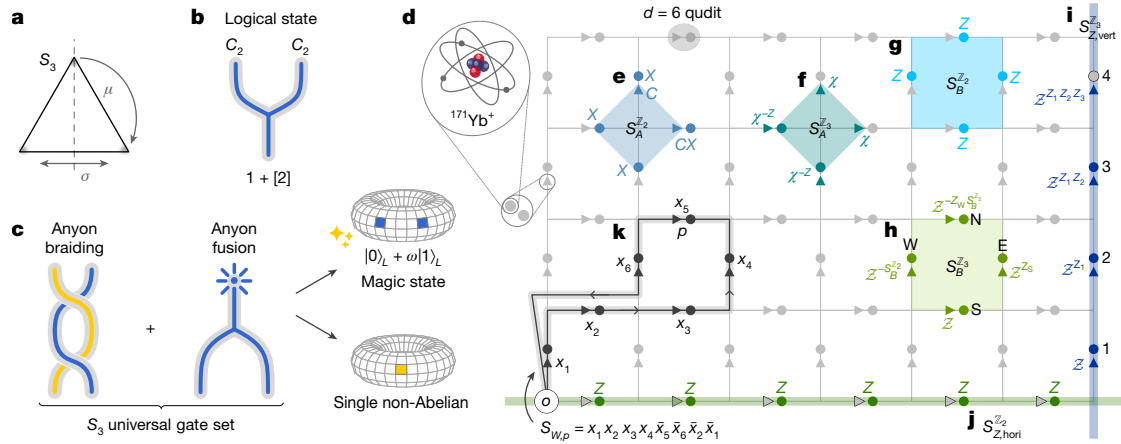


Fig. 1 | Topological order based on S_3 . **a**, Symmetry group $S_3 = \mathbb{Z}_3 \rtimes \mathbb{Z}_2$ of an equilateral triangle, generated by rotation μ and reflection σ . **b**, Logical state encoded in the flux-neutral fusion space of a pair of C_2 fluxes. **c**, Anyon braiding together with fusion realizes a universal gate set for quantum computation. **d**, The code is defined on a square lattice on a torus with oriented edges hosting $d = 6$ qudits, composed of a qubit (solid circle) and a qutrit (solid triangle). The triangle direction specifies the edge orientation. Physically, each qubit is encoded in the hyperfine states of $^{171}\text{Yb}^+$ ions. **e–k**, The stabilizers decompose into Pauli operators X, Z on qubits and generalized Pauli operators X, Z

together with charge conjugation C on qutrits (Supplementary Information section 1B). There are qubit vertex stabilizers $S_A^{Z_2}$ (e), qutrit vertex stabilizers $S_A^{Z_3}$ (these act on qutrits, conditioned on qubits) (f), qubit plaquette stabilizers ($S_B^{Z_2}$) (g) and qutrit plaquette stabilizers ($S_B^{Z_3}$) (h). There are also non-contractible versions of the $S_Z^{Z_3}$ (i) and $S_Z^{Z_2}$ stabilizers (j), which toggle between ground states. The stabilizer $S_{W,p}$ associated with the global flux W is defined as the product of S_3 group elements x_i along the directed loop in **k**, beginning at the origin o and encircling exactly one plaquette p (Supplementary Information section 2).

fusion of non-Abelian anyons in a scalable non-Abelian topological order, in particular the S_3 quantum double model on a trapped-ion quantum computing architecture⁴⁸. Using a two-step process, which can be interpreted as coupling a $\mathbb{Z}_2$ and $\mathbb{Z}_3$ toric code^{15–17,19}, this enables the realization of S_3 topological order ground states as well as the realization of a single non-Abelian anyon on a torus⁴⁹, which we link to the cyclicity of the anyon fusion rules. We nonlocally encode logical information in the internal states of non-Abelian qutrits, and demonstrate a universal gate set on these topological qutrits, consisting of three primitive operations: (1) an entangling gate implemented by coherent anyon braiding; (2) measurement of the topological qutrit in the X -basis; and (3) measurement of the topological qutrit in the Z -basis. Any quantum operation on pairs of qutrits (and therefore also qubits) can be synthesized from these three primitives^{41,47}. Although these primitives cannot be realized in a strictly transversal way^{50,51}, and we use linear-depth circuits with measurements for our demonstrations, they can be scalably implemented with a constant-depth adaptive circuit. We further demonstrate how to prepare reference states (the so-called ‘bureau of standards’⁵²) and how our primitives can be used to prepare a magic state on a topological qutrit (Fig. 1a–c).

The S_3 quantum double

The internal Hilbert space consists of six-state qudits labelled by the elements of the group $S_3 = \langle \mu, \sigma | \mu^3 = \sigma^2 = (\mu\sigma)^2 = 1 \rangle$ (Fig. 1a). We encode this as a qutrit–qubit pair $g = \mu^a \sigma^b$, where a is the qutrit and b the qubit value. On the Quantinuum hardware, the qutrit is itself composed of two qubits, with each qubit being encoded in an ion (Fig. 1d). Errors caused by leakage into the extra qubit state can be mitigated through heralding and post-selection (Supplementary Information section 4).

The S_3 quantum double⁴¹ ground states are stabilized by the following qubit and qutrit projectors:

$$A_v^{Z_2} = \frac{1}{2}(1 + S_A^{Z_2}), \quad A_v^{Z_3} = \frac{1}{3}(1 + (S_A^{Z_3} + \text{h.c.})), \quad (1)$$

$$B_p^{Z_2} = \frac{1}{2}(1 + S_B^{Z_2}), \quad B_p^{Z_3} = \frac{1}{3}(1 + (S_B^{Z_3} + \text{h.c.})), \quad (2)$$

where the qubit and qutrit vertex and plaquette stabilizers, $S_A^{Z_2}$, $S_A^{Z_3}$, $S_B^{Z_2}$, and $S_B^{Z_3}$ are defined in Fig. 1e–h. $S_A^{Z_2}$ ($S_A^{Z_3}$) represent the $\mathbb{Z}_2$ ($\mathbb{Z}_3$) X -type stabilizers imposing Gauss law constraints at a vertex, such that the corresponding projector $A_v^{Z_2}$ ($A_v^{Z_3}$) projects into the zero-charge sector at vertex v . The plaquette stabilizer $S_B^{Z_2}$ ($S_B^{Z_3}$) are Z -type stabilizers detecting $\mathbb{Z}_2$ ($\mathbb{Z}_3$) flux, such that the projector $B_p^{Z_2} B_p^{Z_3}$ projects into the trivial-flux sector at plaquette p . The resulting state can be interpreted as a deconfined S_3 gauge theory, whose quasiparticles exhibit anyonic exchange statistics due to the Aharonov–Bohm effect.

There are eight anyon types in the S_3 quantum double model (Supplementary Information section 1C). Here, we work with three types of non-Abelian anyons: the charge transforming under the two-dimensional $[2]$ representation of S_3 , detected by a violation of $A_v^{Z_3}$ (‘ $[2]$ charge’); the two-dimensional flux that violates a $B_p^{Z_3}$ plaquette, with flux state labelled by the conjugacy class C_3 represented by μ (‘ C_3 flux’); and the three-dimensional flux, which violates both $B_p^{Z_2, Z_3}$ plaquettes, with flux state labelled by the conjugacy class C_2 represented by σ (‘ C_2 flux’).

State preparation

Ground-state preparation

We use the all-to-all connectivity of the trapped-ion platform to realize the S_3 quantum double ground states on a torus, in which there are eight ground states in the S_3 quantum double, in one-to-one correspondence with the eight anyon types. We label the ground states by non-contractible line operators (Fig. 1i, j).

The state-preparation protocol mirrors the decomposition of S_3 group elements into a qutrit and a qubit, using the fact that S_3 can be expressed as a semidirect product $\mathbb{Z}_3 \rtimes \mathbb{Z}_2$. We use a gauging procedure to prepare a $\mathbb{Z}_3$ qutrit toric code³⁰ and subsequently apply another gauging operation (that in isolation would produce a $\mathbb{Z}_2$ qubit toric code) to obtain the ground states of S_3 topological order¹⁷. We implement this construction using a fully unitary circuit (Supplementary Information section 6). The measured projector expectation values are shown in Fig. 2a (see Methods, section ‘Ground-state preparation’ for details of the expectation value estimation). Because the target state $|g_s\rangle$ is fully specified by its projectors, we bound the per-qudit

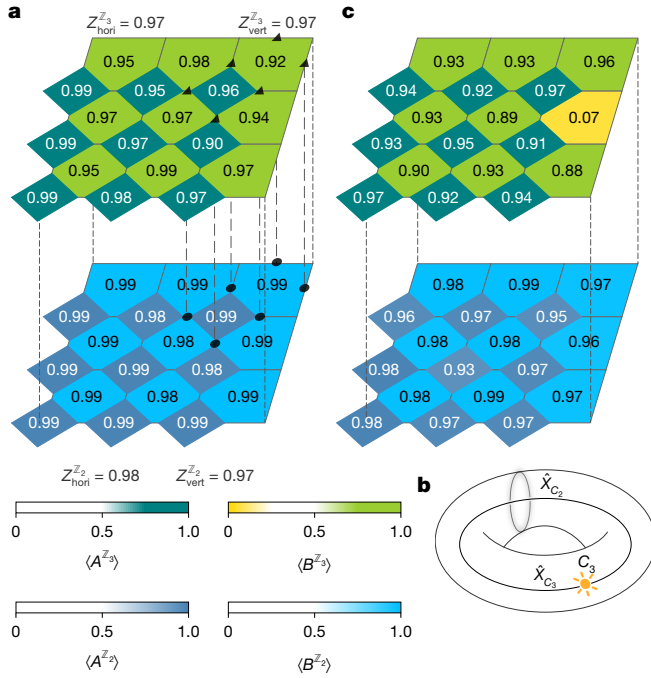


Fig. 2 | Ground and single-anyon states on 54 qubits. a, Z_2 (bottom) and Z_3 (top) vertex and plaquette projector expectation values on a 3×3 periodic square lattice (see equations (1) and (2)). Lighter squares denote plaquette projectors and darker rhombi denote vertex projectors. The mean Z_2 and Z_3 projector values are 0.987(2) and 0.962(4), respectively; the average standard error is 0.014, and the maximum standard error is 0.029. For the non-contractible projectors, the average (maximum) standard error is 0.017 (0.022). **b,** Single C_3 anyon creation by applying a horizontal $\hat{X}_{C_3}$ to a ground state with a vertical C_2 flux loop. **c,** The anyon is identified by a single excited B^{Z_3} plaquette projector. The average standard error is 0.020, and the maximum standard error is 0.034.

fidelity of the prepared state ρ on the 3×3 lattice with 18 sites as (Supplementary Information section 8)

$$0.970(5) \leq \langle \text{gs} | \rho | \text{gs} \rangle^{1/18} \leq 0.988(3). \quad (3)$$

As S_3 is a solvable group, it also admits a constant-depth adaptive preparation^{17,19}. We find a comparable fidelity between 0.930(8) and 0.978(2) (see Extended Data Fig. 2 and Supplementary Information section 6B1). We use the ground states prepared unitarily for the remaining demonstration because of the slightly better fidelity for the current system size; however, the measurement-based preparation protocol guarantees scalability of future experiments.

A single non-Abelian anyon on the torus

A hallmark of non-Abelian topological order is the existence of low-lying excited states hosting a single anyon on the torus. A recent realization of D_4 topological order demonstrated an excited state with a single unpaired Abelian anyon²⁷. Although such an unpaired anyon certifies the non-Abelianness of that state, here we target a more complex and computationally powerful quantum phase: a cyclic topological order. Cyclicity is expected to enable universal computation by braiding and fusion^{9,40}. A diagnostic of cyclic fusion rules is the ability to generate a single non-Abelian anyon on the torus by braiding an identical anyon pair around a non-contractible loop.

Braiding a C_2 flux pair around the vertical non-contractible cycle of the torus generates a non-trivial ground state with C_2 background flux. Starting from this non-trivial ground state, we can braid a C_3 flux pair around the horizontal non-contractible cycle (Fig. 2b). Owing to

the non-trivial braiding between the background C_2 flux and the C_3 anyons, the C_3 pair fuses to a single C_3 anyon after braiding around the horizontal cycle (Fig. 2c). This constitutes the experimental realization of a single non-Abelian flux anyon in a lattice model with topological order. See Supplementary Information section 7 for more details on the preparation protocol.

We note that we will not use this single non-Abelian to store quantum information. Rather, the above demonstrates the cyclicity of the fusion rules, which is a signature of the computational power of the S_3 quantum double. We next demonstrate explicit building blocks of a universal gate set.

Universal topological gate set

Logical qutrit encoding

To robustly store quantum information in nonlocal degrees of freedom, we encode it in the fusion space of spatially separated anyons. In particular, we use a pair of C_2 fluxes with trivial total flux (Fig. 1b), which ensures that the encoded information is unaffected by transporting pairs across the system. Despite its flux neutrality, the pair retains delocalized charge degrees of freedom due to the fusion rule $C_2 \times C_2 = 1 + [2] + \dots$. The associated fusion space already yields a two-dimensional logical subspace. Moreover, the internal structure of the $[2]$ charge sector is itself protected: the two charge states cannot be distinguished or toggled by any local operation that does not enclose both fluxes. We thus see that the separation distance between anyons gives us a code distance. Provided the fluxes remain well separated, we obtain a protected three-dimensional fusion space, realizing a topological qutrit.

We can label this logical qutrit in terms of the fluxes of the constituent C_2 fluxes:

$$|j\rangle_L = |\mu^j \sigma, \mu^j \sigma\rangle, \quad j = 0, 1, 2, \quad (4)$$

defining our $\mathcal{Z}$ -basis. Moreover, the nonlocal charges are diagonal in the logical $\mathcal{X}$ -basis:

$$|\tilde{j}\rangle_L = \frac{1}{\sqrt{3}}(|0\rangle_L + \omega^j |1\rangle_L + \bar{\omega}^j |2\rangle_L), \quad j = 0, 1, 2 \quad (5)$$

where $\omega = e^{2\pi i/3}$. It can be checked that the $|\tilde{0}\rangle_L$ flux pair is in the trivial anyon sector, whereas the $|\tilde{1}\rangle_L$ and $|\tilde{2}\rangle_L$ flux pairs behave as a (delocalized) $[2]$ charge.

For the flux labels in equation (4) to be well-defined for spatially separated C_2 anyons, we can consistently measure the global flux W associated with a loop based at a fixed-but-arbitrary choice of origin (see Fig. 1k for an example of a global flux). We call this the ‘absolute’ encoding. Although this is sufficient to define topological flux and demonstrate universal gate sets (see the following sections), the origin defines an Achilles’ heel through which quantum information can leak out by local charge measurement. This is avoided by a ‘relative’ encoding, where instead of using a fixed origin, the flux values are relative to other pairs. This requires building a ‘bureau of standards’⁵², which we demonstrate in the section ‘Bureau of standards’. We note that this approach is fully topological, in the sense that no individual local measurement can corrupt the logical information, and if so desired, it can be expressed in terms of fusion trees (Supplementary Information section 9).

In ref. 47, it is explained that three primitives (pull-through gate, $\mathcal{Z}$ -basis measurement and $\mathcal{X}$ -basis measurement) constitute a universal gate set. Here we demonstrate the experimental realization of these primitives. The braiding of non-Abelian fluxes is enabled by technical advancements in protocols to coherently move non-Abelian anyons using local flux label while preserving the global logical information (C. F. B. Lo et al., manuscript in preparation) (see Supplementary Information section 3 for theory overview). At the circuit level, the braiding

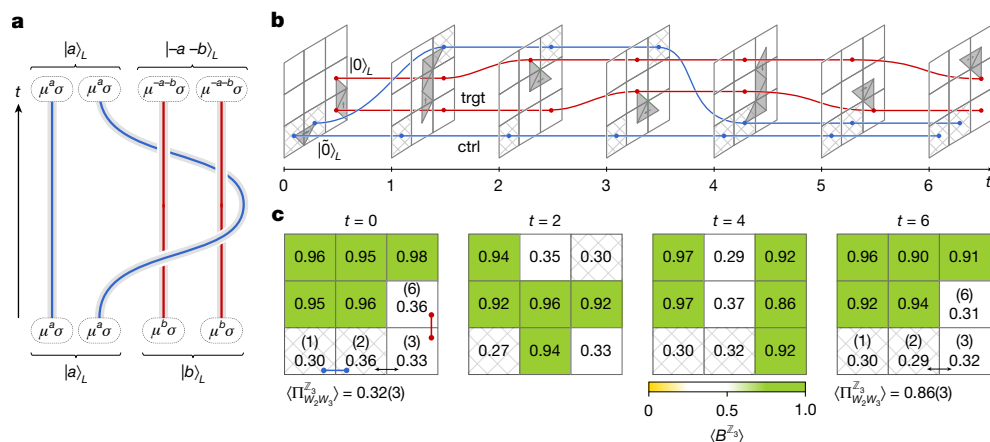


Fig. 3 | Pull-through gate entangles topological qutrits. a, Braiding between the logical pairs (equation (4)) gives an entangling gate (equation (6)) by the conjugation $(\mu^a \sigma) \mu^b \sigma (\mu^a \sigma)^{-1}$. **b**, Spacetime diagram of the protocol on the 3×3 torus. Two pairs of C_2 flux anyons, control (blue worldlines) and target (red worldlines), are initialized in the logical state $|\tilde{0}\rangle_L |0\rangle_L$. One endpoint of the control qutrit is braided (pulled through) around both endpoints of the target and returned to its initial position at $t=6$. Oriented grey triangles indicate the ribbon operators implementing anyon motion. **c**, Expectation values of the

of non-Abelian anyons is implemented by linear-depth (in the distance the anyons travel) unitary circuits with one ancilla; explicit circuits can be found in Supplementary Information section 10.

Pull-through gate

We begin with the entangling primitive, implemented by braiding a constituent C_2 flux of one logical qutrit around the flux pair of another logical qutrit (Fig. 3a). This operation, known as the pull-through gate, exploits the non-Abelian braiding relations of the S_3 quantum double. In general, when fluxes braid around each other, they are each conjugated by their total flux. In this particular case, this leads to the logical two-qutrit gate

$$U|a, b\rangle_L = |a, -a-b\rangle_L. \quad (6)$$

We will demonstrate the action of this pull-through gate on $|\tilde{0}\rangle_L |0\rangle_L$, which creates a qutrit Bell state,

$$U|\tilde{0}\rangle_L |0\rangle_L = \frac{1}{\sqrt{3}}(|0, 0\rangle_L + |1, 2\rangle_L + |2, 1\rangle_L), \quad (7)$$

where the logical information is entangled across the flux pairs (Fig. 3b). Figure 3c shows the evolution of the B^{Z_3} projectors expectation values during the protocol (see Methods, section ‘Pull-through gate’ for protocol details and Extended Data Fig. 3 for further data). The B^{Z_3} projectors identify the locations of logical qutrit endpoints, revealing local information but not the encoded logical state.

To verify the resulting entanglement between the two flux pairs, we certify correlations in the Z -basis, quantified by $\langle W_{p_1}^{Z_3} W_{p_2}^{Z_3} \rangle$, where W_{p_i} denotes the topological W -flux operator encircling plaquette p_i (see Supplementary Information section 2C for the precise path). The observed increase of this correlator from its initial value at $t=0$ to near unity at $t=6$, as both qutrits return to their original lattice positions, provides strong evidence that the prepared state closely approximates the ideal target state.

$\mathcal{X}$ -basis measurement

Apart from braiding logical anyons, fusion of anyons also implements logical operations, in addition to providing a read-out mechanism. The partial $\mathcal{X}$ -basis measurement, which distinguishes between the charge-free $|\tilde{0}\rangle_L$ state and the charged $\{|\tilde{1}\rangle_L, |\tilde{2}\rangle_L\}$ subspace, is implemented

local projectors B^{Z_3} remain close to 1 except on plaquettes hosting logical qutrit endpoints, where they are approximately $1/3$; the average (maximum) standard error is 0.028 (0.044). The nonlocal correlator $\langle \Pi_{W_2 W_3}^{Z_3} \rangle$ certifies Z correlations between the control and target qutrits. Its measured value is consistent with the predicted value of $1/3$ at $t=0$ and increases to 0.86(3) at $t=6$, in good agreement with the theoretical value of 1 for the prepared qutrit Bell state.

by braiding and fusion of C_2 fluxes (Fig. 4a). This coherent detection of $|\tilde{0}\rangle_L$ is sufficient for a universal gate set on the $\{|0\rangle_L, |1\rangle_L\}$ qubit⁴⁷.

We demonstrate the $\mathcal{X}$ -basis measurements on two input states: $|\tilde{0}\rangle_L$ and $|\tilde{1}\rangle_L$. The $|\tilde{0}\rangle_L$ state is charge neutral, whereas the $|\tilde{1}\rangle_L$ state stores a $[2]$ charge nonlocally (that is, the flux pair fuses to a $[2]$ charge). Therefore, the way to measure in the $\mathcal{X}$ -basis is by braiding with an ancilla C_2 flux pair and checking the fusion outcome: if the pair fuses to vacuum, then we have measured in the $|\tilde{0}\rangle_L$ state, whereas if the pair fuses to a $[2]$ charge, then it is orthogonal to the $|\tilde{0}\rangle_L$ state. In the latter case, we can repeat the braiding to remove the $[2]$ charge.

The projection violations for the input states $|\tilde{0}\rangle_L$ are shown in Fig. 4b (top row) and $|\tilde{1}\rangle_L$ (bottom row) (see Extended Data Fig. 4 for further data and Supplementary Information section 10B for implementation details). For the $|\tilde{0}\rangle_L$ input state, the vertex projector A^{Z_3} is not violated after fusion (maintaining a value close to 1), demonstrating that we have measured the state to be in the $|\tilde{0}\rangle_L$ state. For the $|\tilde{1}\rangle_L$ input state, A^{Z_3} is violated after the fusion (with the stabilizer value close to 0), indicating the presence of a $[2]$ charge at the vertex. This demonstrates that the state is in the subspace spanned by the $|\tilde{1}\rangle_L, |\tilde{2}\rangle_L$ state.

Z -basis measurement

As the final primitive that completes the universal gate set, the Z -basis measurement is a comparison measurement: the measurement outcome indicates whether the data state is the same as the reference state or orthogonal to it. At the physical level, the Z -basis measurement determines the flux content of the logical flux pair by braiding and fusion of a $[2]$ charge pair around one flux each from the data and the reference pair (Fig. 4c). If the data and reference flux pairs are in the same internal state, then the $[2]$ charge pair winds around a trivial flux and fuses back to vacuum. If the data and reference flux pairs are in different internal states, then the total flux is a C_3 group element, which results in non-trivial braiding: there is a $3/4$ probability of obtaining a remnant $[-]$ charge after fusion.

We first instantiate the data logical state (here $|0\rangle_L$) in the middle row and the reference logical states $|0\rangle_L$ and $|1\rangle_L$ in (Fig. 4c, bottom and top rows). We then braid $[2]$ charge pairs to do comparison measurement of the data flux pair with the reference pairs; the projection violations are shown in Fig. 4d (see Extended Data Fig. 5 for further data and Supplementary Information section 10C for implementation details). Comparing the data state with the $|1\rangle_L$ state, there is a violation

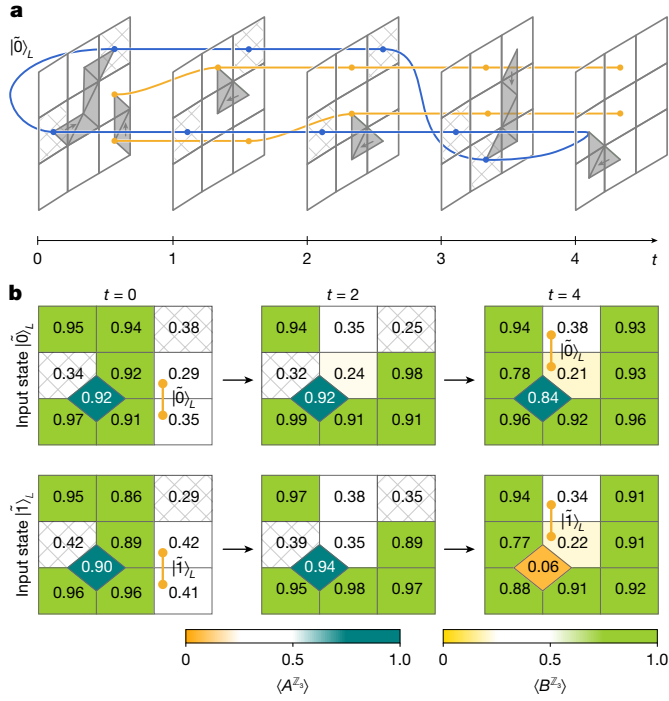


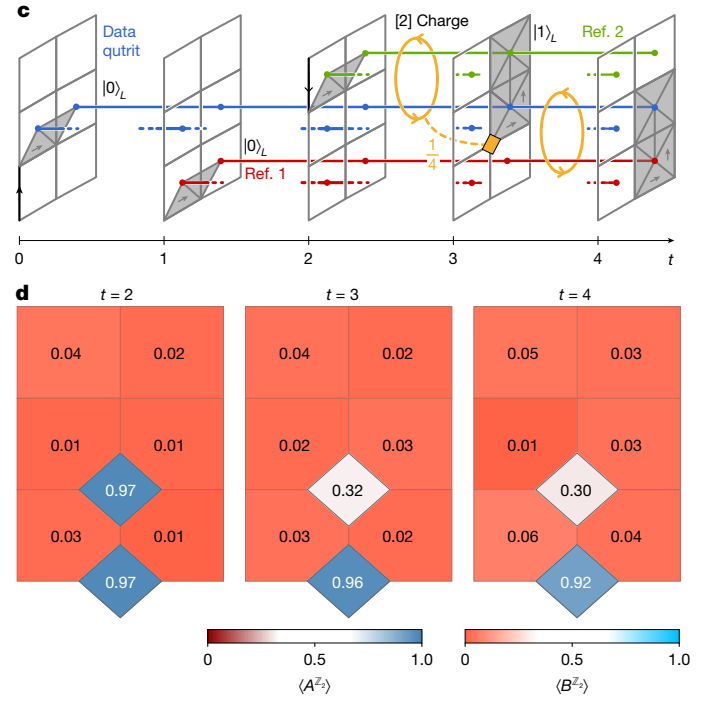
Fig. 4 | Logical $\mathcal{X}$ -basis and $\mathcal{Z}$ -basis measurements. **a**, For logical $\mathcal{X}$ -basis measurement, a symmetric ‘measurement’ flux pair (blue worldlines), created from the vacuum in the state $|\tilde{0}\rangle_L$, is braided around an endpoint of the flux pair of the input qutrit (yellow worldlines) and then fused. The fusion outcome diagnoses the input qutrit state: annihilation to the vacuum corresponds to $|\tilde{0}\rangle_L$, whereas a remnant charge anyon indicates the $\{|\tilde{1}\rangle_L, |\tilde{2}\rangle_L\}$ subspace. **b**, Experimental verification of the logical $\mathcal{X}$ -basis measurement protocol at $t = 0, 2, 4$. Top row indicates input state $|\tilde{0}\rangle_L$, and bottom row indicates $|\tilde{1}\rangle_L$. Local stabilizers remain insensitive to the logical content of the qutrits, with B^{z_2} values near $1/3$ at all flux endpoints (average/maximum standard error $0.034/0.059$). At fusion ($t = 4$), the vertex projector A^{z_2} yields $0.84(4)$ for $|\tilde{0}\rangle_L$, consistent with vacuum annihilation, and $0.06(3)$ for $|\tilde{1}\rangle_L$, indicating a remnant charge anyon. **c**, Schematic of the logical $\mathcal{Z}$ -basis measurement protocol. **d**, Experimental demonstration of the logical $\mathcal{Z}$ -basis measurement protocol. Local plaquette projectors B^{z_2} remain near zero at $t = 2, 3, 4$, confirming persistent qutrit endpoints (average/maximum error $0.014/0.021$). At $t = 3$, the vertex projector A^{z_2} at the first fusion site yields $0.32(4)$, consistent with the expected value $1/4$; at $t = 4$, A^{z_2} at the second fusion site is $0.92(2)$, indicating deterministic annihilation to the vacuum.

in the vertex projector A^{z_2} , where the $[2]$ charge pair is fused at $t = 3$ (Fig. 4d), indicating that the data qutrit is not in the $|\tilde{1}\rangle_L$ state. Finally, comparing the data state and the $|\tilde{0}\rangle_L$ state, there is no additional violation in the A^{z_2} vertex projector (at $t = 4$ in Fig. 4d), supporting that the data flux pair is in the $|\tilde{0}\rangle_L$ state. This comparison measurement between logical states also allows us to transition from the absolute encoding (equation (4)) to the relative encoding.

Bureau of standards

To define the relative encoding, we require a bureau of standards, as mentioned in section ‘Logical qutrit encoding’. Here we demonstrate the preparation of states in the $\mathcal{Z}$ -basis bureau of standards (the procedure is shown schematically in Extended Data Fig. 6a). Three charge-neutral $|\tilde{0}\rangle_L$ states are instantiated by pair-creation from the vacuum. Without loss of generality, we designate the lowest C_2 flux pair as the reference logical $|\tilde{0}\rangle_L$ state, and define all other logical states relative to this reference. Then, $\mathcal{Z}$ -basis comparison measurements are performed among the three possible pairings. By post-selecting on the $[-]$ fusion channel using mid-circuit measurements, indicating each of the three pairs occupies distinct $\mathcal{Z}$ -basis states, the reference $|\tilde{1}\rangle_L$ and $|\tilde{2}\rangle_L$ state are designated. This completes the preparation of the $\mathcal{Z}$ -basis bureau of standards, certified by the inter- and intra-pair correlations (Extended Data Fig. 6b–d).

In the experiment, for simplicity, we have post-selected for the $[-]$ charge outcome for all three $\mathcal{Z}$ -basis measurements to obtain all three logical states. However, post-selection is not necessary. If the



Worldlines on the 3×2 lattice for a data qutrit (blue) and two reference qutrits: reference 1 (red) in $|\tilde{0}\rangle_L$ and reference 2 (green) in $|\tilde{1}\rangle_L$. A $[2]$ charge pair (yellow) encircles the data and reference-2 qutrits between $t = 2$ and $t = 3$; because the logical states differ, the pair annihilates to the vacuum with probability $1/4$ or leaves a remnant charge with probability $3/4$. Between $t = 3$ and $t = 4$, a second $[2]$ charge pair encircles the data and reference-1 qutrits, both in $|\tilde{0}\rangle_L$, resulting in deterministic vacuum annihilation. **d**, Experimental demonstration of the logical $\mathcal{Z}$ -basis measurement protocol. Local plaquette projectors B^{z_2} remain near zero at $t = 2, 3, 4$, confirming persistent qutrit endpoints (average/maximum error $0.014/0.021$). At $t = 3$, the vertex projector A^{z_2} at the first fusion site yields $0.32(4)$, consistent with the expected value $1/4$; at $t = 4$, A^{z_2} at the second fusion site is $0.92(2)$, indicating deterministic annihilation to the vacuum.

measurement outcome is not the $[-]$ fusion channel but the vacuum fusion channel, that implies we have partially prepared the bureau of standards (we documented the other possible outcomes in Extended Data Fig. 6c). We can repeat until success: create fresh C_2 flux pair in the $|\tilde{0}\rangle_L$ state and compare until we obtain the $[-]$ fusion outcome to prepare the full $\mathcal{Z}$ -basis reference states.

Demonstrating magic

To demonstrate the power of the S_3 universal gate set that we have constructed, we exhibit a simple procedure to create a magic state using topological operations. Starting with the basis state $|\tilde{1}\rangle_L \propto |\tilde{0}\rangle_L + \omega |\tilde{1}\rangle_L + \bar{\omega} |\tilde{2}\rangle_L$, we use the $\mathcal{Z}$ -basis measurement primitive to find this in ‘not $|\tilde{2}\rangle_L$ ’, collapsing the state into $|\tilde{0}\rangle_L + \omega |\tilde{1}\rangle_L$ (see Fig. 5 for the protocol). This gives a minimal demonstration of how measurement on qutrits enables the creation of non-Clifford states using only topological operations.

To certify the magic state, we measure the state in both $\mathcal{Z}$ - and $\mathcal{X}$ -basis. First, in the $\mathcal{Z}$ -basis, the resulting state has equal weights only on the $|\tilde{0}\rangle_L, |\tilde{1}\rangle_L$ basis states as expected (Fig. 5c). For the $\mathcal{X}$ -basis, we use the Achilles’ heel of the absolute encoding to our advantage: by interrogating the logical charge violations at the origin, we can infer the $\mathcal{X}$ -basis measurement outcome (see Table 1 for theory prediction and Methods, section ‘Charge measurement at origin’); in particular, the charge measurement at the origin is sensitive to the relative phase ω . In Fig. 5c (inset), the experimental value of the charge violation weight

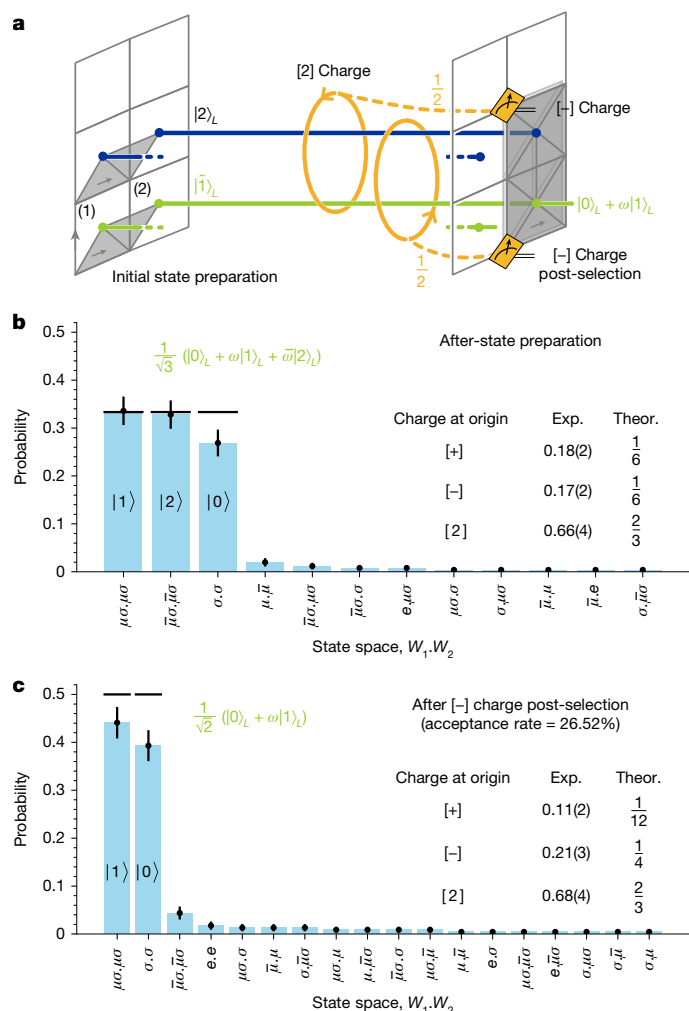


Fig. 5 | Illustrating universality by creating a topological magic state.

a, Spacetime diagram illustrating the procedure for preparing the magic state $\frac{1}{\sqrt{2}}(|0\rangle_L + \omega|1\rangle_L)$ (see Methods, section ‘Magic-state creation protocol details’ for preparation details). **b, c**, The bar plots show the joint probability of measuring the W -fluxes around plaquettes (1) and (2); **b** for the intermediate state $\frac{1}{\sqrt{2}}(|0\rangle_L + \omega|1\rangle_L + \omega|2\rangle_L)$ and **c** for the final state $\frac{1}{\sqrt{2}}(|0\rangle_L + \omega|1\rangle_L)$. Following initial state preparation, the three possible flux configurations, $\sigma, \sigma, \mu\sigma$ and $\bar{\mu}\sigma, \bar{\mu}\sigma$, which correspond to the logical basis states, occur with approximately equal probability of $1/3$. Successful preparation of the magic state is validated in the $\mathcal{Z}$ -basis by a marked increase in the probability of the σ, σ and $\mu\sigma, \mu\sigma$ flux configurations, accompanied by a suppression of the $\bar{\mu}\sigma, \bar{\mu}\sigma$ configuration; inset tables show that charge measurements at the origin in the experiments match with theoretical predictions, validating the state in the $\mathcal{X}$ -basis.

at the origin matches the theory prediction for when the state is $|0\rangle_L + \omega|1\rangle_L$, thus confirming the coherence of the magic state created.

Conclusion

In this work, we demonstrate a purely topological approach to universal quantum computation on quantum hardware. Rather than encoding information in a ground-state subspace, as in conventional topological codes, we encode logical states directly in the nonlocal fusion space of non-Abelian anyons, and we use their braiding and fusion as computational primitives. These abilities are enabled by the realization of a deconfined S_3 gauge theory, the non-Abelian topological order based on the smallest non-Abelian group. Exploiting its cyclic anyon fusion rules, we trap a single non-Abelian anyon on the torus, implement a universal logical gate set and demonstrate the

Table 1 | Theory prediction for the weights of charge violations at the origin for the magic state $|0\rangle_L + \omega|1\rangle_L$ and other possible states

State	1	[-]	[2]
$ \psi_\omega\rangle = \frac{1}{\sqrt{2}}(0\rangle_L + \omega 1\rangle_L)$	$\frac{1}{12}$	$\frac{1}{4}$	$\frac{2}{3}$
$ \psi_+\rangle = \frac{1}{\sqrt{2}}(0\rangle_L + 1\rangle_L)$	$\frac{1}{3}$	0	$\frac{2}{3}$
$\rho = p 0\rangle\langle 0 _L + (1-p) 1\rangle\langle 1 _L$	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{2}{3}$

The state without a relative phase ($|\psi_+\rangle$) and the completely decohered state (ρ) have distinct weights of charge violations.

topological preparation of a magic state. These experimental demonstrations are also enabled by theoretical advances: (1) we showed how a two-anyon encoding is sufficient in contrast to proposed four-anyon encodings^{46,47}; (2) we developed coherent moving protocols to move non-Abelian fluxes deterministically; and (3) we showed that the logical encoding based on bureau of standards can be described in the fusion-tree perspective. Together, these results establish non-Abelian anyons as programmable and computationally powerful degrees of freedom, while the solvability of S_3 supports scalable preparation and control.

The realization of these abilities reflects a convergence of advances in quantum hardware and theoretical understanding of non-Abelian topological phases. Improvements in qubit number and coherence now enable the preparation and manipulation of complex topological states, whereas recent theoretical progress in lattice-scale anyon creation^{19–21}, encoding⁴⁷ and coherent moving (C.F.B.L. et al., manuscript in preparation) establish a setting in which non-Abelian anyons can be controlled rather than merely detected.

Looking ahead, encoding logical information directly in non-Abelian anyonic excitations opens new directions for fault-tolerant quantum computation. Owing to their close similarity to the surface code, non-Abelian quantum double models offer a promising route to extending surface code techniques while enabling new abilities, including universal topological gate sets. As our gate set is topological, the operations are protected as long as one is in the phase. Stabilizing the topological phase requires active error correction, which is beyond the scope of the present work, although a recent work proved a finite decoding threshold even with imperfect measurement during braiding³⁷. Numerical works indicate thresholds may be comparable to those of the surface code^{53,54}, and a systematic comparison between ground-state and fusion-space encoding is an exciting open direction. Our proof-of-principle results lead the way for more detailed resource estimates for the non-Abelian architecture, which will be crucial to evaluate how non-Abelian computing compares with more conventional stabilizer code schemes.

More generally, our work highlights how finite-group non-Abelian topological orders offer a compelling balance between computational power and experimental accessibility. Although more exotic phases may admit braid-only universality, solvable gauge theories such as S_3 occupy a ‘sweet spot’: they are sufficiently rich to support universal computation, yet sufficiently structured to admit efficient preparation and coherent anyon transport. S_3 quantum double provides a versatile platform for studying both scalable quantum computation beyond the stabilizer paradigm and new regimes of real-time dynamics of lattice gauge theories.

Online content

Any methods, additional references, Nature Portfolio reporting summaries, source data, extended data, supplementary information, acknowledgements, peer review information; details of author contributions

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Ground-state preparation

We prepare the ground state with trivial flux along both non-contractible loops of the torus. At the end of the preparation circuit, we non-destructively measure an implicitly prepared $S_A^{\mathbb{Z}_2}$ stabilizer together with an adjacent $S_B^{\mathbb{Z}_2}$ stabilizer. Empirically, we find that, in the event of a state-preparation error, these stabilizers are more likely to be violated³⁰. We, therefore, discard any experimental shot in which either of the two non-destructive measurement outcomes is -1 . Furthermore, we exploit the qutrit encoding—implemented using two physical qubits per qutrit—to detect errors that take qutrits outside the computational subspace; this serves as an independent heralding condition (Supplementary Information section 4).

Before destructive read-out of the edge qudits, a circuit barrier is inserted to ensure complete creation of the ground state. Expectation values of stabilizer and logical projectors are obtained from destructive measurements of each qutrit–qubit pair in the following bases:

- Z and $\mathcal{Z}$, yielding $B^{\mathbb{Z}_2}$, $B^{\mathbb{Z}_3}$ and horizontal and vertical logical Z projectors;
- Z and $\mathcal{X}$, yielding $A^{\mathbb{Z}_3}$ and $B^{\mathbb{Z}_2}$ projectors; and
- X and $\mathcal{C}$, yielding $A^{\mathbb{Z}_2}$ projectors.

Mean values of logical Z projectors of the S_3 quantum double are obtained by averaging across rows for the horizontal logical operators $Z_{\text{hor}}^{\mathbb{Z}_2}$ and $Z_{\text{hor}}^{\mathbb{Z}_3}$, and across columns for the vertical logical operators $Z_{\text{vert}}^{\mathbb{Z}_2}$ and $Z_{\text{vert}}^{\mathbb{Z}_3}$. After applying all heralding criteria, approximately 24% of shots are discarded. For the smaller 3×2 lattice, shown in Extended Data Fig. 1, the corresponding discard rate is about 13%.

For comparison, we also prepare the ground state using an implementation with mid-circuit measurements and feedforward (Supplementary Information section 6B1), which yields slightly lower fidelities (Extended Data Fig. 2).

Pull-through gate

First, we show that the action of flux braiding (conjugation by the total flux) reproduces the qutrit action of the pull-through gate as specified in the main text. We start with the state $|\tilde{0}\rangle_L|0\rangle_L = (|\sigma, \sigma\rangle + |\mu\sigma, \mu\sigma\rangle + |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle)|\sigma, \sigma\rangle$. At the level of the flux, the fluxes are conjugated by the total flux participating in the braiding:

$$\begin{aligned} & U(|\sigma, \sigma\rangle + |\mu\sigma, \mu\sigma\rangle + |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle)|\sigma, \sigma\rangle \\ &= |\sigma, \sigma\rangle|\sigma, \sigma\rangle + |\mu\sigma, \mu\sigma\rangle|\bar{\mu}\sigma, \bar{\mu}\sigma\rangle + |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle|\mu\sigma, \mu\sigma\rangle \\ &= |0\rangle_L|0\rangle_L + |1\rangle_L|2\rangle_L + |2\rangle_L|1\rangle_L. \end{aligned} \quad (8)$$

In implementing the pull-through gate, as well as all other experiments presented in this work, logical qutrit states are prepared in the $\mathcal{X}$ -basis (that is, $|\star\rangle_L$) using an ancilla qutrit (Supplementary Information section 10A). Given the constrained budget of available qutrits (56 in total), ancilla qutrits are obtained by using a qubit-reuse technique, in which physical qubits that are no longer acted on in subsequent stages of the circuit are measured early and reset for use as ancillas.

After successful preparation, the ancilla qutrit is disentangled and ideally returns to the qutrit zero state. Measurement outcomes corresponding to a nonzero ancilla state—predominantly arising from memory errors in deeper circuits (Extended Data Table 1)—are therefore used to herald unsuccessful logical qutrit preparation, and the corresponding shots are discarded.

The correlator $\langle W_{p_1}^{\mathbb{Z}_3} W_{p_2}^{\mathbb{Z}_3} \rangle$, which we use to quantify correlations, ensures that

- when the control global flux is σ , the target is σ ;
- when the control global flux is $\mu\sigma$, the target is $\bar{\mu}\sigma$; and
- when the control global flux is $\bar{\mu}\sigma$, the target is $\mu\sigma$.

See Extended Data Fig. 3 for further data on non-trivial $\mathcal{Z}$ correlators.

Magic-state creation

Magic-state creation protocol details. The protocol begins by initializing the lower flux pair to be in the input state $|\hat{1}\rangle_L = \frac{1}{\sqrt{3}}(|0\rangle_L + \omega|1\rangle_L + \bar{\omega}|2\rangle_L)$ and the middle flux pair with the reference state $|2\rangle_L$. Two $\mathcal{Z}$ -basis comparison measurements with respect to $|2\rangle_L$ are performed using $[2]$ charges. After braiding the first $[2]$ charge pair around the end of both flux pairs, we post-select on the $[-]$ fusion outcome. This implies that the measured state is orthogonal to the $|2\rangle_L$ state, so that the state is projected to the orthogonal subspace. To ensure the correct coherence of the state in the subspace spanned by $|0\rangle_L$ and $|1\rangle_L$, we need to ensure no net charge is extracted from the flux pair that constitutes the logical state. With only one round of $[2]$ charge braiding, there is one local $[-]$ charge in the system, and by neutrality, the flux pair is storing another $[-]$ charge nonlocally. Therefore, we repeat the $\mathcal{Z}$ comparison measurement and wind another $[2]$ charge pair and again post-select on the $[-]$ fusion outcome, such that there are two local $[-]$ charges in the system; this ensures that the charge content of the resulting state is correct such that we obtain the state $|0\rangle_L + \omega|1\rangle_L$ (Fig. 5, bottom). See Extended Data Fig. 7 for further data and Supplementary Information section 10E for circuit-level implementation details.

Charge measurement at origin. Recall from section ‘ $\mathcal{X}$ -basis measurement’ that the $\mathcal{X}$ -basis measurement determines whether the state is $|\tilde{0}\rangle_L$ or in the orthogonal subspace by measuring the $[2]$ charge of the flux pair. In the absolute logical state, the origin contains the charge information of the logical state. This is due to the neutrality condition, the charges at the origin are exactly the partners of the charges shared by the flux pairs, such that fusing all anyons will return to vacuum. Therefore, by interrogating the logical charge violations at the origin, we can infer the $\mathcal{X}$ -basis measurement outcome.

The prediction for the expected charge violation weights for the $|0\rangle_L + \omega|1\rangle_L$ magic state and other possible states is shown in Table 1.

Data availability

The data generated in this study are available at Zenodo⁵⁵ (<https://doi.org/10.5281/zenodo.18054264>) under open access.

Code availability

The code used for numerical simulations is available at Zenodo⁵⁵ (<https://doi.org/10.5281/zenodo.18054264>).

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Author contributions C.F.B.L. and M.I. wrote the code generating the circuits for all experiments. The experiment was built and carried out by D.G., M.M., P.E.S. and M.D.U. The data analysis and interpretation were done by C.F.B.L., M.I., A.L., N.T., H.D., A.V. and R.V. C.F.B.L., A.L., N.T., A.V. and R.V. contributed to the ideation, theory and experiment design, including the definition of the operators for creation, coherent moving, and fusion of anyons; in particular, C.F.B.L., A.L. and R.V. contributed to the theory for the bureau of standards for the relative logical encoding. C.F.B.L. drafted the initial paper, which was refined by contributions from all authors, especially M.I., A.L., N.T., A.V. and R.V.; C.F.B.L., M.I. and A.L. contributed to the Supplementary Information; in particular, C.F.B.L. contributed to the protocol implementations, M.I. contributed to state preparation protocols and circuit compilation techniques, and A.L. contributed to the bureau of standards theory.

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Competing interests H.D. is a shareholder of Quantinuum. All other authors declare no competing interests.

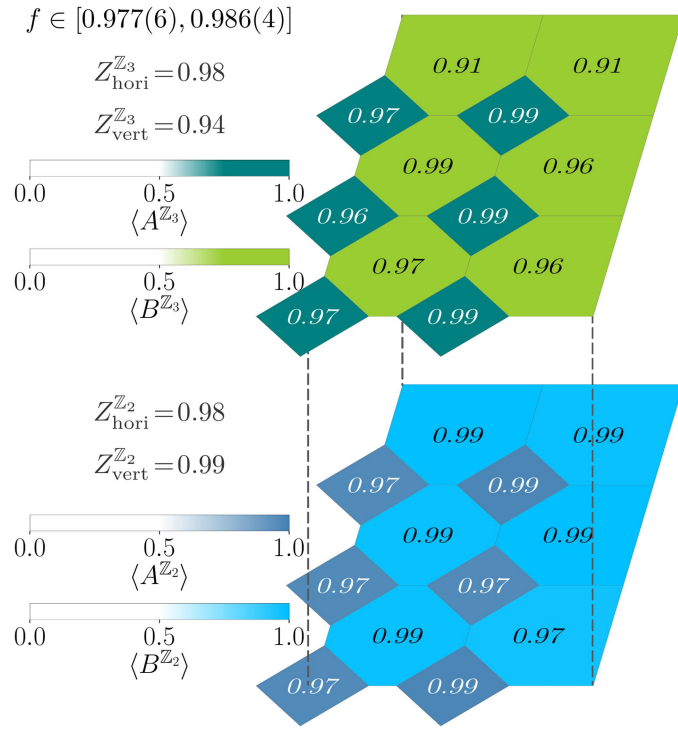
Additional information

Supplementary information The online version contains supplementary material available at <https://doi.org/10.1038/s41586-026-10709-y>.

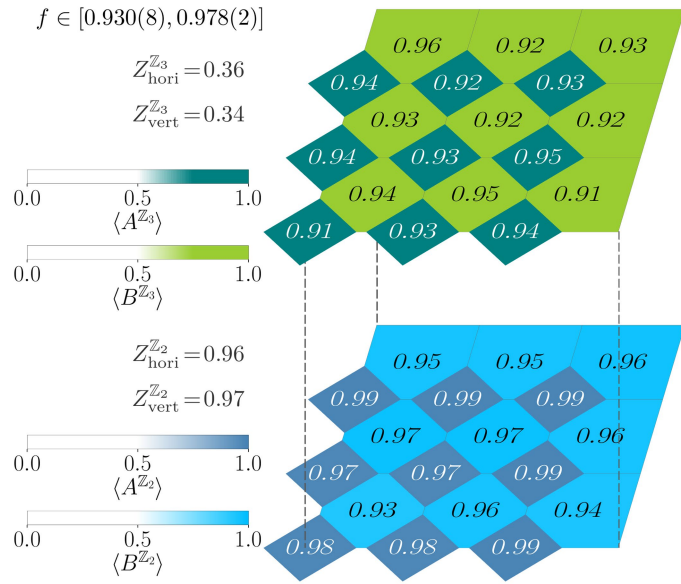
Correspondence and requests for materials should be addressed to Chiu Fan Bowen Lo, Ruben Verresen or Mohsin Iqbal.

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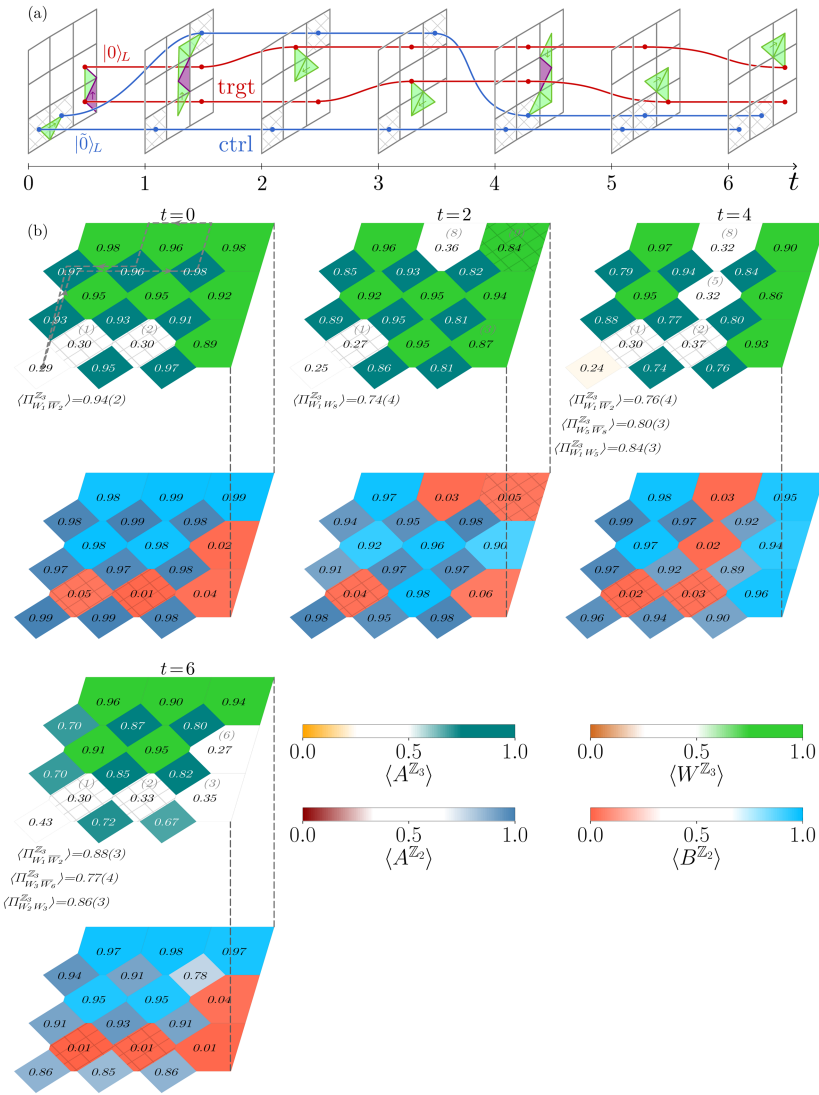
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Extended Data Fig. 1 | Ground state on the 3×2 lattice. Experimentally measured values of local and Z-logical projectors for a state prepared via a fully unitary protocol (Supplementary Information Section S6B2). The $\mathbb{Z}_3$ toric code is first prepared unitarily, followed by the unitary gauging of the $\mathbb{Z}_2$ charge-conjugation symmetry. The per-qudit fidelity of the resulting state is denoted by f . The average (maximum) standard error is 0.018 (0.034) for local projectors and 0.018 (0.031) for Z-logical projectors.

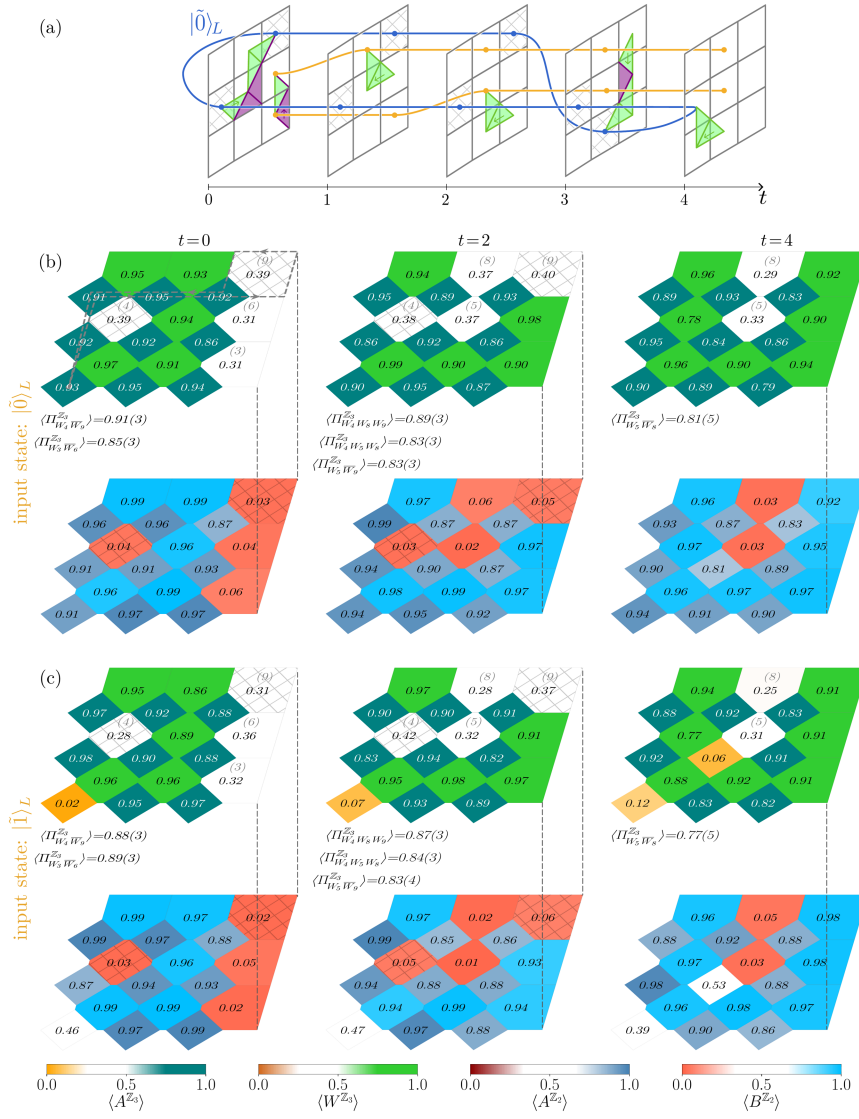


Extended Data Fig. 2 | Ground state preparation using mid-circuit measurements and feed-forward. Expectation values of local and Z-logical projectors are shown. The $\mathbb{Z}_3$ toric-code ground state is prepared unitarily by enforcing plaquette stabilizers (in contrast to the vertex stabilizers used in the main-text experiments), yielding a +1 eigenstate of the $\mathbb{Z}_3$ X-logical operators. This is verified by $Z_{L,\text{hor}}^{\mathbb{Z}_3}$ and $Z_{L,\text{vert}}^{\mathbb{Z}_3}$ values near 1/3. The $\mathbb{Z}_2$ charge-conjugation symmetry is then gauged using an adaptive circuit (Supplementary Information Section S6B1) with mid-circuit measurements and conditional error correction based on stabilizer outcomes. Given that Π is the projector onto the ground-state manifold, $f = \text{Tr}(\rho\Pi)^{1/18}$ denotes the per-qudit fidelity of the prepared state ρ . The average (maximum) standard error for local projectors is 0.012 (0.017), and for Z-logical projectors the average and maximum standard errors are 0.019 and 0.028, respectively.



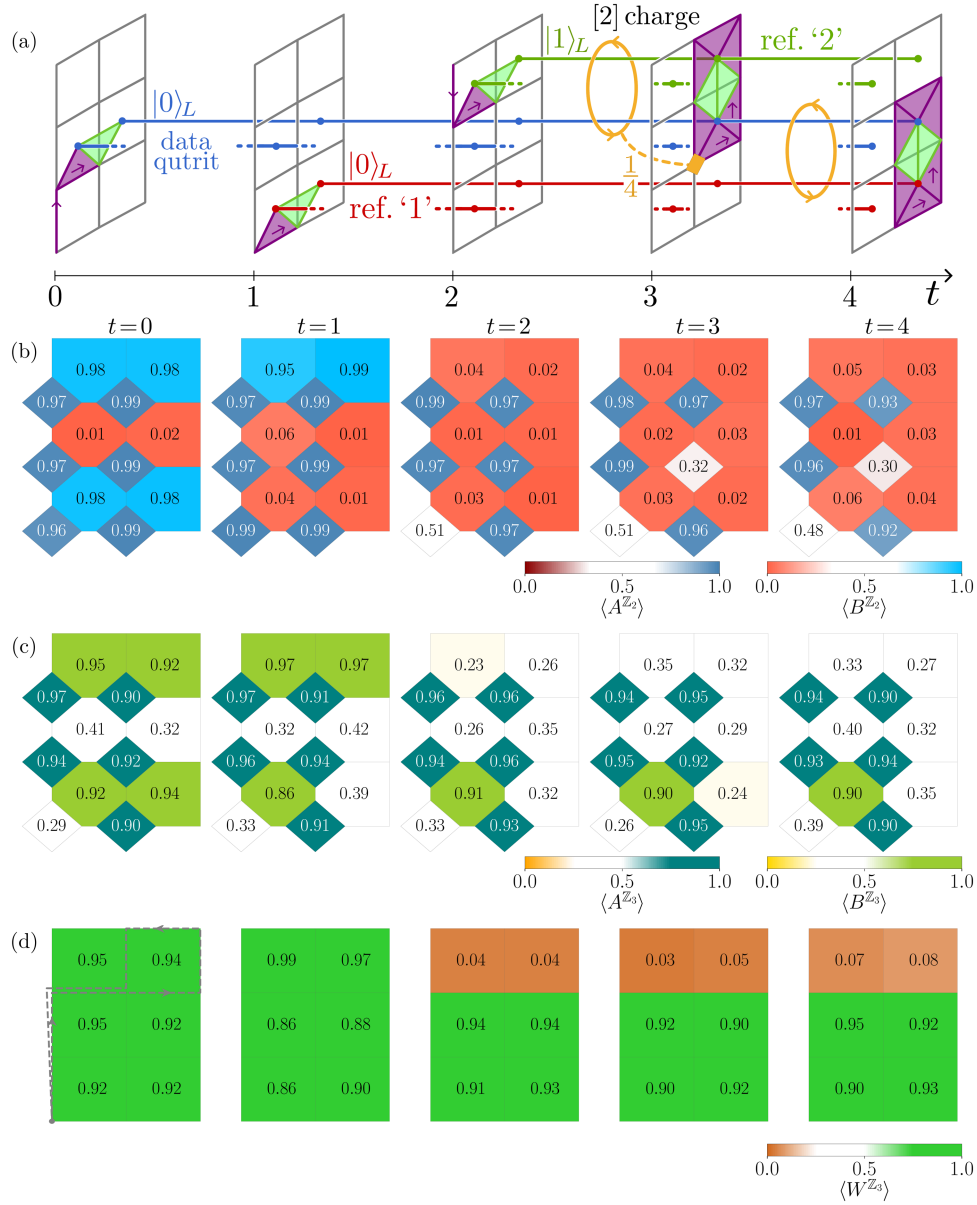
Extended Data Fig. 3 | Pull-through gate protocol. (a) A control qutrit, initialized in the $|0\rangle_L$ state, is shown with blue worldlines; plaquettes hatched by $\times$ indicate the locations of its endpoints at each time slice. The control qutrit is moved around a target qutrit initialized in the $|0\rangle_L$ state and shown with red worldlines. The ribbon operator implementing the motion is constructed from elementary triangle operators (Supplementary Information Section S1D). (b) Measured expectation values of the local projectors $A^{\mathbb{Z}_2}$, $A^{\mathbb{Z}_3}$, $B^{\mathbb{Z}_2}$ and $B^{\mathbb{Z}_3}$,

together with the nonlocal projector $W^{\mathbb{Z}_3}$, shown at different stages of the protocol. For all local and $W^{\mathbb{Z}_3}$ measurements, the average standard error is 0.023, with a maximum of 0.052. Also shown are the W -flux correlators $\Pi_{W_{p_1}W_{p_2}}^{\mathbb{Z}_3}$, which certify inter- and intra-pair delocalized $\mathbb{Z}$ correlations of the control and target qutrits throughout the protocol. In the noise-free limit these correlators equal 1; the measured values are consistent with this expectation, indicating high-fidelity state preparation.



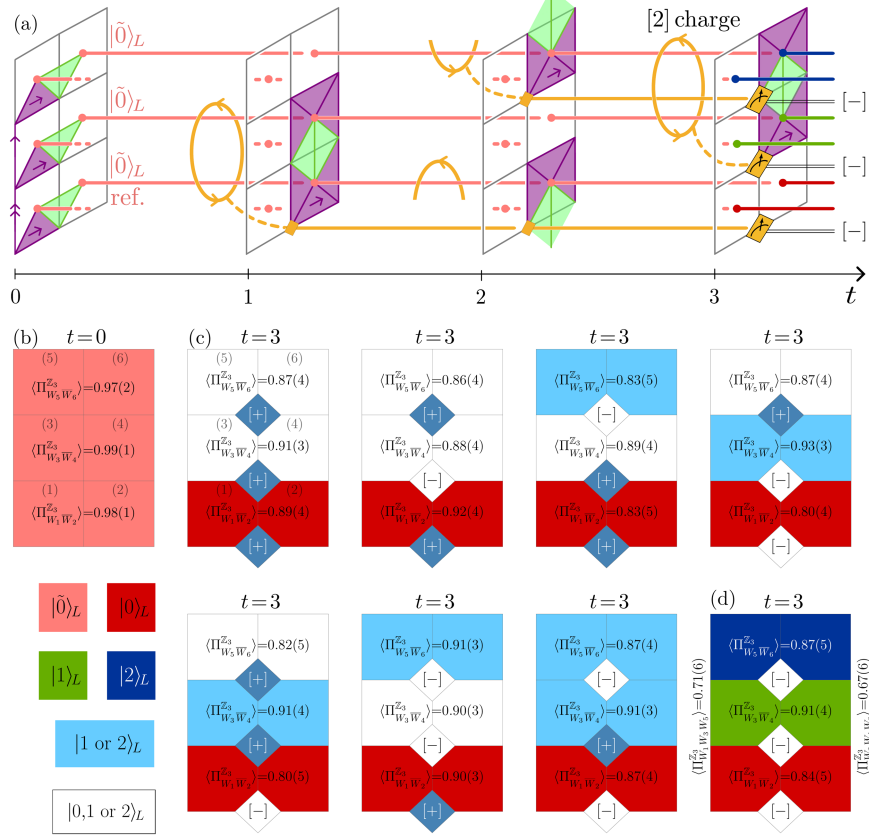
Extended Data Fig. 4 | $\mathcal{X}$ -basis measurement protocol. (a) Spacetime diagram of a ‘measurement’ qutrit braiding around the flux endpoint of an input qutrit. Worldlines are shown in blue (measurement) and yellow (input). Plaquettes marked by $\times$ denote the measurement qutrit flux endpoints on the 3×3 lattice. (b,c) Measurement results for the $\mathbb{Z}_2$ and $\mathbb{Z}_3$ projectors. The input qutrit is prepared in the $|\tilde{0}\rangle_L$ state in panel (b) and in the $|\tilde{1}\rangle_L$ state in panel (c). In each panel, the lower grid shows the local A^{Z_2} and B^{Z_2} projectors, while the upper grid shows the local A^{Z_3} and nonlocal $\Pi_W^{Z_3}$ projectors (see Supplementary Information Section S2C for precise W -flux paths). The vertex projectors A^{Z_2}

and A^{Z_3} remain close to +1 at all sites except the origin and the remnant charge site, while the local B^{Z_2} plaquette stabilizers remain near 0 at qutrit endpoints and reveal no logical information. The nonlocal W -flux projectors, which distinguish the absolute logical states $|0\rangle_L, |1\rangle_L$, and $|2\rangle_L$, remain near $1/3$ due to the superposition of these states. Two-body W -flux correlators $\Pi_{W_{p_1}W_{p_2}}^{Z_3}$ are close to 1, indicating that the prepared state closely approximates the noiseless ideal. The local $\langle A^{Z_2}, A^{Z_3}, B^{Z_2} \rangle$ and nonlocal $\langle W^{Z_3} \rangle$ projectors have an average (maximum) standard error of 0.028 (0.072).



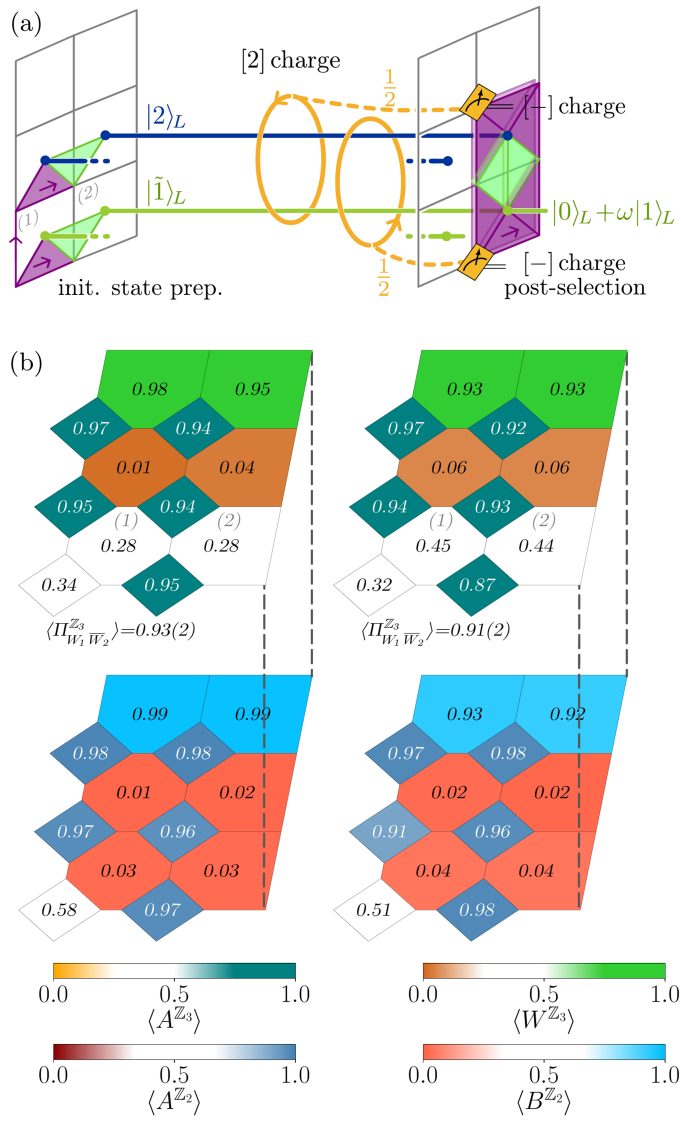
Extended Data Fig. 5 | Z-basis measurement protocol. (a) A data qutrit (blue) and two reference qutrits, labeled '1' (red) and '2' (green), are initialized in $|0\rangle_L$, $|0\rangle_L$, and $|1\rangle_L$ states at $t=0$, $t=1$, and $t=2$, respectively. [2] charge pairs loop around the endpoints of data-reference qutrit pairs ($t=2$ - $t=3$ for data and reference '2' qutrits, $t=3$ - $t=4$ for data and reference '1' qutrits). (b) Expectation values of local $\mathbb{Z}_2$ vertex ($A^{\mathbb{Z}_2}$) and plaquette ($B^{\mathbb{Z}_2}$) projectors. $B^{\mathbb{Z}_2} \rightarrow 0$ upon logical qutrits initialization indicate presence of logical qutrit endpoints. The vertex projector at the first [2] charge pair fusion site yields a value $\approx 1/4$, confirming remnant $[-]$ charge. Subsequent [2] charge pair fusion between $t=3$

and $t=4$ leaves $A^{\mathbb{Z}_2}$ values unchanged, indicating no additional charge accumulation. (c) $\mathbb{Z}_3$ vertex ($A^{\mathbb{Z}_3}$) and plaquette ($B^{\mathbb{Z}_3}$) projector expectation values. $A^{\mathbb{Z}_3}$ remains constant (≈ 1 , except $\approx 1/3$ at origin). $B^{\mathbb{Z}_3}$ projectors are measured $\approx 1/3$ at all qutrit initialization sites provide no information about the internal state of the qutrits. (d) Nonlocal $W^{\mathbb{Z}_3}$ projector values distinguish internal states: $W^{\mathbb{Z}_3}$ for $|0\rangle_L$ state of data and reference '1' qutrits and ≈ 0 for $|1\rangle_L$ state of reference '2' qutrit. Across all local and nonlocal projectors, the average standard error is 0.030, with a maximum standard error of 0.064.



Extended Data Fig. 6 | Z bureau of standards protocol. (a) Schematic of the protocol used to define a logical qutrit in the Z -basis via relative flux comparisons. Three C_2 flux pairs are initialized in the $|\tilde{0}\rangle_L$ state, with the lowest pair chosen as the reference $|\tilde{0}\rangle_L$. The logical states of the remaining pairs are determined through three comparison measurements implemented by braiding a $[2]$ charge pair. (b–d) Experimental verification of the prepared states. (b) At $t=1$, the $\Pi_{W_{p1}W_{p2}}^{Z_3}$ projectors take values close to unity, demonstrating that the prepared states closely approximate the ideally prepared (i.e., noiseless) $|\tilde{0}\rangle_L$ pairs. (c) The data is shown for different sets of

'partially' prepared basis states at $t=3$, where each set is defined by post-selection on a specific outcome of the remnant charge (A^{Z_2}) measurement. (d) The final panel at $t=3$ displays three distinct states where all three remnant charges are measured in $[-]$ states, post-selected with an acceptance rate of ~ 0.115 , comparable with the theoretical prediction of 0.125 ; the total acceptance rate (including heralding) is ~ 0.06 . Strong intra-pair and inter-pair Z correlations are confirmed by measurements of $\Pi_{W_{p1}W_{p2}}^{Z_3}$ and $\Pi_{W_{p1}W_{p2}W_{p3}}^{Z_3}$, respectively.



Extended Data Fig. 7 | Magic state preparation protocol. (a) Schematic illustrating the creation of a magic state by inducing projection via measurement (Supplementary Information Section S10E). (b) Expectation values of local projectors A^{z_2} , A^{z_3} , and B^{z_2} , and the nonlocal projector W^{z_3} after initial state preparation and magic state preparation. For plaquettes (1) and (2), the measured $\Pi_{W_1 W_2}^{z_3}$ values are 0.28 and 0.28 after initial state preparation, in close agreement with the theoretical value of $1/3$. Following the complete protocol, these values change to 0.45 and 0.44, consistent with the theoretical value of $1/2$. The high fidelities of the prepared states is further supported by measurements of $\Pi_{W_1 W_2}^{z_3}$, which are close to unity. Across all W^{z_3} and local projector measurements, the average standard error is 0.017, and the maximum standard error is 0.057.

Extended Data Table 1 | Circuit specifications

Experiment	Lattice size	No. of native one-qubit gates	No. of native two-qubit gates	Depth (two-qubit gates)
Ground state unitary prep.	3×3	429	238	43
Ground state measurement-based prep.	3×3	545	274	50
Single anyon state prep.	3×3	453	308	76
$\mathcal{Z}$ -basis measurement at $t = 4$	3×2	423	269	74
$\mathcal{Z}$ -basis bureau of standards at $t = 3$	3×2	646	437	107
Magic state preparation	3×2	507	346	134
$\mathcal{X}$ -basis measurement at $t = 5$	3×3	938	744	232
Pull-through gate at $t = 6$	3×3	1,008	845	307

Gate counts and depths for the circuits used in each experiment. Destructive measurement at the end of the circuit was performed to compute the plaquette projectors. Furthermore, each circuit was optimized using local qutrit basis transformations (Supplementary Information Section S5). For reference, the execution time for a single shot of the $\mathcal{X}$ -basis measurement at $t = 5$ was approximately 5.0 s, and for the pull-through gate experiment at $t = 6$ it was approximately 5.9 s on Quantinuum’s H2-1 quantum computer.

Supplementary information

Universal gates from braiding and fusing anyons on quantum hardware

In the format provided by the
authors and unedited

Supplementary Information

Note: the section labels (with Roman numerals) refer to sections in the main text as follows:

- Section II: The S_3 quantum double
- Section III: State preparation
 - Section III A: Ground-state preparation
 - Section III B: A single non-Abelian anyon on the torus
- Section IV: Universal topological gate set
 - Section IV A: Logical qutrit encoding
 - Section IV B: Pull-through gate
 - Section IV C: $\mathcal{X}$ -basis measurement
 - Section IV D: $\mathcal{Z}$ -basis measurement
 - Section IV E: Bureau of standards
- Section V: Demonstrating magic

S1. BASICS OF S_3 QUANTUM DOUBLE MODEL

For more details, we refer the reader to the comprehensive pedagogical notes in Refs. S1 and S2.

A. Review of S_3 group theory and representation theory

We use the group presentation for $S_3 = \langle \mu, \sigma | \mu^3 = \sigma^2 = e, \mu\sigma = \sigma\mu^{-1} \rangle$, where μ and σ corresponds to the 120° rotation and reflection of an equilateral triangle. Another common notation is to recognize that S_3 is a symmetric group on three elements, so group elements are permutations that can be represented in terms of cycles. The identification is $\mu = (123)$, $\sigma = (23)$ (the remaining group elements can be determined from these two generators).

There are three conjugacy classes that partitions S_3 : $C_1 = \{e\}$, $C_2 = \{\sigma, \mu\sigma, \bar{\mu}\sigma\}$, $C_3 = \{\mu, \bar{\mu}\}$.

There are three irreducible representations (irreps) of S_3 : $[+]$, $[-]$, $[2]$. The trivial representation $[+]$ sends every group element to 1. The sign representation $[-]$ sends group element $g \mapsto \text{sgn}(g)$, where sgn gives the sign of the permutation; so even permutations (e.g. C_3 group elements) are mapped to $+1$, whereas odd permutations (e.g. C_2 group elements) are mapped to -1 . The standard representation $[2]$ is a two-dimensional representation, and the representation matrices are:

$$\begin{aligned} \Gamma_{[2]}(e) &= \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} & \Gamma_{[2]}(\mu) &= \begin{pmatrix} \omega & 0 \\ 0 & \bar{\omega} \end{pmatrix} & \Gamma_{[2]}(\bar{\mu}) &= \begin{pmatrix} \bar{\omega} & 0 \\ 0 & \omega \end{pmatrix} \\ \Gamma_{[2]}(\mu\sigma) &= \begin{pmatrix} 0 & \omega \\ \bar{\omega} & 0 \end{pmatrix} & \Gamma_{[2]}(\bar{\mu}\sigma) &= \begin{pmatrix} 0 & \bar{\omega} \\ \omega & 0 \end{pmatrix} & \Gamma_{[2]}(\sigma) &= \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \end{aligned}$$

where the indexing order is $\{|2+\rangle, |2-\rangle\}$ and $\omega = \exp(2\pi i/3)$.

The representation theory of the $\mathbb{Z}_3$ subgroup of S_3 also warrants a mention, as it will be important for defining some of the anyon types in S_3 quantum double. Since $\mathbb{Z}_3$ is an Abelian group, all irreps are one-dimensional. The three possible irreps are labeled by the third root of unity. The trivial representation $[1]$ maps all group elements to 1. The $[\omega]$ representation is given by the map $a \mapsto \omega^a$. The $[\bar{\omega}]$ representation is given by the map $a \mapsto (\bar{\omega})^a$.

B. Qutrit-qubit encoding

We use the qutrit-qubit encoding for the S_3 group elements, such that for a given S_3 group element basis $|g\rangle = |\mu^a \sigma^b\rangle$, it is encoded as a qutrit $|a \pmod{3}\rangle$ and a qubit $|b \pmod{2}\rangle$, where a, b are integers.

The left- and right-group multiplication operators are

$$L^g |h\rangle = |gh\rangle, \quad R^g |h\rangle = |h\bar{g}\rangle, \quad (\text{S1})$$

where $\bar{g}$ is the inverse of g . In terms of qutrit and qubit gates, the left- and right- group multiplication of S_3 generators are

$$L^\mu = \mathcal{X} \otimes I, \quad L^\sigma = \mathcal{C} \otimes X, \quad (\text{S2})$$

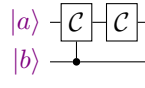
$$R^\mu = \mathcal{X}^{-Z}, \quad R^\sigma = I \otimes X, \quad (\text{S3})$$

where $\mathcal{X}$ and Z denote the qutrit clock matrices, and $\mathcal{C}$ is the qutrit charge-conjugation operator,

$$Z = \begin{pmatrix} 1 & & \\ & \omega & \\ & & \omega^2 \end{pmatrix}, \quad \mathcal{X} = \begin{pmatrix} & 1 & \\ 1 & & \\ & & 1 \end{pmatrix}, \quad \mathcal{C} = \begin{pmatrix} 1 & & \\ & 1 & \\ & & 1 \end{pmatrix}, \quad (\text{S4})$$

and where Z and X are the Pauli matrices acting on the qubit degrees of freedom.

Given an S_3 group element $g = \mu^a \sigma^b$, its inverse is $\bar{g} = \mu^{-(-1)^b a} \sigma^b$. This inverse operation is implemented by the following circuit:


(S5)

Note that taking an inverse of a S_3 group element does not change the qubit value b since inverse operation does not change its conjugacy class.

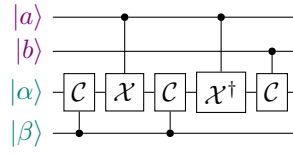
Controlled left multiplication of a group element $h = \mu^\alpha \sigma^\beta$ (in teal) by the group element $g = \mu^a \sigma^b$ (in violet) (left panel) and its inverse $\bar{g}$ (right panel) is


(S6)

Similarly, controlled right multiplication by a group element g (left panel) and its inverse $\bar{g}$ (right panel) is


(S7)

Finally, controlled conjugation sends $g \rightarrow \bar{z}gz$, where $z = \mu^a \sigma^b$ is the control (in violet) and $g = \mu^\alpha \sigma^\beta$ is the target (in teal). So the circuit is


(S8)

C. Anyon types, fusion rules, and braiding

Given a finite group G , the quantum double of G (also known as the G Drinfeld double, denoted as $\mathcal{D}(G)$) [S3], has anyon type specified by (C, χ) , where C is a conjugacy class and χ is an irreducible representation of the centralizer $Z(r)$ for a representative r in C . The 8 possible anyon types in S_3 quantum double are given below:

We denote the vacuum anyon $(C_1, [+])$ by 1 in the main text. In the following, we also simply refer to it as $[+]$.

The most relevant fusion rules are as follows:

$$[-] \times [-] = [+], \quad [2] \times [2] = [+] + [-] + [2] \quad (\text{S9})$$

Conjugacy class	Centralizer	Charge	d
C_1	$Z(e) = S_3$	$[+]$	1
		$[-]$	1
		$[2]$	2
C_2	$Z(\sigma) \cong \mathbb{Z}_2$	$[+]$	3
		$[-]$	3
C_3	$Z(\mu) \cong \mathbb{Z}_3$	$[1]$	2
		$[\omega]$	2
		$[\bar{\omega}]$	2

Table I. Conjugacy classes, centralizers, charges, and the quantum dimensions (d) of the 8 anyon types in $\mathcal{D}(S_3)$.

$$C_2 \times C_2 = [+] + [2] + C_3 + (C_3, [\omega]) + (C_3, [\bar{\omega}]), \quad C_2 \times [2] = C_2 + (C_2, [-]) \quad (\text{S10})$$

$$C_3 \times C_3 = [+] + [-] + C_3, \quad C_3 \times C_2 = C_2 + (C_2, [-]) \quad (\text{S11})$$

The full categorical data describing the S_3 quantum double, including fusion rules, the F symbols, and the R matrices, can be found in Ref. S4 and in the Supplementary Categorical Data of Ref. S5.

For braiding, it suffices to describe two scenarios: 1) braiding two fluxes and 2) braiding a flux around a charge. Braiding involving dyons can be decomposed into aforementioned scenarios. For scenario 1, when two fluxes are braided around each other, the global fluxes labels are both conjugated by the total flux. For scenario 2, the action of a flux g on a charge's internal Hilbert space is given by the representation matrix $\Gamma(g)$, where Γ is the irreducible representation of the charge.

D. Ribbon operators

1. Basic construction

The operators that create pairs of non-Abelian anyons on top of the quantum double ground state are known as *ribbon operators*. A ribbon operator $F_\rho^{A;u,u'}$ is specified by four data: (i) the path along the lattice where the ribbon acts, ρ , (ii) the anyon type $A = (C, \chi)$, where C is the conjugacy class and χ is an irreducible representation, (iii) the local internal state $u = (v, i)$ at the start, and (iv) the local internal state $u' = (v', j)$ at the end, where u, v are elements in the conjugacy class C and i, j are the matrix indices for the irreducible representation χ . In general, a ribbon operator creates vertex and plaquette violations at both ends of the ribbon.

The easiest way to express the ribbon operator $F_\rho^{A;u,u'}$ is in terms of basis elements $F_\rho^{(z,v)}$, whose action is depicted in Fig. 1b.

$$F^{A;u,u'}_\rho = \frac{|\chi|}{|\mathbf{Z}(r)|} \sum_{n \in \mathbf{Z}(r)} \Gamma_{jj'}^\chi(n) F^{(q_v n \bar{q}_{v'}, v)}, \quad A = (C, \chi), \quad u = (v, j) \quad (\text{S12})$$

Here, $r \in C$ is a representative element of the conjugacy class corresponding to the anyon in question. $\mathbf{Z}(r)$ is the centralizer of r , the subgroup of all elements in that commute with r . Furthermore, we have

$$v = q_v r \bar{q}_v, \quad v' = q_{v'} r \bar{q}_{v'} \quad (\text{S13})$$

The operator $F_\rho^{A;u,u'}$ creates two excitations at either end of ρ , but commutes with all stabilizers along the bulk of the ribbon. The anyons on either end will be of type $A = (C, \chi)$, and the internal state of each anyon will be labeled by u and u' , respectively.

Ribbon operators are constructed out of fundamental building blocks known as “triangle” operators: these are operators that act on a single edge of the lattice, and are strung together through a recursive formula to yield a full-length ribbon operator. Fig. 2 contains a table listing the eight different triangle operators, distinguished by three different characteristics: (i) whether they act on the direct or dual lattice, (ii) whether the direction of the operator is aligned with the lattice, and (iii) a property known as “local orientation” [S6].

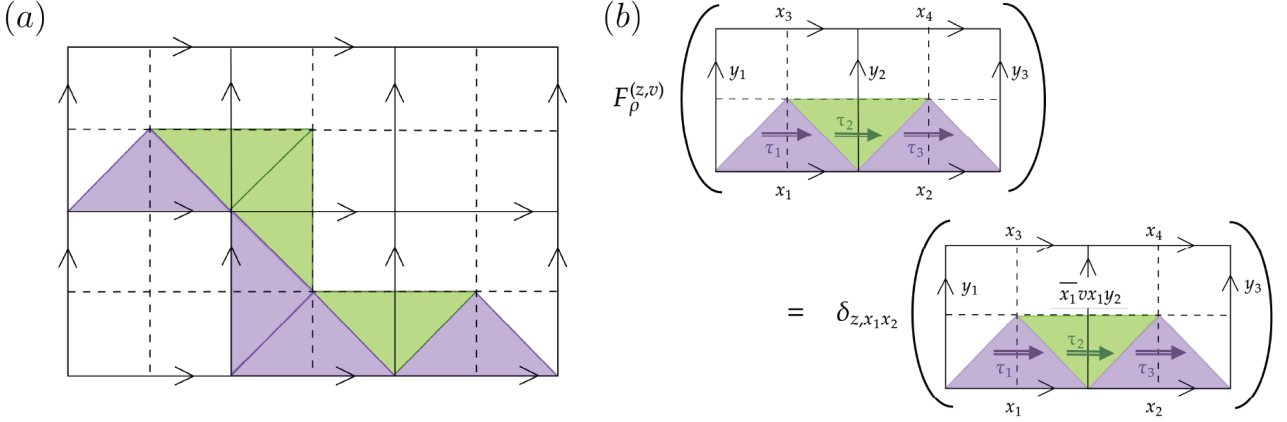


Figure 1. (a) Ribbon operators create non-Abelian anyons at their ends; they are thicker as compared to their abelian counterparts, and act both on the direct and dual edges of the lattice. (b) The action of a ribbon operator in the “group element” basis. This ribbon creates an anyon with local flux v and nonlocal flux $zv\bar{z}$, and will generically create a superposition of charges. The anyon basis ribbons are superpositions of these ribbons.

A full-length ribbon operator $F^{(z,v)}$ is built out of triangle operators recursively. Consider the three triangle ribbon pictured in Fig. 1. Denote the triangles by τ_i . We can derive the action of the full ribbon using the following formula:

$$F^{(z,v)}(\rho = \rho_1 \cup \rho_2) = \sum_{k \in G} F^{(k,v)}(\rho_1) F^{(\bar{k}z, \bar{k}vk)}(\rho_2). \quad (\text{S14})$$

Triangles τ_1 and τ_3 act on the direct lattice, are aligned with the lattice orientation, and have counter-clockwise local orientation, meaning the corresponding triangle operators $F^{(z,v)}(\tau_1)$ and $F^{(z,v)}(\tau_3)$ act by $|x\rangle \rightarrow \delta_{z,x} |x\rangle$. Triangle τ_2 acts on the dual lattice, is opposite to the lattice orientation, and has counter-clockwise local orientation, and so the corresponding operator acts as $F^{(z,v)}(\tau_2) |x\rangle = \delta_{z,e} |vx\rangle$.

We start by breaking up the ribbon into two pieces: $\rho = \tau_1 \cup (\tau_2 \cup \tau_3)$:

$$F^{(z,v)}(\rho) = \sum_{k \in G} F^{(k,v)}(\tau_1) F^{(\bar{k}z, \bar{k}vk)}(\tau_2 \cup \tau_3) = \sum_{k \in G} \delta_{x_1, k} F^{(\bar{k}z, \bar{k}vk)}(\tau_2 \cup \tau_3) = F^{(\bar{x}_1 z, \bar{x}_1 vx_1)}(\tau_2 \cup \tau_3). \quad (\text{S15})$$

Applying the recursion relation again:

$$\begin{aligned} F^{(z,v)}(\rho) &= \sum_{k \in G} F^{(k, \bar{x}_1 vx_1)}(\tau_2) F^{(\bar{k}\bar{x}_1 z, \bar{k}\bar{x}_1 vx_1 k)}(\tau_3) \\ &= \sum_{k \in G} \delta_{k, e} |(\bar{x}_1 vx_1) y_2\rangle \langle y_2 | \delta_{x_2, \bar{k}\bar{x}_1 z} \\ &= \delta_{x_1 x_2, z} |(\bar{x}_1 vx_1) y_2\rangle \langle y_2|. \end{aligned} \quad (\text{S16})$$

This is the same action quoted in Fig. 1(b).

2. Unitary ribbon

We now specialize to useful ribbons used in the main text. We observe that by summing over the local internal state on one end of the ribbon, the resulting ribbon operator is unitary. Such unitary ribbon can be decomposed into three steps: 1) ungauging, 2) finite-depth local unitary (FDLU) circuit in a simpler TO, and 3) regauging [S7]. For the ungauging step, it is a sequential unitary circuit that “unzips” along the length of the ribbon to a simpler TO. Then, the FDLU circuit is either a string operator or symmetry fractionalization pattern in the simpler TO. Finally, the regauging circuit “zips up” along the length of the ribbon back to the original TO.

The unitary ribbon for a $[2]$ charge vacuum pair is given by:

$$F^{(+, [2]); 0} = \sum_j F_{\rho}^{(+, [2]); 0, j} = \prod_{x \in \rho} Z_x^{\prod_{g < x} Z_g}. \quad (\text{S17})$$

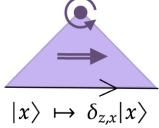
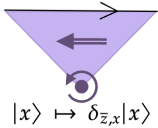
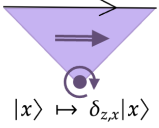
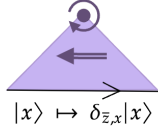
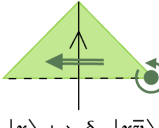
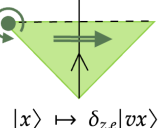
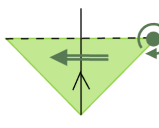
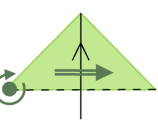
		Alignment	
Type of Triangle	Local Orientation	Aligned	Opposite
Direct	Counterclockwise	 $ x\rangle \mapsto \delta_{z,x} x\rangle$	 $ x\rangle \mapsto \delta_{z,x} x\rangle$
	Clockwise	 $ x\rangle \mapsto \delta_{z,x} x\rangle$	 $ x\rangle \mapsto \delta_{z,x} x\rangle$
Dual	Counterclockwise	 $ x\rangle \mapsto \delta_{z,e} x\bar{v}\rangle$	 $ x\rangle \mapsto \delta_{z,e} vx\rangle$
	Clockwise	 $ x\rangle \mapsto \delta_{z,e} v\bar{x}\rangle$	 $ x\rangle \mapsto \delta_{z,e} xv\rangle$

Figure 2. Definition of triangle operators, based on the type (direct or dual), alignment (aligned or opposite), and local orientation (clockwise or counterclockwise) of the triangle operator [S2]. Direct and dual edge is denoted by solid and dotted line respectively. The orientation of the edge is indicated by the single arrow; the direction of the ribbon is indicated by the double arrow. The local orientation of the triangle operator is determined by first fixing the hinge (dark circle) at the plaquette center touching the triangle operator, and then turn along the ribbon direction.

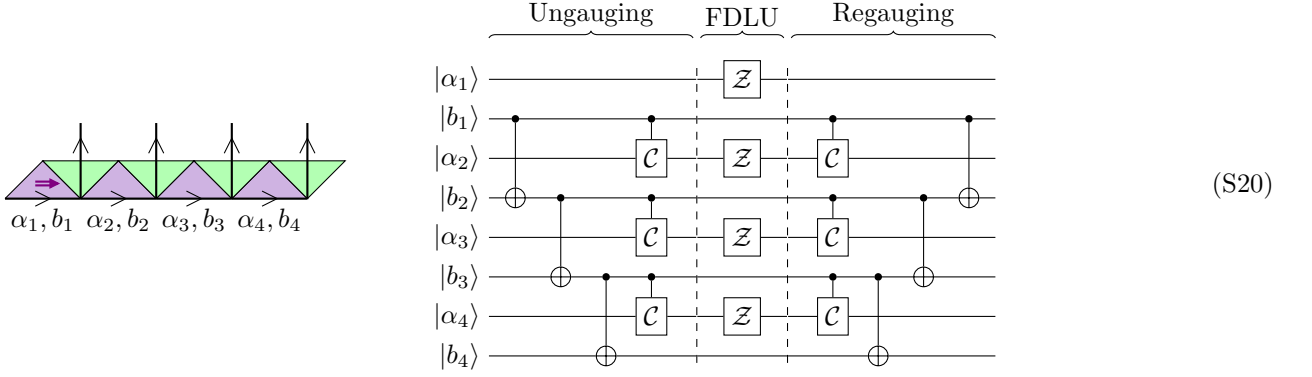
Here x, g are edges on the lattice along the ribbon ρ , and $g < x$ means that g occurs before edge x , traveling along ρ from the start of the ribbon to the end. The gate $\mathcal{Z}_x^{Z_g}$ is a controlled gate with the qutrit on edge x as the target and the qubit on edge g as the control:

$$\mathcal{Z}_x^{Z_g} |\mu^a\rangle_x |\sigma^b\rangle = \omega^{(-1)^b a} |\mu^a\rangle_x |\sigma^b\rangle, \quad a = 0, 1, 2, \quad b = 0, 1. \quad (\text{S18})$$

There can also be multiple controls on some set of edges $\{g\}$:

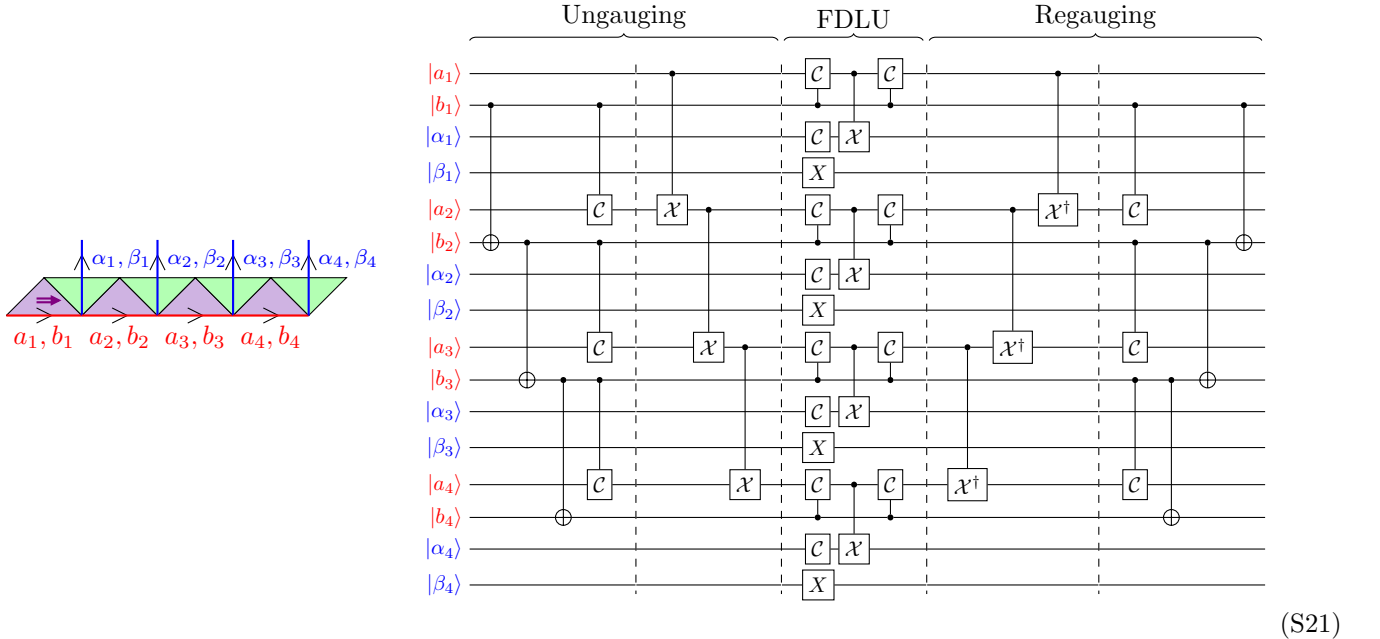
$$\mathcal{Z}_x^{\prod_g Z_g} |\mu^a\rangle_x \prod_g |\sigma^{b_g}\rangle_g = \omega^{(-1)^{\sum_g b_g a}} |\mu^a\rangle_x \prod_g |\sigma^{b_g}\rangle_g, \quad a = 0, 1, 2, \quad b_g = 0, 1. \quad (\text{S19})$$

$F^{(+,[2]):0}$ can be expressed as a circuit by using the identities in Section S1 B (where the purple arrow in the first purple triangle indicates the direction of the ribbon):



The [2] charge ribbon operator circuit above illustrate how a unitary ribbon operator can be decomposed into three steps (as divided by dotted line): 1) ungauging, 2) finite-depth local unitary (FDLU) circuit in a simpler TO, and 3) regauging [S7]. For the ungauging step, it is a sequential unitary circuit that "unzip" the direct edges to a simpler TO (in this case ungauging to $\mathbb{Z}_3$ toric code). Then, the FDLU circuit is just the usual string operator for the e anyons in $\mathbb{Z}_3$ toric code. Finally, the regauging layer is used to "zip up" the direct edges back to S_3 TO.

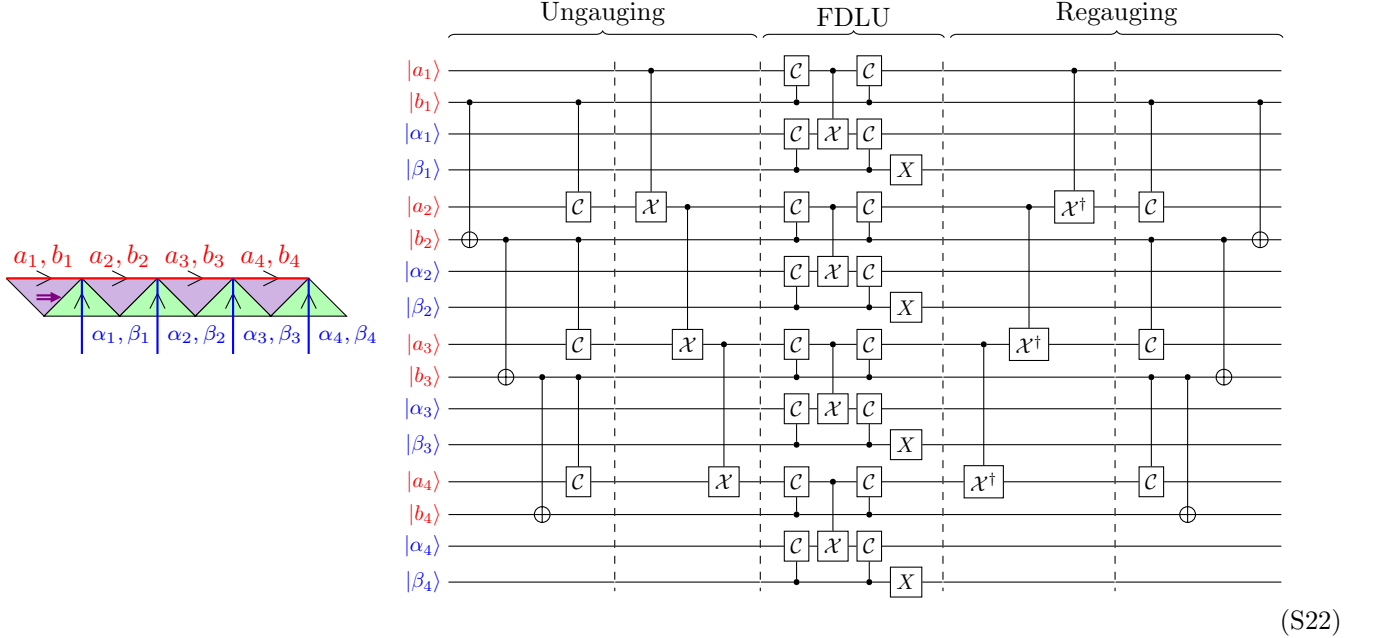
Unitary ribbons for a $[C_2]$ flux is more complicated than the [2] charge ribbon. For example, to implement the unitary ribbon $F^{C_2;\sigma} = \sum_{c \in C_2} F^{C_2;\sigma,c}$, where the local internal state is $|\sigma\rangle$ at the start of the ribbon, the circuit for a C_2 ribbon (with direction indicated by the purple arrow in the first purple triangle) with counterclockwise local orientation (the issue of local orientation will be explained below) is given by:



Note that the ungauging and regauging steps are more complicated than that of the [2] charge ribbon, because it does not stop at the level of the $\mathbb{Z}_3$ toric code but continue to the level of $\mathbb{Z}_3$ paramagnet [S7].

As a remark, to implement a C_2 ribbon where the starting local internal state is $|\mu\sigma\rangle$ and $|\bar{\mu}\sigma\rangle$, the entire circuit needed to be conjugated with $\mathcal{X}$ and $\mathcal{X}^\dagger$ respectively on the first direct edge (a_1, b_1 in this case), as it modifies the action of individual dual triangle to $L^{\mu\sigma}$ and $L^{\bar{\mu}\sigma}$ respectively (for the counterclockwise, aligned case).

It is important to keep track of the local orientation of the ribbon to ensure consistent action of ribbon operator [S2, S6]. The example ribbons displayed above are of counterclockwise local orientation. The circuit for a C_2 ribbon with clockwise local orientation is given by:



The circuit implementation for clockwise local orientation can be derived from that of counterclockwise local orientation by applying the inverse operation circuit (Eq. S5) on the dual edges.

3. Generalized ribbon operator

In Ref. S2, a generalized ribbon operator is introduced that relaxes the condition that a ribbon operator only create violations at two sites. It is observed the charge violations can be decoupled from the flux violations at the ends of a ribbon; i.e., the charge violations can be moved arbitrarily far away from the flux violations. This can be understood from the gauging perspective discussed in Section S1D2: the ungauging/regauging circuit can be along a length beyond the support of the FDLU part of the circuit. Then the charge violations are located at the starting and ending location of the ungauging/regauging circuit, whereas the flux violations are located at the starting and ending location of the FDLU circuit.

Generalized ribbon operators are crucial for defining the absolute logical encoding used in demonstrating the universal gate set in Section IV B, IV C, and IV D. A generalized ribbon operator ensures global fluxes are well-defined (so that they can be compared consistently) and allows the fluxes at the two ends of a ribbon operator to be at arbitrary locations and moved for braiding. When defining a flux, we need to specify the base point where the flux loop starts and ends. Changing base point conjugate the flux by the value of the path between the old and new base point. Therefore, to compare the flux internal states between distinct fluxes, the same base point has to be used to ensure consistency.

For the experiment, we fix the base point to be at the origin such that all global fluxes are defined with respect to the origin; an example is shown in Fig. 1(k).

S2. QUBIT-QUTRIT ENCODING OF $\mathcal{D}(S_3)$ STABILIZERS AND PROJECTORS

In Section S1B we detailed the encoding of S_3 group elements in terms of qubit-qutrit pairings and described the corresponding group multiplication rules. Here we provide a description of the *standard* $\mathcal{D}(S_3)$ projectors, originally defined in Ref. S3, expressed in terms of their action on the qubit and qutrit degrees of freedom. In particular, we derive explicit expressions for the $\mathcal{D}(S_3)$ vertex and plaquette projectors in terms of the operators $A^{\mathbb{Z}_2}$, $A^{\mathbb{Z}_3}$, $B^{\mathbb{Z}_2}$, and $B^{\mathbb{Z}_3}$ introduced in Eqs. (1) and (2). Similar definitions for the W -flux and non-contractible projectors are also given.

We begin by writing $\mathcal{D}(S_3)$ hamiltonian

$$H = - \sum_v A_v - \sum_p B_p^e, \quad (\text{S23})$$

where the vertex projector $A_v = \frac{1}{6} \sum_{g \in S_3} A_v^g$ enforces the zero-charge constraint at each vertex. Pictorially, the action of A_v^g on a group-element configuration is

$$A_v^g \left(\begin{array}{c} x_N \bullet \\ \uparrow \\ \text{---} \bullet v \text{---} \\ \downarrow \\ x_S \bullet \end{array} \begin{array}{c} x_E \bullet \\ \leftarrow \\ \text{---} \bullet v \text{---} \\ \rightarrow \\ x_W \bullet \end{array} \right) = \begin{array}{c} gx_N \bullet \\ \uparrow \\ \text{---} \bullet v \text{---} \\ \downarrow \\ x_S \bar{g} \bullet \end{array} \begin{array}{c} gx_E \bullet \\ \leftarrow \\ \text{---} \bullet v \text{---} \\ \rightarrow \\ x_W \bar{g} \bullet \end{array}. \quad (\text{S24})$$

Here, the solid circle and triangle on each edge denote the qubit–qutrit encoding. Similarly, the plaquette projectors B_p^e enforce trivial total flux around each plaquette. Their action on basis states is

$$B_p^e \left(\begin{array}{ccc} & x_N \bullet & \\ x_W \bullet & & x_E \bullet \\ & x_S \bullet & \end{array} \right) = \delta_{e, x_W x_N \bar{x}_E \bar{x}_S} \begin{array}{ccc} & x_N \bullet & \\ x_W \bullet & & x_E \bullet \\ & x_S \bullet & \end{array}, \quad (\text{S25})$$

where $x_W x_N \bar{x}_E \bar{x}_S$ denotes the flux around the plaquette.

A. Vertex projectors

From Eq. (S24), A_v^g left-multiplies the group elements on the north (N) and east (E) edges, and right-multiplies those on the south (S) and west (W) edges. The full vertex projector can be written as

$$A_v = \frac{1}{6} \sum_{g \in S_3} \begin{array}{c} L^g \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ R^g \bullet \end{array} \begin{array}{c} L^g \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ R^g \bullet \end{array}. \quad (\text{S26})$$

Using Eqs. (S2) and (S3), the terms A_v^σ and A_v^μ can be expressed as

$$A_v^\sigma = \begin{array}{c} L^\sigma \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ R^\sigma \bullet \end{array} \begin{array}{c} L^\sigma \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ R^\sigma \bullet \end{array} = \begin{array}{c} CX \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ X \bullet \end{array} \begin{array}{c} CX \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ X \bullet \end{array}, \quad A_v^\mu = \begin{array}{c} L^\mu \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ R^\mu \bullet \end{array} \begin{array}{c} L^\mu \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ R^\mu \bullet \end{array} = \begin{array}{c} \chi \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ \chi^{-Z} \bullet \end{array} \begin{array}{c} \chi \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ \chi^{-Z} \bullet \end{array}. \quad (\text{S27})$$

Expanding A_v , we obtain

$$\begin{aligned} A_v &= \frac{1}{6} (1 + A^\mu + A^{\bar{\mu}} + A^\sigma + A^{\mu\sigma} + A^{\bar{\mu}\sigma}) \\ &= \frac{1}{6} (1 + A^\sigma) (1 + A^\mu + A^{\bar{\mu}}) \\ &= \frac{1}{2} \left(1 + \begin{array}{c} CX \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ X \bullet \end{array} \begin{array}{c} CX \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ X \bullet \end{array} \right) \cdot \frac{1}{3} \left[1 + \left(\begin{array}{c} \chi \bullet \\ \uparrow \\ \text{---} \bullet \\ \downarrow \\ \chi^{-Z} \bullet \end{array} \begin{array}{c} \chi \bullet \\ \leftarrow \\ \text{---} \bullet \\ \rightarrow \\ \chi^{-Z} \bullet \end{array} + \text{h.c.} \right) \right] \\ &= A^{\mathbb{Z}_2} A^{\mathbb{Z}_3}. \end{aligned} \quad (\text{S28})$$

B. Plaquette projector

Since B_p^e enforces trivial flux around each plaquette, it can be written as

$$B_p^e = \sum_{g_W, g_N, g_E, g_S \in S_3} \delta_{e, g_S g_E \bar{g}_N \bar{g}_W} \left| \begin{array}{ccc} & g_N \bullet & \\ g_W \bullet & & g_E \bullet \\ & g_S \bullet & \end{array} \right\rangle \left\langle \begin{array}{ccc} & g_N \bullet & \\ g_W \bullet & & g_E \bullet \\ & g_S \bullet & \end{array} \right|. \quad (\text{S29})$$

Using the qubit-qutrit encoding $g = \mu^a \sigma^b = \sigma^b \mu^{(-1)^b a}$ and $\bar{g} = \sigma^b \mu^{-a} = \mu^{(-1)^b a} \sigma^b$, the flux $g_S g_E \bar{g}_N \bar{g}_W$ becomes

$$\begin{aligned} \text{flux} &= \mu^{a_S} \sigma^{b_S} \mu^{a_E} \sigma^{b_E} \sigma^{b_N} \mu^{-a_N} \sigma^{b_W} \mu^{-a_W} \\ &= \mu^{a_S} \mu^{(-1)^{b_S} a_E} \sigma^{b_S+b_E+b_N} \mu^{-a_N} \mu^{(-1)^{b_W} a_W} \sigma^{b_W} \\ &= \mu^{a_S} \mu^{(-1)^{b_S} a_E} \mu^{(-1)^{b_S+b_E+b_N} a_N} \mu^{(-1)^{b_S+b_E+b_N+b_W} a_W} \sigma^{b_S+b_E+b_N+b_W}. \end{aligned} \quad (\text{S30})$$

Each term in the product can be written in terms of Z and $\mathcal{Z}$ operators on the qubit and qutrit degrees of freedom, that is,

$$\sigma^{b_S+b_E+b_N+b_W} \rightarrow \begin{array}{c} \bullet \\ \uparrow Z_W \\ \bullet \\ \downarrow Z_N \\ \bullet \\ \uparrow Z_E \\ \bullet \\ \downarrow Z_S \\ \bullet \end{array} := S^{\mathbb{Z}_2}, \quad (\text{S31})$$

$$\begin{aligned} \mu^{(-1)^{b_S+b_E+b_N+b_W} a_W} &\rightarrow \begin{array}{c} \bullet \\ \uparrow Z_W^{-S^{\mathbb{Z}_2}} \\ \bullet \\ \downarrow \\ \bullet \end{array}, \quad \mu^{(-1)^{b_S+b_E+b_N} a_N} \rightarrow \begin{array}{c} \bullet \\ \uparrow Z_N^{-Z_S Z_E Z_N} \\ \bullet \\ \downarrow \\ \bullet \end{array} = \begin{array}{c} \bullet \\ \uparrow Z_N^{-Z_W S^{\mathbb{Z}_2}} \\ \bullet \\ \downarrow \\ \bullet \end{array}, \\ \mu^{(-1)^{b_S} a_E} &\rightarrow \begin{array}{c} \bullet \\ \uparrow Z_E^{Z_S} \\ \bullet \\ \downarrow \\ \bullet \end{array}, \quad \text{and} \quad \mu^{a_S} \rightarrow \begin{array}{c} \bullet \\ \uparrow Z_S \\ \bullet \\ \downarrow \\ \bullet \end{array}. \end{aligned} \quad (\text{S32})$$

Collecting all contributions, the flux around the plaquette can be measured by

$$S^{\mathbb{Z}_2} \cdot \begin{array}{c} \bullet \\ \uparrow Z_W^{-S^{\mathbb{Z}_2}} \\ \bullet \\ \downarrow \\ \bullet \end{array} \cdot \begin{array}{c} \bullet \\ \uparrow Z_N^{-Z_S Z_E Z_N} \\ \bullet \\ \downarrow \\ \bullet \end{array} \cdot \begin{array}{c} \bullet \\ \uparrow Z_E^{Z_S} \\ \bullet \\ \downarrow \\ \bullet \end{array} = S^{\mathbb{Z}_2} \cdot S^{\mathbb{Z}_3}. \quad (\text{S33})$$

Moreover, projectors onto the $+1$ eigenspaces of $S^{\mathbb{Z}_2}$ and $S^{\mathbb{Z}_3}$ enforce the trivial flux constraint, and Eq. (S29) can be written as (see Ref. S8):

$$\begin{aligned} B_p^e &= \frac{1}{2} \left(1 + \begin{array}{c} \bullet \\ \uparrow Z_W \\ \bullet \\ \downarrow Z_N \\ \bullet \\ \uparrow Z_E \\ \bullet \\ \downarrow Z_S \\ \bullet \end{array} \right) \cdot \frac{1}{3} \left[1 + \left(\begin{array}{c} \bullet \\ \uparrow Z_W^{-S^{\mathbb{Z}_2}} \\ \bullet \\ \downarrow \\ \bullet \end{array} \cdot \begin{array}{c} \bullet \\ \uparrow Z_N^{-Z_S Z_E Z_N} \\ \bullet \\ \downarrow \\ \bullet \end{array} \cdot \begin{array}{c} \bullet \\ \uparrow Z_E^{Z_S} \\ \bullet \\ \downarrow \\ \bullet \end{array} + \text{h.c.} \right) \right] \\ &= B^{\mathbb{Z}_2} B^{\mathbb{Z}_3}. \end{aligned} \quad (\text{S34})$$

C. W -flux and non-contractible Z projectors of $\mathcal{D}(S_3)$

The W -flux is the topological flux defined by the ordered group product along a closed loop, based at a fixed origin, that encircles an excitation. Fig. 3 enumerates the different paths used to define the W -flux stabilizer S_{W_p} , each of which encloses exactly one plaquette p . The corresponding projector W_p is defined as the projector onto the $+1$ eigenspace of S_{W_p} . Using steps analogous to those employed in deriving the plaquette operator B_p , the W -flux projector can also be written as

$$W_p = W_p^{\mathbb{Z}_2} W_p^{\mathbb{Z}_3}. \quad (\text{S35})$$

Here $W_p^{\mathbb{Z}_2} = B_p^{\mathbb{Z}_2}$ (as defined in Eq. (S31)) for the paths listed in Fig. 3, while $W_p^{\mathbb{Z}_3}$ has a form analogous to $S_B^{\mathbb{Z}_3}$ and acts non-trivially on all qudits along the loop.

Correlators of the W -flux are obtained by forming products of stabilizers, for example $S_{W_{p_1}} S_{W_{p_2}}$, and projecting onto their joint $+1$ eigenspace, yielding the projector $\Pi_{W_{p_1} W_{p_2}}$. This projector similarly factorizes into independent $\mathbb{Z}_2$ and $\mathbb{Z}_3$ components.

$$\Pi_{W_{p_1} W_{p_2}} = \Pi_{W_{p_1} W_{p_2}}^{\mathbb{Z}_2} \Pi_{W_{p_1} W_{p_2}}^{\mathbb{Z}_3} \quad (\text{S36})$$

Similarly, the ‘ Z -logical’ projectors of the $\mathcal{D}(S_3)$ code are defined as the ordered product of group elements along non-contractible loops. For the horizontal and vertical directions, these take the form $Z_{\text{hor}} = Z_{\text{hor}}^{\mathbb{Z}_2} Z_{\text{hor}}^{\mathbb{Z}_3}$ and $Z_{\text{vert}} = Z_{\text{vert}}^{\mathbb{Z}_2} Z_{\text{vert}}^{\mathbb{Z}_3}$, respectively (see Fig. 1e–f for the corresponding stabilizer definitions).

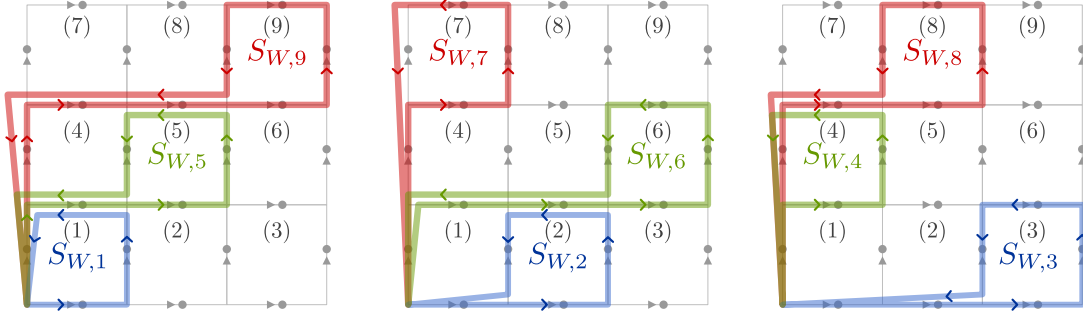


Figure 3. **W-flux.** Paths along which W -flux stabilizers are evaluated to benchmark the experimental results presented in the main text.

S3. COHERENT MOVING OF NON-ABELIAN ANYONS: THEORY

In our experiment, we only need to coherently move non-Abelian pure fluxes. We will focus on the protocol for coherently moving non-Abelian flux for the rest of the section. This is a special case of the general procedure of coherent moving [S9].

Suppose we want to coherently move an anyon from location A to B . Operationally, we demand a deterministic process such that there is no remnant anyon at the site A after coherent moving. We also require that, at no point in the protocol, is the global flux label (i.e. what defines the “absolute” logical encoding) stored locally; the coherent moving protocol only involves the local flux label instead. We implement coherent moving unitarily with the use of one ancilla qudit, where d matches the quantum dimension of the non-Abelian anyon.

Ribbon operators create pairs of excitation at their end points; to ensure that an anyon is deterministically moved from the starting site to the ending site, we need to ensure that the fusion outcome is always vacuum at the starting site. This is done by judicious choice of internal state at the ends of a ribbon operator. For a pair of non-Abelian fluxes, the state in the vacuum fusion channel is the symmetric superposition over all representatives in a conjugacy class C , i.e.,

$$|\text{vac}\rangle = \frac{1}{|C|} \sum_{c \in C} |c, \bar{c}\rangle. \quad (\text{S37})$$

The pairing of $c, \bar{c}$ ensures flux neutrality since the total flux is $c \cdot \bar{c} = e$; the symmetrization ensures charge neutrality as any permutation of the terms leave the state invariant, so the state is not charged under any symmetry action (which acts by conjugation). For moving non-Abelian *pure* flux, the procedure is simplified because pure fluxes already satisfy the property that the local internal flux state is a symmetric superposition over all possible basis states. Therefore, we only need to ensure the correct pairing $c, \bar{c}$ for the internal state of the ribbon operators.

The coherent moving protocol of non-Abelian flux consists of three main steps:

1. entangle the ancilla with the local internal state of the flux;
2. conditioned on the ancilla, apply the corresponding ribbon operator;
3. disentangle ancilla.

For step 1, a $C_L \mathcal{X}_a$ gate is implemented to entangle the ancilla with the local internal state of the flux.

For step 2, a C_2 ribbon operator is conjugated by $C_a \mathcal{X}_{dir}$, where dir is the first direct edge of the ribbon operator, such that starting local internal state created by the ribbon operator matches with the ancilla state. This ensures that the ribbon operator create the same C_2 flux at the starting site to cancel out the flux there.

For step 3, we want to implement a $C_L \mathcal{X}_a$ gate again to disentangling the ancilla from the local internal state of the flux at the start. However, the flux at the start is annihilated to vacuum! The key is to recognize that the flux label *based at the same vertex* remains the same due to the fact that ribbon operator can only create vacuum pair of anyons, i.e. ribbon operator satisfies the neutrality condition by construction [S2]. Therefore, the control qutrit is the global internal state (based at the starting vertex of the ribbon operator) of the new C_2 flux at the ending site.

See Section S10 A and Section S10 B for examples of coherent moving circuit implementations.

S4. PHYSICAL ENCODING AND IMPLEMENTATION OF PRIMITIVE GATES FOR $\mathcal{D}(S_3)$

This Section describes our implementation of qudit encoding and gates, realized with the native gate set of Quantinuum H2-series devices [S10, S11]. The native gates for the H2 architecture are as follows:

$$U_{1q}(\alpha, \beta) = e^{-\frac{1}{2}i\pi\beta Z} e^{-\frac{1}{2}i\pi\alpha X} e^{+\frac{1}{2}i\pi\beta Z}, \quad R_Z(\alpha) = e^{-\frac{1}{2}i\pi\theta Z} \quad \text{and} \quad (\text{S38})$$

$$\text{ZZPhase}(\theta) = e^{-\frac{1}{2}i\pi\theta(Z \otimes Z)}. \quad (\text{S39})$$

The number of ZZPhase gates significantly affects the overall implementation cost and serves as a key metric for evaluating gate decompositions.

Each six-dimensional qudit, representing the $\mathcal{D}(S_3)$ degrees of freedom on the edges of our square lattice, is encoded using a qutrit and a qubit. The qutrit, in turn, is encoded using two qubits. Therefore, each qudit requires a total of 3 qubits. The qutrit encoding is defined as:

$$|0\rangle_{\blacktriangle} := |00\rangle, \quad |1\rangle_{\blacktriangle} := |10\rangle, \quad |2\rangle_{\blacktriangle} := |11\rangle, \quad |\text{nc}\rangle := |01\rangle. \quad (\text{S40})$$

The non-computational state, $|\text{nc}\rangle$, is used for error detection; its presence indicates that the two encoding qubits have fallen outside the qutrit subspace.

Qutrits are primarily manipulated using the $\mathcal{X}$, $\mathcal{Z}$, and $\mathcal{C}$ gates, as defined in Eqs. (S4). $\mathcal{X}$ and $\mathcal{Z}$ satisfy the commutation relation $\mathcal{X}\mathcal{Z} = \omega\mathcal{Z}\mathcal{X}$. The charge-conjugation gate $\mathcal{C}$ is Hermitian, $\mathcal{C}^\dagger = \mathcal{C}$, and conjugates $\mathcal{X}$ and $\mathcal{Z}$:

$$\mathcal{C}\mathcal{X}\mathcal{C} = \mathcal{X}^\dagger \quad \text{and} \quad \mathcal{C}\mathcal{Z}\mathcal{C} = \mathcal{Z}^\dagger. \quad (\text{S41})$$

The qutrit $\mathcal{Z}$ gate can be efficiently implemented using two single-qubit native R_Z gates:

$$\mathcal{Z} := \begin{array}{c} \text{---} \boxed{R_Z(\frac{2}{3})} \text{---} \\ \text{---} \boxed{R_Z(\frac{2}{3})} \text{---} \end{array} = \bar{\omega} \text{diag}(1, \omega, \omega, \bar{\omega}). \quad (\text{S42})$$

The action of $\mathcal{Z}$ on $|\text{nc}\rangle$ is trivial (up to ω phase). The extra $\bar{\omega}$ phase needs to be corrected to properly implement a controlled- $\mathcal{Z}$ ($\mathcal{C}\mathcal{Z}$) gate. In contrast, implementing $\mathcal{X}$ requires either two CNOT gates or a single ZZPhase gate:

$$\mathcal{X} := \begin{array}{c} \text{---} \bullet \text{---} \oplus \text{---} \boxed{X} \text{---} \\ \text{---} \oplus \text{---} \bullet \text{---} \end{array} = \begin{array}{c} \text{---} \boxed{U_{1q}(1.5, 1.5)} \text{---} \\ \text{---} \end{array} \text{ZZPhase}(0.5) \begin{array}{c} \text{---} \boxed{U_{1q}(0.5, 0)} \text{---} \boxed{R_Z(3.5)} \text{---} \\ \text{---} \boxed{U_{1q}(1, 1.5)} \text{---} \boxed{R_Z(2.5)} \text{---} \end{array} = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix} \quad (\text{S43})$$

The charge-conjugation gate $\mathcal{C}$ can also be implemented using only a single ZZPhase gate.

$$\mathcal{C} := \begin{array}{c} \text{---} \bullet \text{---} \\ \text{---} \oplus \text{---} \end{array} = \begin{array}{c} \text{---} \text{---} \boxed{R_Z(0.5)} \text{---} \\ \text{---} \boxed{U_{1q}(2.5, 0.5)} \text{---} \text{ZZPhase}(0.5) \text{---} \boxed{U_{1q}(0.5, 0)} \text{---} \boxed{R_Z(1.5)} \text{---} \end{array} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \quad (\text{S44})$$

$\mathcal{X}$ can be transformed into $\mathcal{Z}$ gate by performing a qutrit Fourier transform, denoted by $\mathcal{H}$ which acts on the qutrit space as follows:

$$\mathcal{H}|i\rangle_{\blacktriangle} = \frac{1}{\sqrt{3}}(|0\rangle_{\blacktriangle} + \omega^i|1\rangle_{\blacktriangle} + \omega^{2i}|2\rangle_{\blacktriangle}), \quad (\text{S45})$$

The native implementation of $\mathcal{H}$ requires three ZZPhase gates, preserving the $|\text{nc}\rangle$ state. Note that we treat the action of $\mathcal{H}$ on $|0\rangle_{\blacktriangle}$ as a *state preparation procedure*, which can be implemented using one ZZPhase gate:

$$\mathcal{H}|00\rangle = \begin{array}{c} |0\rangle \text{---} \boxed{U_{1q}(0.5, 0)} \text{---} \\ |0\rangle \text{---} \boxed{U_{1q}(\cos^{-1}(-1/3)/\pi, 1.5)} \text{---} \end{array} \text{ZZPhase}(0.25) \begin{array}{c} \text{---} \boxed{U_{1q}(1.5, 1.25)} \text{---} \boxed{R_Z(0.75)} \text{---} \\ \text{---} \boxed{U_{1q}(1, 1)} \text{---} \boxed{R_Z(3)} \text{---} \end{array} \Bigg\} \frac{1}{\sqrt{3}}(|00\rangle + |10\rangle + |11\rangle). \quad (\text{S46})$$

Furthermore, $\mathcal{H}$ commutes with $\mathcal{C}$, since $\mathcal{H}^2 = \mathcal{C}$.

Similarly, H_{12} diagonalizes $\mathcal{C}$, leaving $|0\rangle_{\blacktriangle}$ unchanged while applying a ‘‘Hadamard gate’’ to both $|1\rangle_{\blacktriangle}$ and $|2\rangle_{\blacktriangle}$. In terms of native gates, H_{12} requires a single ZZPhase gate:

$$H_{12} := \begin{array}{c} \text{---} \\ \text{---} \end{array} \begin{array}{c} \boxed{U_{1q}(3.25, 1.5)} \\ \boxed{ZZPhase(0.5)} \end{array} \begin{array}{c} \boxed{R_Z(0.5)} \\ \boxed{U_{1q}(2.75, 0)} \end{array} \begin{array}{c} \text{---} \\ \text{---} \end{array} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ 0 & 0 & \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}. \quad (\text{S47})$$

Since $\mathcal{C}$ squares to the identity, the CC gate is defined with a qubit control and qutrit target. To avoid applying three qubit CCZ gate to implement CC , we use the fact that H_C disentangles the $\mathcal{C}$ gate, i.e., $H_{12} \mathcal{C} H_{12}^\dagger = IZ$. The implementation of CC requires three ZZPhase gates:

$$CC := \begin{array}{c} \text{control qubit} \text{---} \\ \text{target qutrit} \left\{ \begin{array}{c} \boxed{H_{12}} \\ \boxed{Z} \\ \boxed{H_{12}^\dagger} \end{array} \right. \end{array} \quad (\text{S48})$$

The application of CZ requires four CR_z gates:

$$CZ := \begin{array}{c} \text{control qutrit} \left\{ \begin{array}{c} \boxed{R_Z(\frac{2}{3})} \\ \boxed{R_Z(\frac{2}{3})} \end{array} \right. \\ \text{target qutrit} \left\{ \begin{array}{c} \boxed{R_Z(\frac{2}{3})} \\ \boxed{R_Z(\frac{2}{3})} \\ \boxed{R_Z(\frac{2}{3})} \\ \boxed{R_Z(\frac{2}{3})} \end{array} \right. \end{array}, \quad (\text{S49})$$

where each of the CR_z gate can be implemented using a single ZZPhase gate. The initial single qubit R_z gates on control qutrit compensate for the extra phase $\bar{\omega}$ in Eq. (S43). Applying CZ while transforming the target qutrit by $\mathcal{H}$ yields the $C\mathcal{X}$ gate:

$$C\mathcal{X} := \begin{array}{c} \text{---} \\ \text{---} \end{array} \boxed{CZ} \begin{array}{c} \boxed{\mathcal{H}} \\ \boxed{\mathcal{H}^\dagger} \end{array} \quad (\text{S50})$$

The $C\mathcal{X}$ gate requires a total of $3 + 4 + 3 = 10$ ZZPhase gates.

We define a $C_2\text{NOT}$ gate with a qutrit as the control and a qubit as the target.

$$C_2\text{NOT} := \begin{array}{c} \text{control qutrit} \left\{ \text{---} \right. \\ \text{target qubit} \text{---} \oplus \end{array} \quad (\text{S51})$$

It acts trivially on the target qubit when the control qutrit is in the $|0\rangle$ or $|1\rangle$ state, and flips the target qubit when the control is in the $|2\rangle$ state. It is implemented by a single CNOT gate conditioned on the second physical qubit of the control qutrit, because in the physical encoding of the qutrit (see Eq. (S40)), the second physical qubit is in the $|1\rangle$ state only for the $|2\rangle = |11\rangle$ qutrit state.

Moreover, we define the CX_{12} gate with a qutrit control and a qutrit target.

$$CX_{12} := \begin{array}{c} \text{control qutrit} \left\{ \text{---} \right. \\ \text{target qutrit} \left\{ \begin{array}{c} \boxed{H_{12}} \\ \boxed{Z} \\ \boxed{H_{12}^\dagger} \end{array} \right. \end{array} \quad (\text{S52})$$

The gate acts as the identity when the control qutrit is in the $|0\rangle$ or $|1\rangle$ state. When the control is in the $|2\rangle$ state, it swaps the target states $|1\rangle$ and $|2\rangle$ while leaving $|0\rangle$ invariant.

S5. CIRCUIT OPTIMIZATION BY LOCAL BASIS TRANSFORMATION

For each experiment, before compiling the circuit unitary U into native gates using TKET [S12], we locally transform every qutrit to the dual ($\mathcal{X}$) basis by applying $\mathcal{H}$, as defined in Eq. (S45), and construct $\tilde{U}$ such that

$$\tilde{U} = \mathcal{H}^{\otimes n \dagger} U \mathcal{H}^{\otimes n}, \quad (\text{S53})$$

where n is the number of qutrits. We then compile and execute $\mathcal{H}^{\otimes n} \tilde{U} \mathcal{H}^{\otimes n}$ instead of compiling and executing U directly. This local basis transformation reduces the number of gates after compilation and, empirically, also reduces memory errors. For comparison,

$$\begin{aligned} \mathcal{X}\text{-basis measurement: } & \begin{matrix} n_{2q} \text{ gates} = 792 \\ \text{depth}_{2q} = 261 \end{matrix} \xrightarrow[\text{transformation}]{\text{local basis}} \begin{matrix} 744 \\ 232 \end{matrix}, \text{ and} \\ \text{Pull-through gate: } & \begin{matrix} n_{2q} \text{ gates} = 911 \\ \text{depth}_{2q} = 354 \end{matrix} \xrightarrow[\text{transformation}]{\text{local basis}} \begin{matrix} 845 \\ 307 \end{matrix}. \end{aligned} \quad (\text{S54})$$

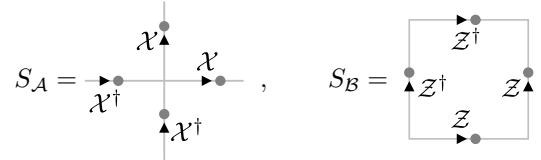
It is straightforward to compute $\tilde{U}$ because $\mathcal{H}$ commutes with the charge-conjugation gate $\mathcal{C}$ and maps a $\text{C}\mathcal{X}$ gate to another $\text{C}\mathcal{X}$ gate with the control and target qutrits swapped.

S6. GROUND STATE PREPARATION

This section details the methodology for preparing the ground state of the $\mathcal{D}(S_3)$ model. The preparation strategy we use follows the method given in the appendix of Ref. S13 and consists of a multi-step procedure. The first step is the preparation of the $\mathbb{Z}_3$ toric code ground state. Subsequently, this Abelian state is transformed into the desired non-Abelian $\mathcal{D}(S_3)$ order by gauging its charge-conjugation symmetry. A detailed exposition of these steps follows.

A. Preparation of $\mathbb{Z}_3$ toric code ground state

The ground state of the $\mathbb{Z}_3$ toric code with qutrits on the edges of the square lattice is stabilized by,



$$S_A = \begin{array}{c} \bullet \\ \uparrow \mathcal{X} \\ \bullet \\ \downarrow \mathcal{X}^\dagger \\ \bullet \end{array} \quad , \quad S_B = \begin{array}{ccc} & \bullet & \\ \uparrow \mathcal{Z}^\dagger & & \uparrow \mathcal{Z}^\dagger \\ \bullet & \square & \bullet \\ \downarrow \mathcal{Z} & & \downarrow \mathcal{Z} \\ \bullet & & \bullet \end{array} . \quad (\text{S55})$$

While these stabilizers can be prepared in constant depth using an adaptive circuit with mid-circuit measurements and ancilla qutrits, this approach incurs a significant qubit overhead. On the Quantinuum System Model H2, which has 56 qubits, preparing the $\mathcal{D}(S_3)$ ground state on a 3×3 lattice already requires 54 qubits, leaving little room for additional ancillae. We therefore employ a purely unitary preparation protocol that requires neither mid-circuit measurements nor ancillae for the $\mathbb{Z}_3$ toric code state preparation itself. The protocol is described below.

1. All qutrits are initialized in the state $|0\rangle$, which trivially satisfies all plaquette stabilizers $S_B = +1$.
2. To satisfy the vertex stabilizers $S_A = +1$, we proceed in an iterative manner:
 - (i) Select a vertex stabilizer and designate one of its qutrits as the “representative.” The representative qutrit is prepared in the state $|+\rangle_\bullet := \frac{1}{\sqrt{3}}(|0\rangle_\bullet + |1\rangle_\bullet + |2\rangle_\bullet)$, using the state-preparation circuit of Eq. (S46).
 - (ii) Apply a sequence of $\text{C}\mathcal{X}$ or $\text{C}\mathcal{X}^\dagger$ gates from the representative qutrit to the remaining qutrits within the same vertex stabilizer. A $\text{C}\mathcal{X}$ gate is applied when both the control and target qutrits are acted on by $\mathcal{X}\mathcal{X}$ or $\mathcal{X}^\dagger\mathcal{X}^\dagger$ in the vertex stabilizer; otherwise, a $\text{C}\mathcal{X}^\dagger$ gate is applied.

This procedure is repeated for all vertex stabilizers except one, ensuring that a qutrit is selected as a representative only if it is in the $|0\rangle$ state and has not been acted upon previously. The final vertex stabilizer is implicitly satisfied due to the global constraint $\prod_v S_{A_v} = 1$. The vertex preparation order used to obtain the $\mathbb{Z}_3$ toric code ground state on 3×2 and 3×3 lattices is illustrated in Fig. 4a.

The resulting state is the logical ground state $|00\rangle_L$ of the $\mathbb{Z}_3$ toric code. To prepare the logical state $|++\rangle_L$, where both logical $\mathcal{X}_L$ operators have eigenvalue $+1$, a dual procedure is employed. All qutrits are first initialized in the $|+\rangle_{\blacktriangle}$ state, satisfying all vertex stabilizers $S_{\mathcal{A}} = +1$. Then, plaquette stabilizers $S_{\mathcal{B}} = +1$ are enforced using a protocol analogous to steps (i) and (ii) above.

B. Gauging the $\mathbb{Z}_2$ charge-conjugation symmetry

The $\mathbb{Z}_3$ toric code Hamiltonian possesses a global $\mathbb{Z}_2$ charge-conjugation symmetry, denoted $\prod \mathcal{C}$ (see Eqs. (S41) and (S55)). Gauging this global $\mathbb{Z}_2$ symmetry means promoting it to a local gauge symmetry, where an independent $\mathbb{Z}_2$ transformation can be performed locally at each “site”. This procedure *enriches* the topological order, transforming the $\mathbb{Z}_3$ toric code into the $\mathcal{D}(S_3)$ quantum double model [S13, S14].

The gauging process can be implemented via an adaptive quantum circuit involving measurements and feed-forward, or via a purely unitary circuit composed of a deterministic sequence of gates. For the experiments presented in the main text, we prepare the $\mathcal{D}(S_3)$ ground state using a unitary protocol. The data for the $\mathcal{D}(S_3)$ ground state prepared using an adaptive circuit is shown in Fig. 2.

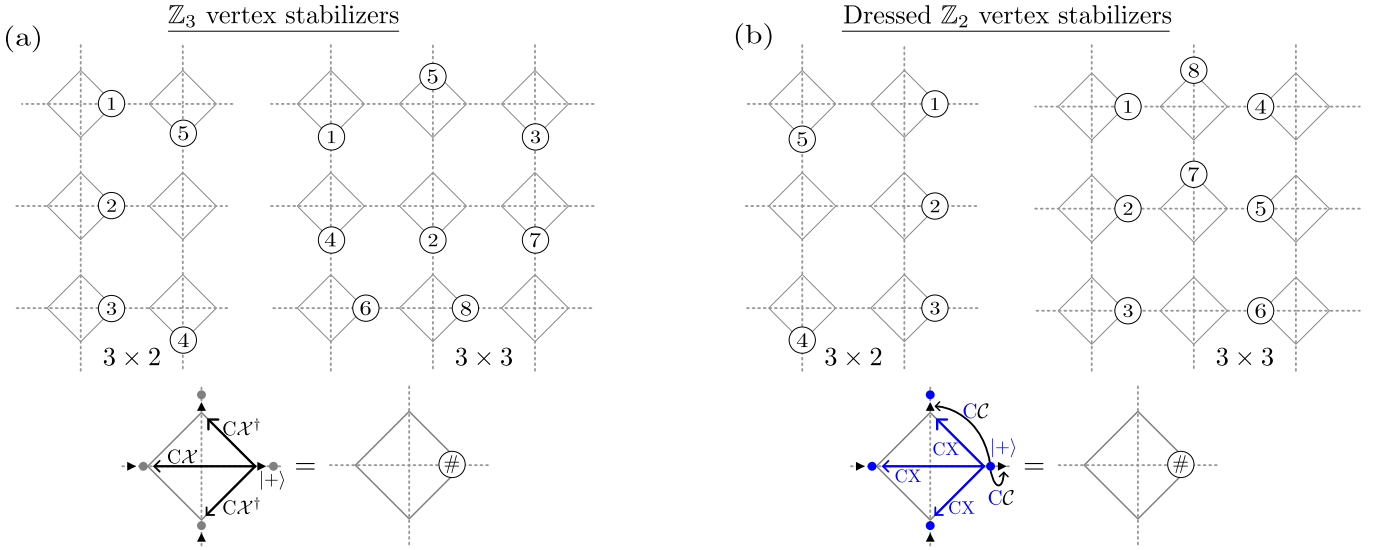
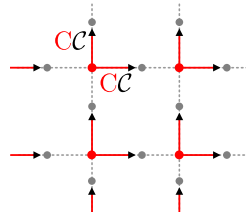


Figure 4. **Vertex Stabilizers Unitary Preparation.** Ordering for preparing vertex stabilizers in (a) the $\mathbb{Z}_3$ toric code and (b) the dressed $\mathbb{Z}_2$ toric code, for the 3×2 and 3×3 lattices. The dashed lines outline the square lattice. Rotated squares at each vertex represent the $S_{\mathcal{A}}$ stabilizers. Numbered edges denote the control qutrits, and the sequence of numbers indicates the order in which these vertex stabilizers are prepared.

1. Using adaptive circuit

The gauging protocol augments the Hilbert space and enforces new local constraints:

1. Introduce an auxiliary qubit at each vertex of the lattice initialized in the $|+\rangle$ state, stabilized by X .
2. Apply a layer of CC gates that entangle the original $\mathbb{Z}_3$ qutrits with the new $\mathbb{Z}_2$ vertex qubits. Specifically, for each vertex, we apply two CC gates, where $\mathbb{Z}_2$ vertex qubit acts as the control and the right and top neighboring $\mathbb{Z}_3$ edge qutrits act as targets.



(S56)

Stabilizers of the joint state are now given by

$$\begin{array}{c} \text{X}^{Z_1} \\ \text{X}^{-Z_3} \text{X}^{Z_1} \\ \text{X}^{-Z_2} \end{array}, \quad \begin{array}{c} \text{Z}^{-Z_3} \\ \text{Z}^{-Z_1} \text{Z}^{Z_2} \\ \text{Z}^{Z_1} \end{array}, \quad \begin{array}{c} \text{C} \\ \text{X} \\ \text{C} \end{array}. \quad (\text{S57})$$

3. We further add a qubit on each edge initialized in the $|0\rangle$ state, stabilized by Z . We apply CX gates between vertex and neighboring qubits with vertex qubits as the control and edge qubits as the target.

$$(\text{S58})$$

The expanded set of stabilizers for the state augmented with edge qubits is given by

$$\begin{array}{c} \text{X}^{Z_1} \\ \text{X}^{-Z_3} \text{X}^{Z_1} \\ \text{X}^{-Z_2} \end{array}, \quad \begin{array}{c} \text{Z}^{-Z_3} \\ \text{Z}^{-Z_1} \text{Z}^{Z_2} \\ \text{Z}^{Z_1} \end{array}, \quad \begin{array}{c} \text{X} \\ \text{C} \\ \text{X} \end{array}, \quad \begin{array}{c} \text{Z} \\ \text{Z} \\ \text{Z} \end{array}, \quad \begin{array}{c} \text{Z} \\ \text{Z} \\ \text{Z} \end{array}. \quad (\text{S59})$$

These stabilizers of the state do not all commute with each other. However, by writing these stabilizers as projectors over the trivial (+1) eigenspace, we obtain a set of commuting projectors with a *simpler* form (see [S13] for details)

$$\begin{aligned} & \frac{1}{3} \left[1 + \left(\begin{array}{c} \text{X} \\ \text{X}^{-Z_3} \text{X}^{Z_1} \\ \text{X}^{-Z_2} \end{array} + \text{h.c.} \right) \right], \quad \frac{1}{3} \left[1 + \left(\begin{array}{c} \text{Z}^{-Z_3} \text{Z}^{Z_1} \\ \text{Z}^{\dagger} \text{Z}^{Z_2} \text{Z}^{Z_1} \\ \text{Z} \end{array} + \text{h.c.} \right) \right] \\ & \frac{1}{2} \left[1 + \begin{array}{c} \text{X} \\ \text{C} \\ \text{X} \end{array} \right], \quad \frac{1}{2} \left[1 + \begin{array}{c} \text{Z} \\ \text{Z} \\ \text{Z} \end{array} \right], \quad \frac{1}{2} \left[1 + \begin{array}{c} \text{Z} \\ \text{Z} \\ \text{Z} \end{array} \right] \end{aligned} \quad (\text{S60})$$

Moreover, the vertex qubit dependence of $\mathbb{Z}_3$ projectors can be eliminated by writing vertex qubit variables in

terms of edge qubits as determined by CX gates between vertex and edge qubits.

$$\begin{aligned}
 & \frac{1}{3} \left[1 + \left(\begin{array}{c} \mathcal{X} \\ \text{---} \bullet_3 \mathcal{X}^{-Z_3} \text{---} \mathcal{X} \text{---} + \text{h.c} \\ \mathcal{X}^{-Z_2} \\ \bullet_2 \end{array} \right) \right], \quad \frac{1}{3} \left[1 + \left(\begin{array}{c} \mathcal{Z}^{-Z_2} \\ \mathcal{Z}^\dagger \quad \mathcal{Z}^{Z_1} \\ \bullet_2 \quad \bullet_1 \mathcal{Z} \end{array} + \text{h.c} \right) \right] \\
 & \frac{1}{2} \left[1 + \begin{array}{c} \mathcal{X} \\ \mathcal{C} \\ \text{---} \mathcal{X} \text{---} \mathcal{C} \mathcal{X} \text{---} \\ \mathcal{X} \end{array} \right], \quad \frac{1}{2} \left[1 + \begin{array}{c} \mathcal{Z} \quad \mathcal{Z} \quad \mathcal{Z} \end{array} \right], \quad \frac{1}{2} \left[1 + \begin{array}{c} \mathcal{Z} \\ \mathcal{Z} \\ \mathcal{Z} \end{array} \right]
 \end{aligned} \tag{S61}$$

4. Finally, we measure vertex qubits in X -basis and enforce local $\mathbb{Z}_2$ gauge invariance. This projects the state into the subspace where the original $\mathbb{Z}_3$ vertex constraints are now tied to the state of the $\mathbb{Z}_2$ gauge qubits. These measurements yield ± 1 outcomes. Based on the outcomes, we apply corrective Z 's to the $\mathbb{Z}_2$ edge qubits. This feed-forward ensures that the system is deterministically prepared in the $+1$ eigenspace of all the projectors, namely

$$\begin{aligned}
 A^{\mathbb{Z}_3} &= \frac{1}{3} \left[1 + \left(\begin{array}{c} \mathcal{X} \\ \text{---} \bullet_3 \mathcal{X}^{-Z_3} \text{---} \mathcal{X} \text{---} + \text{h.c} \\ \mathcal{X}^{-Z_2} \\ \bullet_2 \end{array} \right) \right], \quad B^{\mathbb{Z}_3} = \frac{1}{3} \left[1 + \left(\begin{array}{c} \mathcal{Z}^{-Z_2} \\ \mathcal{Z}^\dagger \quad \mathcal{Z}^{Z_1} \\ \bullet_2 \quad \bullet_1 \mathcal{Z} \end{array} + \text{h.c} \right) \right] \\
 A^{\mathbb{Z}_2} &= \frac{1}{2} \left[1 + \begin{array}{c} \mathcal{X} \\ \mathcal{C} \\ \text{---} \mathcal{X} \text{---} \mathcal{C} \mathcal{X} \text{---} \\ \mathcal{X} \end{array} \right], \quad B^{\mathbb{Z}_2} = \frac{1}{2} \left[1 + \begin{array}{c} \mathcal{Z} \\ \mathcal{Z} \quad \mathcal{Z} \\ \mathcal{Z} \end{array} \right]
 \end{aligned} \tag{S62}$$

The resulting projectors can be shown to be equivalent to the canonical quantum double projectors (see Section S2) within the ground state subspace of the $\mathcal{D}(S_3)$ model. In particular, the vertex projectors and the $B^{\mathbb{Z}_2}$ are exactly the same, whereas the $B^{\mathbb{Z}_3}$ given in Eq. (S34) in Section S2 can be reduced to the above expression when $B^{\mathbb{Z}_2} = +1$. While this method is optimal in terms of circuit depth, its reliance on mid-circuit measurements, feed-forward, and ancilla qubits affects the fidelity of the prepared state.

2. Using unitary circuit

We developed a purely unitary protocol that requires no measurements or feed-forward. The charge-conjugation gauging procedure, as described above, effectively amounts to measuring *dressed* vertex stabilizers of the $\mathbb{Z}_2$ toric code using an ancilla, after preparing the $\mathbb{Z}_3$ toric code, which is given as

$$\begin{array}{c} \mathcal{X} \\ \mathcal{C} \\ \text{---} \mathcal{X} \text{---} \mathcal{C} \mathcal{X} \text{---} \\ \mathcal{X} \end{array} \tag{S63}$$

We can omit the need for introducing the ancilla (red vertex) qubit in the stabilizer measurement by using the following circuit identity.

The diagram shows two equivalent quantum circuits. The left circuit has a red line (ancilla) and four blue lines (qubits). The red line starts with a $|0\rangle$ state, followed by a Hadamard gate H , and then four CNOT gates with targets on the blue lines. After the CNOTs, the red line has a Hadamard gate H and is measured, with the result labeled $0/1$. The right circuit is identical but the red line is replaced by a blue line initialized to $|0\rangle$, and the measurement is omitted. Both circuits have control lines labeled c at the bottom. The equation is labeled (S64) on the right.

where blue and red lines correspond to qubits, and black lines correspond to qutrits. This identity shows that a circuit involving an ancilla measurement (left) can be replaced by a purely unitary sequence of gates (right), provided that one of the target qubits is initialized in the $|0\rangle$ state. This constraint forces the preparation of the dressed $\mathbb{Z}_2$ vertex stabilizers to be performed sequentially, similar to the unitary preparation of the $\mathbb{Z}_3$ toric code ground state in Section S6 A. We select a representative qubit, prepare it in the qubit $|+\rangle$ state, and use it as a control to entangle its neighbors. This process is repeated for all vertices in a specific order, shown in Fig. 4b, ensuring that the final state is the ground state of the $\mathcal{D}(S_3)$ model.

C. Adaptive vs unitary preparation protocol

1. Fidelity estimate

The overall fidelity in the Quantinuum's H2 hardware is upper-bounded by the fidelity of the native two-qubit gate (around 0.998), which is the dominant source of error. Therefore, a quick back-of-the-envelope estimation of the ground state preparation fidelity is to use the following formula

$$(\text{Fidelity per edge}) = (\text{Fidelity of two-qubit gate})^{\frac{d}{N}} \quad (\text{S65})$$

where d is the total number of native two-qubit gates used in the preparation circuit, and $N = 3 \times 3 \times 2 = 18$ is the number of edges of the square lattice (for our 3×3 geometry).

For the unitary preparation protocol, the circuit depth is reported in the Supplementary Table S3, for the ground state unitary preparation, the native two-qubit gate depth $d = 238$. Therefore, the fidelity per edge is $(0.998)^{238/18} \approx 0.974$, which is within the fidelity bound of 0.970(5)–0.988(3) given in Eq. 3 in the main text.

For the measurement-based adaptive preparation protocol, the native two-qubit gate depth $d = 274$. Therefore, the fidelity per edge is $(0.998)^{274/18} \approx 0.970$, which is within the fidelity bound of 0.930(8)–0.978(2) given right below Eq. 3 in the main text. The fidelity is slightly worse than unitary preparation protocol because of the constant overhead involved for measurement-based protocol. Indeed, the measurement-based protocol requires slightly more gates since we need to entangle to ancillas (which we subsequently measure).

2. Scalability

For the measurement-based adaptive preparation protocol, the native two-qubit gate depth $d = 274$. Therefore, the fidelity per edge is $(0.998)^{274/18} \approx 0.970$, which is within the fidelity bound of 0.930(8)–0.978(2) given right below Eq. 3 in the main text. The fidelity is slightly worse than unitary preparation protocol because of the constant overhead involved for measurement-based protocol. Indeed, the measurement-based protocol requires slightly more gates since we need to entangle to ancillas (which we subsequently measure).

For the current system sizes, error contributions are dominated by gate count rather than runtime; therefore, the unitary preparation has (slightly) better fidelity. Naturally, as system size is scaled up, the runtime of the unitary preparation scales linearly with system size and will eventually form a bottleneck, such that the adaptive protocol then wins out. More precisely, for the adaptive protocol, the number of gates grows proportional to system size ($\sim L^2$), while the preparation time remains constant, such that the fidelity-per-qubit will remain constant (which can then be combined with quantum error correction if this local fidelity is beyond some stability threshold). In contrast, for the unitary preparation, while the number of gates also grows as $\sim L^2$, the preparation time scales as $\sim L$, leading to a fidelity $\sim e^{-aL^3}$, which means $e^{-aL} \sim e^{-a\sqrt{N}}$ per qubit.

S7. C_2 AND C_3 FLUX INSERTION AND SINGLE NON-ABELIAN ANYON

This section details the quantum circuits required to prepare the ground states of the $\mathcal{D}(S_3)$ model threaded with non-trivial topological flux. Specifically, we focus on the C_2 and C_3 flux sectors. Furthermore, we outline the steps for preparing a single C_3 non-Abelian anyon, a capability that is fundamentally impossible in Abelian topological phases.

To access the non-trivial C_2 and C_3 sectors, we employ the symmetry un-gauging and re-gauging strategy described in Ref. S7. Construction circuits are provided for a 3×3 lattice on a torus.

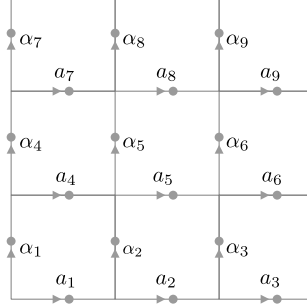


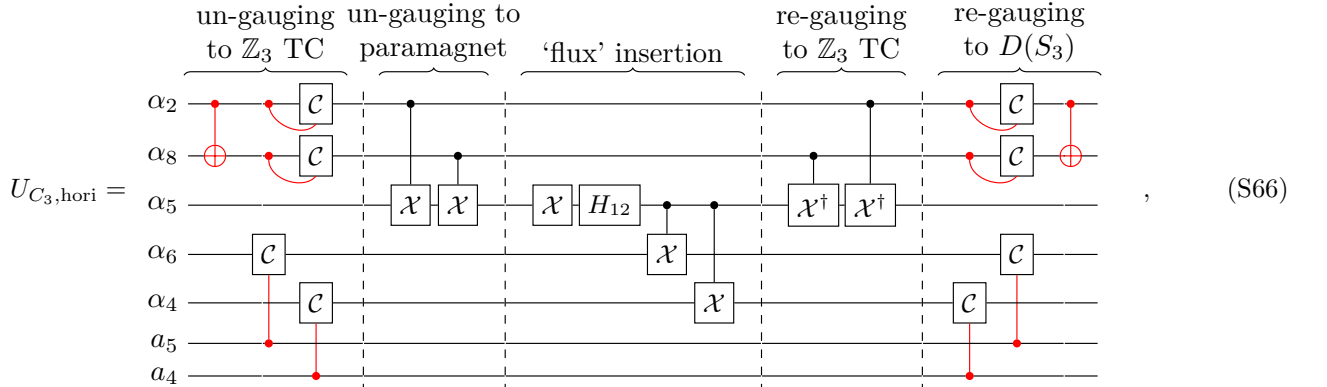
Figure 5. 3×3 square lattice with labeled edges

A. C_3 flux

The preparation of a ground state with C_3 flux, for instance $|\text{gs}_{C_3, \text{hori}}\rangle$, where the flux is threaded through bonds in the horizontal direction, involves:

- Ungauging the $\mathbb{Z}_2$ charge-conjugation symmetry.
- Ungauging the $\mathbb{Z}_3$ symmetry to access the ‘paramagnet’ qutrits.
- Applying a sequence of gates along a non-trivial horizontal loop to prepare the qutrits in the state $\frac{1}{\sqrt{2}}(|11\dots\rangle_{\blacktriangle} + |22\dots\rangle_{\blacktriangle})$.
- Re-gauging the $\mathbb{Z}_2$ and $\mathbb{Z}_3$ symmetries to return to the $\mathcal{D}(S_3)$ model with the desired C_3 flux.

The following is the optimized circuit for preparing $|\text{gs}_{C_3, \text{hori}}\rangle$ on a 3×3 lattice (see Fig. 5 for qudit labels) by applying it to the trivial ground state, $|\text{gs}\rangle$.



In the circuits above, each black line represents a qubit–qutrit pair encoding an S_3 qudit. Gate actions on the qubit lines are shown in red, and the gates act as follows:

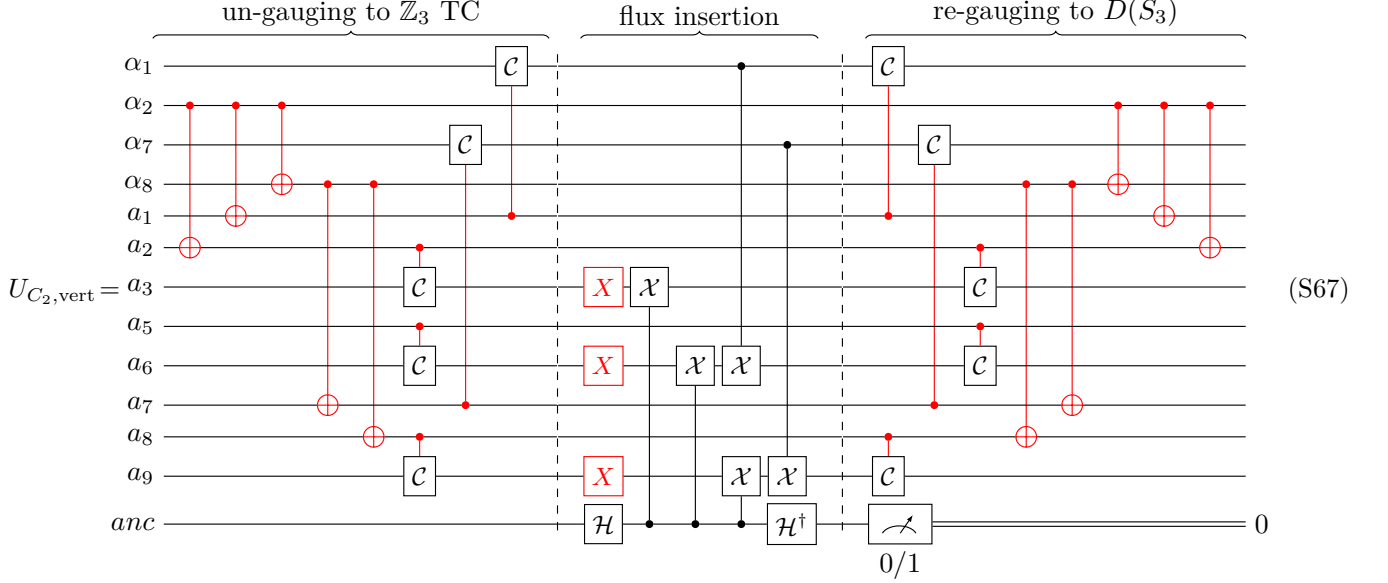
- CNOT acts on the qubit lines of the control and target qudits.
- $C\mathcal{X}$ acts on the qutrit lines of the control and target qudits.

- CC acts on the qubit line of the control and qutrit line of the target qutrit. Hence, the "self-loop" CC gates at the un-gauging and re-gauging step act on distinct qubit and qutrit lines of a qudit.
- $\mathcal{X}$ and H_{12} act on the qutrit degrees of freedom of the qudits.

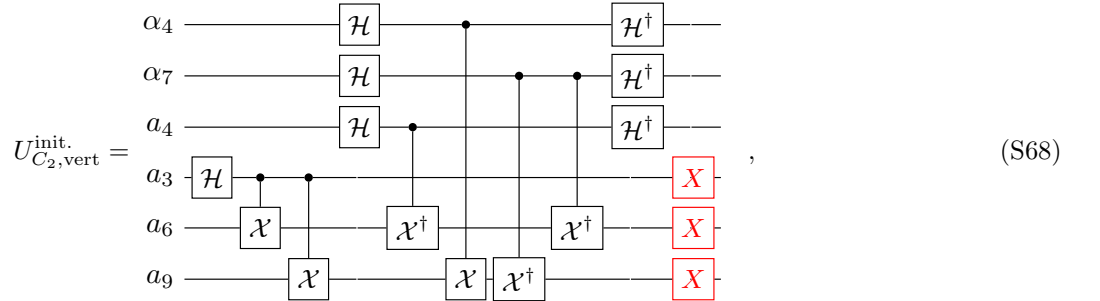
For detailed definitions and physical implementations, see Section S4.

B. C_2 flux

The construction of the ground state with C_2 flux requires ungauging to the $\mathbb{Z}_3$ toric code and then applying the projector $\frac{1}{3}(II \dots + \mathcal{X}\mathcal{X} \dots + \mathcal{X}^\dagger \mathcal{X}^\dagger \dots)$ across a non-trivial loop on qutrits, as well as X gates on the qubits. This requires an ancilla qutrit and complex feed-forward operations for deterministic preparation, as shown below, which creates the ground state $|\text{gs}_{C_2, \text{vert}}\rangle$ by acting on horizontal bonds across a non-trivial loop in the vertical direction.



To simplify and avoid mid-circuits measurement, we constructed an optimized initialization circuit. The key insight is that the flux insertion operation in the above circuit can be commuted through the $\mathbb{Z}_3$ gauging unitary used in the standard $\mathcal{D}(S_3)$ state preparation. This shifts the flux-insertion step to an earlier stage, where it acts on a simple product state, thereby eliminating the need for an ancilla. The resulting initialization circuit,



is applied to the qudits *before* the main $\mathcal{D}(S_3)$ preparation unitary (see Section S4). This circuit effectively implements the +1 post-selected projector by using one of the qutrits as a control (i.e. a_4 in Eq. (S68)) for a series of $C\mathcal{X}$ gates. The final $\mathcal{D}(S_3)$ ground state prepared after this initialization has C_2 flux. Note that the correctness of this optimized initialization circuit depends on the ordering of plaquettes used during the $\mathbb{Z}_3$ toric code preparation (see Fig. 4).

C. Single C_3 flux anyon on the torus

Distinct, non-commuting flux sectors enable the creation of a single non-Abelian anyon through a mechanism unique to non-Abelian topological order. When a non-Abelian flux pair is created in a ground state that already carries a

different non-trivial flux, braiding between the two flux types changes the fusion channel of the pair. Consequently, fusing the pair does not return the system to the vacuum but instead leaves behind a single isolated anyon.

In our implementation, we begin from the ground state $|\text{gs}_{C_2, \text{vert}}\rangle$ (see Fig. 6a) and create a pair of C_3 fluxes that encircle the torus non-trivially along the horizontal direction (see Fig. 6b). Because the C_2 and C_3 fluxes braid non-trivially, the C_3 pair is forced out of the vacuum fusion channel. Upon fusion, a remnant C_3 non-Abelian anyon with quantum dimension 2 remains localized at the fusion site, consistent with the fusion rule

$$C_3 \times C_3 = [+] + [-] + C_3. \quad (\text{S69})$$

This process can also be understood at the level of the $\mathbb{Z}_3$ toric code by ungauging the $\mathbb{Z}_2$ charge-conjugation symmetry. In this picture, the C_3 flux pair becomes the superposition of m and $\bar{m}$ anyons in the state $|m\rangle_1 |\bar{m}\rangle_2 + |\bar{m}\rangle_1 |m\rangle_2$, while the C_2 flux becomes a charge-conjugation defect [S15]. The C_2 non-contractible loop becomes a charge-conjugation domain wall, which toggles $m \leftrightarrow \bar{m}$. So moving the second particle of the pair across the charge-conjugation domain wall maps the state $|m\rangle_1 |\bar{m}\rangle_2 + |\bar{m}\rangle_1 |m\rangle_2 \rightarrow |m\rangle_1 |m\rangle_2 + |\bar{m}\rangle_1 |\bar{m}\rangle_2$. Finally, fusing particle 1 and 2 together gives $|m\rangle + |\bar{m}\rangle$; this is why the remnant particle is a C_3 flux. We note that the location of the single anyon excitation is where the fusion happened, not where the C_3 flux crosses the support of the $\hat{X}_{C_2}$ operator.

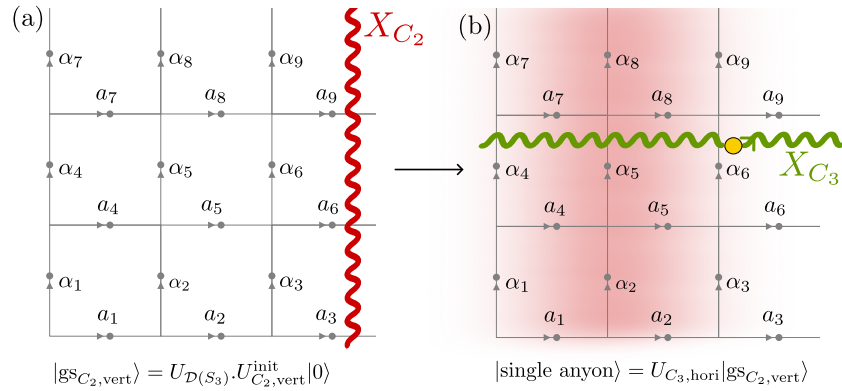


Figure 6. **Steps to create a single C_3 anyon**

We can also create a single non-Abelian [2] charge on the torus. One way to see it is by replacing the $m, \bar{m}$ pair with $e, \bar{e}$ pair while keeping the charge-conjugation domain wall the same. Running through the same argument as above, we obtain the state $|e\rangle + |\bar{e}\rangle$. Upon regauging the charge-conjugation symmetry, $|e\rangle, |\bar{e}\rangle$ becomes internal states of the [2] charge, thus we obtain a single [2] charge. Another more efficient perspective is that S_3 quantum double admits a $\mathbb{Z}_2$ anyon automorphism duality that interchange C_3 flux with [2] charge (this is inherited from the EM duality at the level of the qutrit toric code). Therefore by applying this duality transformation to the single C_3 flux state, we obtain the single [2] charge state.

S8. FIDELITY BOUNDS FOR $\mathcal{D}(S_3)$ GROUND STATE

This section establishes lower and upper bounds on the fidelity of the prepared state, ρ , with respect to the target ground state $|\text{gs}\rangle$. The ground state is defined as the +1 eigenstate of the local vertex projectors ($A_v^{\mathbb{Z}_2}, A_v^{\mathbb{Z}_3}$) and local plaquette projectors ($B_p^{\mathbb{Z}_2}, B_p^{\mathbb{Z}_3}$). Furthermore, within the ground state manifold, $|\text{gs}\rangle$ is uniquely identified by the +1 eigenspace of the horizontal and vertical non-contractable S_3 flux projectors, Z_{hori} and Z_{vert} (see Section S2). We can express the projector onto the ground state as:

$$|\text{gs}\rangle\langle\text{gs}| = \underbrace{\prod_v A_v^{\mathbb{Z}_2}}_{:=P_A^{\mathbb{Z}_2}} \cdot \underbrace{\prod_v A_v^{\mathbb{Z}_3}}_{:=P_A^{\mathbb{Z}_3}} \cdot \underbrace{\prod_p B_p^{\mathbb{Z}_2} \cdot \prod_p B_p^{\mathbb{Z}_3} \cdot Z_{\text{hori}} \cdot Z_{\text{vert}}}_{:=P_B} \quad (\text{S70})$$

While all projectors commute, allowing for simultaneous measurement, we group them into ‘product’ projectors $P_A^{\mathbb{Z}_2}$, $P_A^{\mathbb{Z}_3}$, and P_B . These groupings are chosen to be as large as possible, subject to the constraint that their destructive measurement at the end of the circuit incurs minimal gate overhead.

For the experimentally prepared density matrix ρ , the fidelity with respect to the target state $|\text{gs}\rangle$ is:

$$\begin{aligned}\langle \text{gs} | \rho | \text{gs} \rangle &= \text{Tr}(\rho |\text{gs}\rangle \langle \text{gs}|) \\ &= \text{Tr}\left(\rho P_A^{\mathbb{Z}_2} P_A^{\mathbb{Z}_3} P_B\right)\end{aligned}\tag{S71}$$

We define the expectation values of the product projectors as:

$$\begin{aligned}p_A^{\mathbb{Z}_2} &:= \text{Tr}(\rho P_A^{\mathbb{Z}_2}) \\ p_A^{\mathbb{Z}_3} &:= \text{Tr}(\rho P_A^{\mathbb{Z}_3}) \\ p_B &:= \text{Tr}(\rho P_B)\end{aligned}\tag{S72}$$

To derive a lower bound for the fidelity, we use the following inequality which holds for a set of n commuting projectors $\{P_i\}$:

$$\prod_{i=1}^n P_i \geq \left(\sum_{i=1}^n P_i \right) - (n-1)\mathbb{I}.\tag{S73}$$

Taking the trace of both sides with the density matrix ρ , and applying Eqs. (S71) and (S72), we obtain:

$$\text{Tr}\left(\rho P_A^{\mathbb{Z}_2} P_A^{\mathbb{Z}_3} P_B\right) \geq (p_A^{\mathbb{Z}_2} + p_A^{\mathbb{Z}_3} + p_B - 2)\tag{S74}$$

Substituting the experimentally measured expectation values for the 3×3 lattice gives:

$$\langle \text{gs} | \rho | \text{gs} \rangle \geq 0.93(3) + 0.84(3) + 0.81(4) - 2\tag{S75}$$

$$= 0.58(6),\tag{S76}$$

and the per-qudit fidelity is

$$f := \langle \text{gs} | \rho | \text{gs} \rangle^{1/18} \geq 0.970(5).\tag{S77}$$

Similarly, to get the upper bound for fidelity, we use the fact that the product of projectors is bounded by any individual projector:

$$\prod_{j=1}^n P_j \leq P_i \quad \forall i.\tag{S78}$$

Taking the trace with ρ yields:

$$\text{Tr}\left(\rho P_A^{\mathbb{Z}_2} P_A^{\mathbb{Z}_3} P_B\right) \leq \min\left\{p_A^{\mathbb{Z}_2}, p_A^{\mathbb{Z}_3}, p_B\right\}.\tag{S79}$$

Substituting the expectation values, we obtain:

$$\langle \text{gs} | \rho | \text{gs} \rangle \leq 0.81(4).\tag{S80}$$

In summary, the fidelity per qudit, f , for the ground state of the $\mathcal{D}(S_3)$ model on the 3×3 lattice is bounded by:

$$0.970(5) \leq f \leq 0.988(3).\tag{S81}$$

S9. BUREAU OF STANDARDS: THEORY

As touched upon in the main text, we are encoding a logical qutrit inside the zero-flux fusion subspace of two C_2 fluxes, defined by the following fusion rule:

$$C_2 \times C_2 = [+] + [2].\tag{S82}$$

One qutrit basis state is supplied by the $[+]$ channel, and the other two are furnished by the $[2]$ channel—see Fig. 7(a). We note that this is a somewhat different logical encoding strategy that might be familiar to some readers, as we are

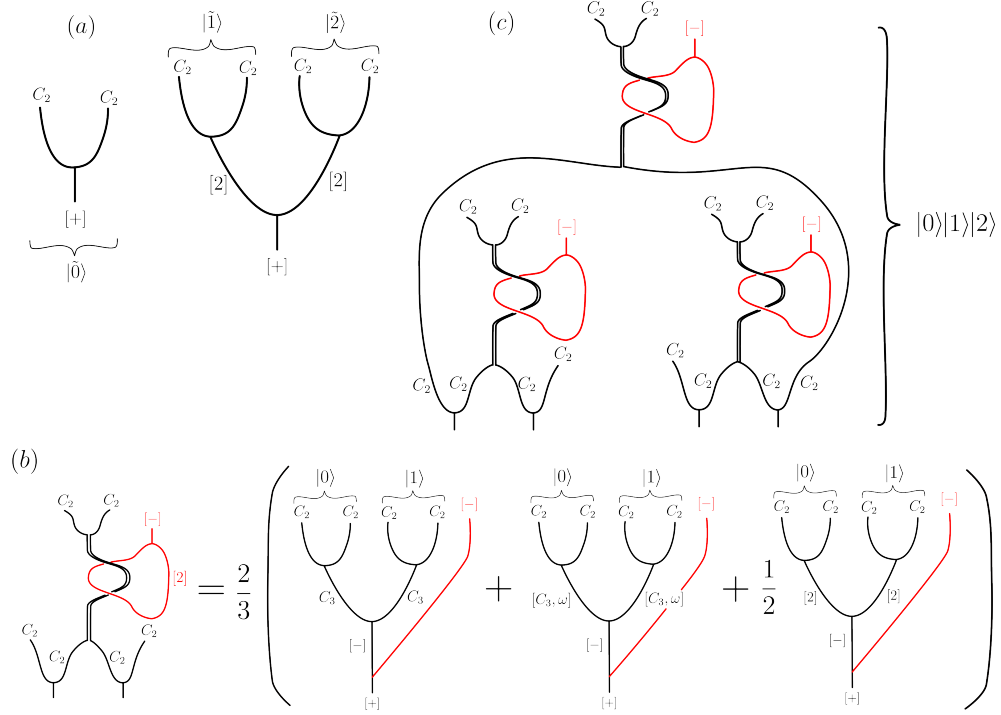


Figure 7. **Reformulating our approach as a fusion tree.** (a) Qutrit dual basis fusion trees; a $|\tilde{0}\rangle$ state is the vacuum state of two C_2 fluxes, while a reference pair of $|\tilde{1}\rangle$ and $|\tilde{2}\rangle$ state can be prepared by splitting a vacuum pair of $[2]$ charges into four C_2 fluxes. Any further logical qubits will have to be compared with the $|\tilde{1}\rangle$ and $|\tilde{2}\rangle$ states to set their values in the dual basis. (b) The first step in setting up the computation basis bureau of standards. We assume the left pair is set as our reference $|0\rangle$, and we want to compare the right pair of fluxes in the computational basis. If the $[2]$ charge braiding yields a remnant $[-]$, the right pair cannot be a $|0\rangle$, and so we can set it to be a $|1\rangle$. By re-writing the fusion and braiding diagram, we can express this $|0\rangle|1\rangle$ state in a more conventional fusion tree form, as a superposition over different intermediate fusion channels for the two logical flux pairs. (c) The fusion and braiding diagram corresponding to a full set of reference states, showing the three comparison measurements that ensure all three pairs of fluxes are in different computational basis states.

resolving the $[2]$ charge outcome into its internal states. This type of encoding is sometimes known as a “split” fusion channel encoding [S4], to contrast with the usual fusion channel encoding where every anyon type has a single logical state associated with it.

The split fusion channel encoding treats the internal states of the $[2]$ charge as distinct logical states: however, these internal states have no definite labeling, and are only physically meaningful as relative states. This subtlety necessitates the use of a “bureau of standards” which enables us to set a consistent convention of which state is which.

It is important to emphasize that, although we are using the internal states of the $[2]$ charge as logical states, our encoding is *still topologically protected*, and cannot be accessed (and hence corrupted) locally. This is because the $[2]$ charge is itself stored nonlocally between the C_2 fluxes that form the logical qutrit. Ultimately, this split fusion channel kind of encoding is equivalent to a more traditional fusion tree approach: we must simply take into account the reference states which form the bureau of standards, and the probe anyons that perform the comparison measurements, when writing down the fusion tree.

We build the bureau of standards for the Z -basis by starting with a pair of C_2 fluxes created from the vacuum. At this point, these fluxes must be in the $|\tilde{0}\rangle$ state:

$$|\tilde{0}\rangle = \frac{1}{\sqrt{3}} (|\sigma, \sigma\rangle + |\mu\sigma, \mu\sigma\rangle + |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle). \quad (\text{S83})$$

This initial pair of fluxes will act as our reference $|0\rangle$ state: any subsequent pair of fluxes must be entangled with this pair so that each branch of the wavefunction has a consistent convention. For instance, if we create a second pair of

fluxes to act as reference $|1\rangle$, we could entangle with the first “ $|0\rangle$ ” in the following way:

$$\frac{1}{\sqrt{3}}(|\sigma, \sigma\rangle |\mu\sigma, \mu\sigma\rangle + |\mu\sigma, \mu\sigma\rangle |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle + |\bar{\mu}\sigma, \bar{\mu}\sigma\rangle |\sigma, \sigma\rangle). \quad (\text{S84})$$

To create this entangled state, we first notice that if we take one flux from each pair, the combined flux of the two anyons will be C_3 -valued. So we can start with our reference $|0\rangle$ state and a new vacuum pair of C_2 fluxes, performing a $[2]$ charge braiding experiment around one half of each pair. If the $[2]$ charge fuses to a remnant $[-]$ charge, then the total flux encircled must have been a C_3 , and we have projected into the desired state. In other words, we have performed a Z -basis comparison between the reference $|0\rangle$ and the new pair, found that they must be in different states, and so labeled the new pair as a $|1\rangle$.

This procedure can be expressed as a fusion and braiding diagram, see Fig. 7(b). Using the F - and R -symbol data for S_3 , this diagram can be converted into a pure fusion diagram, demonstrating that the split fusion channel encoding can be understood as a usual fusion tree encoding with an expanded set of anyons. Fig. 7(c) shows the braiding and fusion diagram expressing the full bureau of standards protocol to create all three computational basis reference states, which can be simplified into a fusion diagram involving the three pairs of logical fluxes and the three remnant $[-]$ charges.

S10. PROTOCOL IMPLEMENTATION DETAILS

A. Pull-through gate

1. Theory prediction for local plaquette projector

In Fig. 3 caption, we note that the theoretical prediction for the expectation value local plaquette projector $B^{\mathbb{Z}_3}$ is $\frac{1}{3}$. This is due to the fact that, for the local flux label, the C_2 flux we create has a equal and symmetric weight on all three labels: $\sigma, \mu\sigma, \bar{\mu}\sigma$. The $B^{\mathbb{Z}_3}$ projector projects to the sector where the qutrit value is 0, i.e. to the label σ . Therefore, the expectation value of $B^{\mathbb{Z}_3}$ is $\frac{1}{3}$.

2. Theory prediction for nonlocal correlator

In the caption of Fig. 3, we also note that the nonlocal correlator $\langle \Pi_{W_2W_3}^{\mathbb{Z}_3} \rangle$ has a predicted value of $1/3$ at $t = 0$. The logical state is $|\tilde{0}\rangle_L |0\rangle_L = \frac{1}{\sqrt{3}}(|0\rangle_L + |1\rangle_L + |2\rangle_L) |0\rangle_L$. We see that there is equal weight between the $\mathcal{Z}$ -basis eigenstates. Therefore, the expectation value of the nonlocal correlator

$$\langle \Pi_{W_2W_3}^{\mathbb{Z}_3} \rangle = \langle \tilde{0} | \langle 0 | \Pi_{W_2W_3}^{\mathbb{Z}_3} | \tilde{0} \rangle | 0 \rangle = \left(\frac{1}{\sqrt{3}} \langle 0 | \langle 0 | \right) \left(\frac{1}{\sqrt{3}} | 0 \rangle | 0 \rangle \right) = \frac{1}{3}. \quad (\text{S85})$$

At $t = 6$, the logical state is the qutrit Bell state $\frac{1}{\sqrt{3}}(|0\rangle_L |0\rangle_L + |1\rangle_L |2\rangle_L + |2\rangle_L |1\rangle_L)$. The $\Pi_{W_2W_3}^{\mathbb{Z}_3}$ correlator is designed to evaluate the $\mathbb{Z}_3$ part of the global flux and check whether it multiply to the trivial group element. For example, the state $|1\rangle_L |2\rangle_L$ has global flux $\mu\sigma$ and $\bar{\mu}\sigma$, such that the total qutrit flux is $\mu \cdot \bar{\mu} = e$; similarly for the $|2\rangle_L |1\rangle_L$ state. Therefore, the state is manifestly in the $+1$ eigenstate of the $\Pi_{W_2W_3}^{\mathbb{Z}_3}$ correlator.

3. Circuits

The system is on a 3×3 square lattice, where the indexing of qutrits (denoted by a_i, α_i) and qubits (denoted by b_i, β_i) on each edge is as follows:

4. Step 1: Initialize control qutrit

We initialize a pair of C_2 flux (at the (1,1) and (2,1) plaquettes) to be the control qutrit in the $|\tilde{0}\rangle_L$ state at $t = 0$ depicted in the main text Fig. 3. For the circuit below, the ancilla is initialized in the $|0\rangle_a$ state, such that the first qutrit $\mathcal{H}$ transforms the ancilla to be in the $|\tilde{0}\rangle_a$ state. Then, the flux’s internal state is entangled with the ancilla via

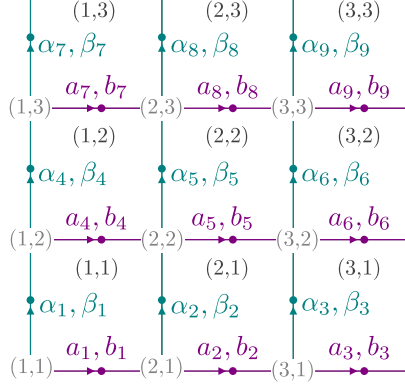
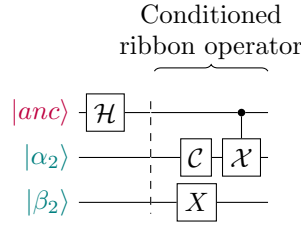
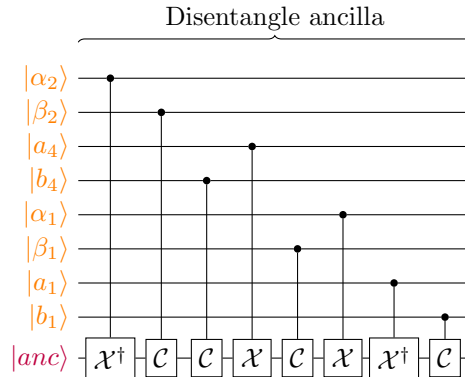


Figure 8. Labeling of vertices, edges, and plaquettes on a 3×3 square lattice on torus. The coordinate of vertex and plaquette are of the form (x, y) indexed from 1. The edges are indexed from left to right and bottom to top (alternating between horizontal and vertical edges).

action of ribbon operator with internal state conditioned on the ancilla state; we refer to this as state injection (see circuit below). This creates the state $|0\rangle_L |0\rangle_a + |1\rangle_L |1\rangle_a + |2\rangle_L |2\rangle_a$. The action is derived from the unitary circuit representation of C_2 ribbon operators (an example is shown in Eq. S21). Significant simplifications in the circuit (in particular, there are no ungauging and regauging layers) derives from the fact that there is no direct triangle in the ribbon.



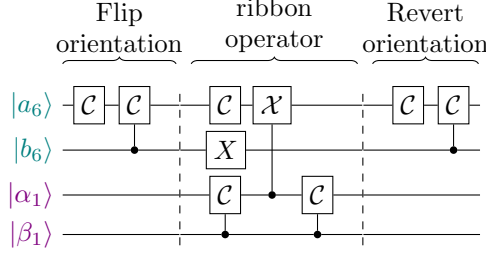
After creating the entangled state between the data qutrit and the ancilla, we need to disentangle the ancilla to initialize the data qutrit in the $|\bar{1}\rangle_L$ state. To disentangle the ancilla, we need to apply a $C_L \mathcal{X}_a^\dagger$ gate so as to obtain the state $|\bar{1}\rangle_L |0\rangle_a$. To control on the logical qutrit, we need to condition on the value of the flux loop (based at the origin) around the $(3, 1)$ plaquette. The following circuit implements the $C_L \mathcal{X}_a^\dagger$ gate to disentangles ancilla from the logical state.



5. Initialize target qutrit

We initialize a second C_2 flux pair (at the $(3,1)$ and $(3,2)$ plaquettes) to serve as the target qutrit in the $|0\rangle$ state at $t = 0$ depicted in the main text Fig. 3. As this is a $\mathcal{Z}$ -basis eigenstate, no ancilla is needed to initialize the state. The

circuit is conjugated by gates that change the orientation of the edge for the qudits $|a_6, b_6\rangle$ such that the orientation of the triangle operator is aligned.



6. Coherent moving of flux example

Here, we provide an example circuit for moving flux coherently, when we move control qutrit from the (2,1) plaquette to the (3,3) plaquette as illustrated in Fig. 9. This is the step at $t = 1$ depicted in the main text Fig. 3. For all the remaining steps of the protocol (from $t = 2$ to $t = 6$), similar circuits are used to coherently move the fluxes to complete the braiding.

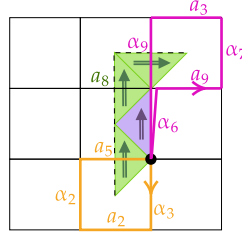
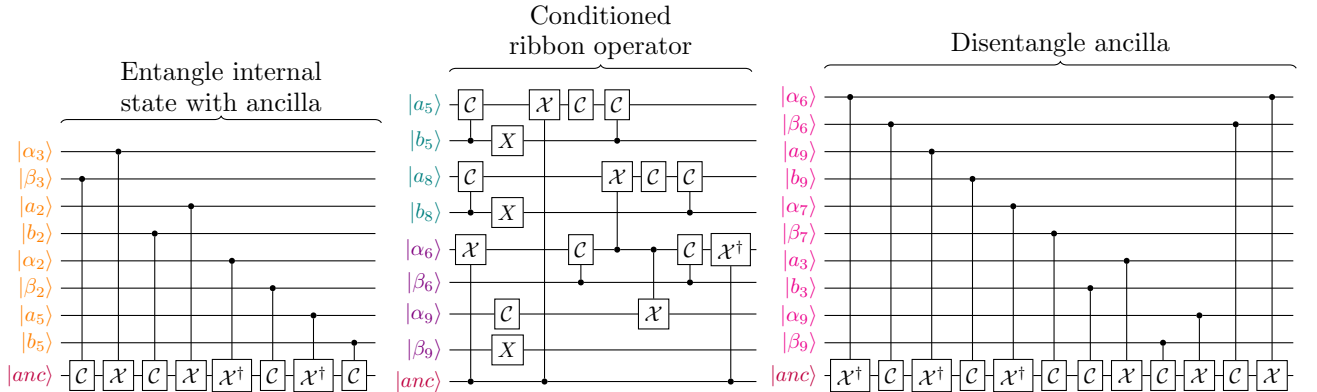


Figure 9. Schematic for the pull-through step at $t = 1$, moving a flux of the control flux pair.

First, the local flux internal state of the flux (measured in a clockwise loop depicted in orange based at the black dot) is entangled with the ancilla. Then, conditioned on the ancilla, a corresponding ribbon operator is applied to cancel out the flux at the initial plaquette. Finally, the ancilla is disentangled by applying a $C_L \mathcal{X}_a^\dagger$ gate (flux value measured in a counterclockwise loop depicted in magenta) based at the same base point (black dot). We emphasize that the base point has to be the same as the start of the ribbon to ensure that the ancilla is properly disentangled (see Section S3 for explanation).

The ancilla is initialized in the $|0\rangle$ state. The circuit on the left panel entangles the ancilla state with the internal local flux state of one of the control qutrit by implementing a $C_L \mathcal{X}_a$ gate. For the circuit in the middle panel, conditioned on the ancilla state, it applies the corresponding ribbon operator to cancel out the flux at the initial plaquette. Since the ribbon starts with a dual triangle, whose action is decoupled from the rest of the ribbon, its action is directly conditional on the ancilla state, just as the first direct triangle. Finally, for the circuit on the right panel, the ancilla is disentangled by implementing a $C_L \mathcal{X}_a^\dagger$ gate.



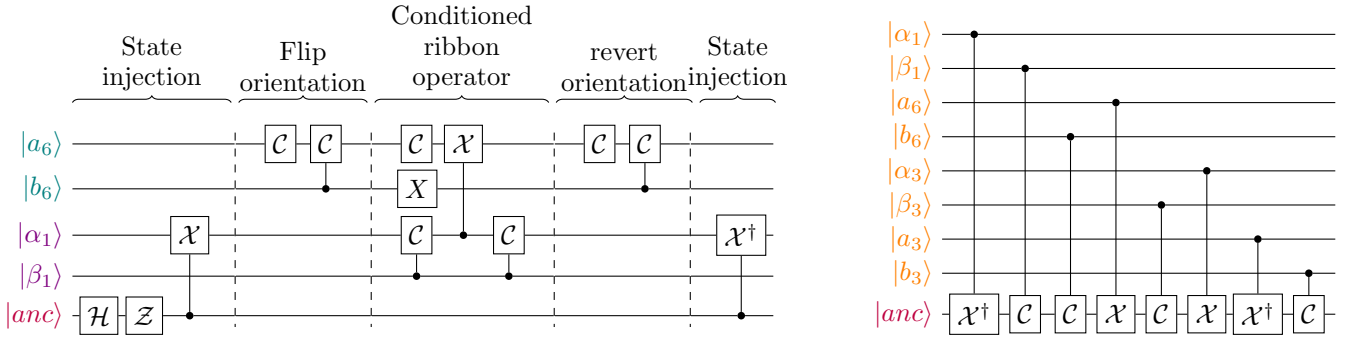
B. $\mathcal{X}$ -basis measurement

Below, we detail the circuit for implementing the $\mathcal{X}$ -basis measurement protocol for the case where the data logical qutrit is in the $|\bar{1}\rangle$ state. For the case where the data logical qutrit is in the $|\bar{0}\rangle$ state, the only difference is in step 1, where we omit the $\mathcal{Z}$ gate in the initial state injection stage; the rest of the circuit remains the same.

1. Step 1: Create data logical $|\bar{1}\rangle$

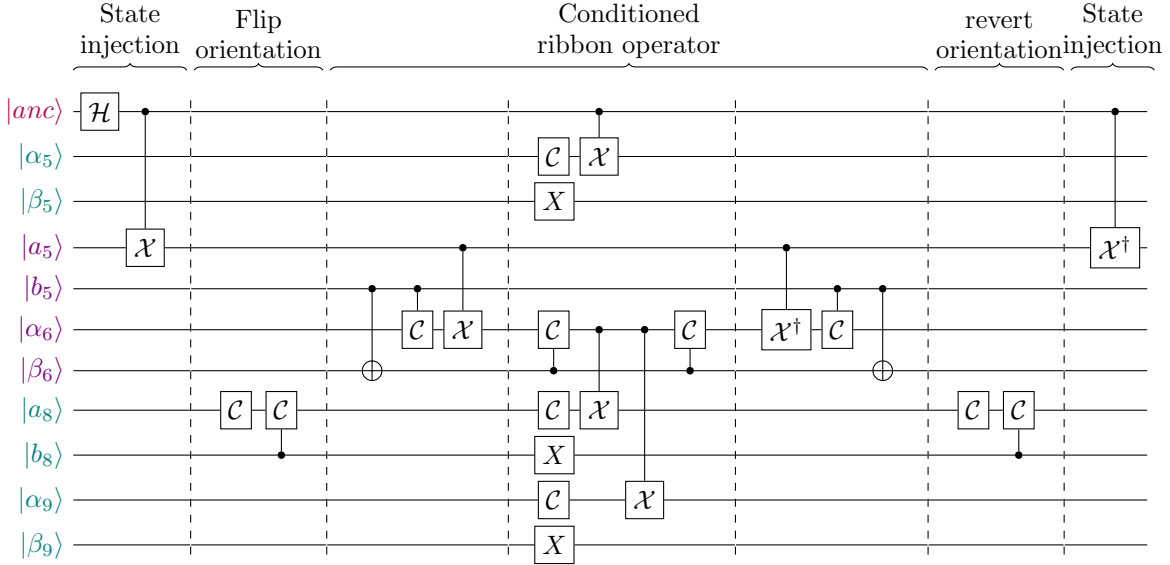
First, we create a C_2 flux pair at the (3,1) and (3,2) plaquettes (see Fig. 8 for labels); this flux pair serves as the data qutrit and is initialized in the $|\bar{1}\rangle_L$ logical state by the end of step 1 using state injection similar to the procedure in the step 1 of pull-through gate (see Section S10 A 4). The circuit implementing the entangling is shown on the left panel, where the qudits in violet (teal) font are acted on by direct (dual) triangle of the ribbon operator.

The circuit on the right panel disentangle the ancilla by applying a $C_L \mathcal{X}_a^\dagger$ gate to obtain the (periodic boundary condition is used to shorten the flux loop).

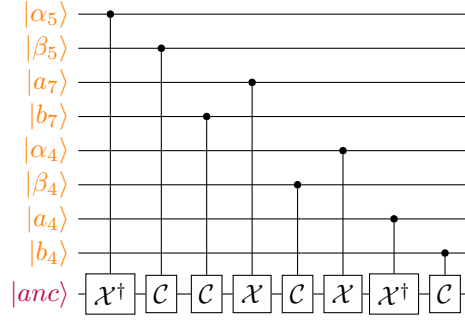


2. Step 2: Create measurement pure flux

We then create the second flux pair in the $|\bar{0}\rangle_L$ logical state at the (1,2) and (3,3) plaquette, using similar techniques as in step 1.



Similar to the second part of step 1, we disentangle the ancilla from the logical state by applying a $C_L \mathcal{X}_a^\dagger$ gate, conditioned on the value of the flux loop (based at the origin) around the (1,2) plaquette using periodic boundary condition.



3. Step 3-5: Coherent moving fluxes

We schematically illustrate the coherent moving circuits for step 3-5, similar to those implemented in Section S10 A 6.

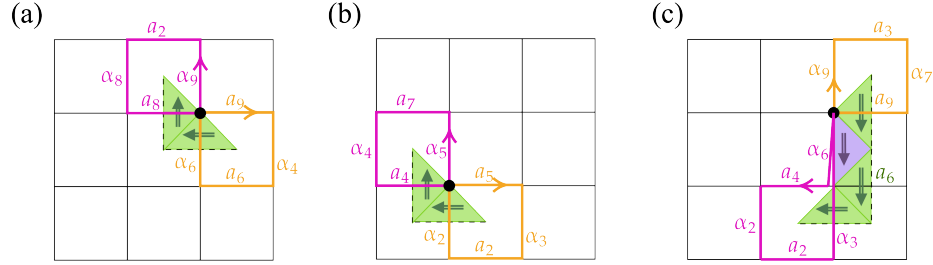
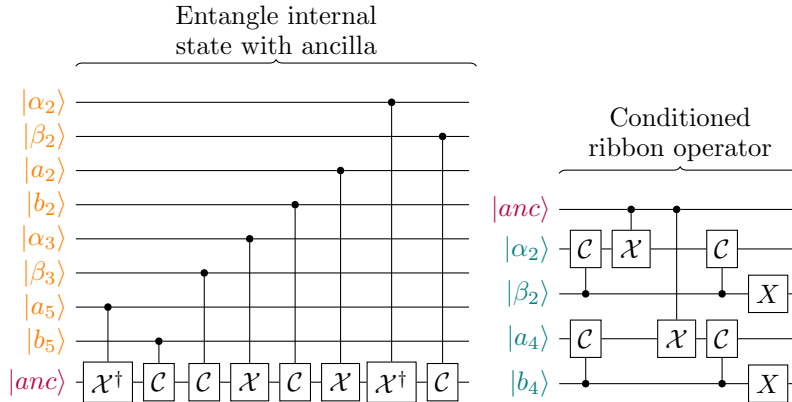


Figure 10. **Schematics for the coherent moving circuits for $\mathcal{X}$ -basis measurement.** (a) Step 3, (b) step 4, and (c) step 5.

4. Step 6: Fuse C_2 flux

For the last step of the $\mathcal{X}$ -basis measurement protocol, it involves fusing two C_2 flux, which warrants more explanation as the circuit towards the end is different from that of the coherent moving circuits in previous steps.

The entangling (left panel) and conditioned ribbon operator (middle panel) of the circuit proceeds in the same way as the coherent moving circuits, as shown below:



At this point, the flux pair have fused together and leaves behind a remnant charge at the fusion site. The ancilla is entangled with the system in a non-obvious basis. First, a qutrit Hadamard $\mathcal{H}^\dagger$ is applied, because this would reset the ancilla from the $|0\rangle$ state back to the $|0\rangle$ state if the ancilla is not entangled; the state is then in the basis where the internal state of the remnant charge is entangled with the ancilla as follows:

$$|0\rangle_a |+\rangle + |1\rangle_a |2-\rangle + |2\rangle_a |2+\rangle. \quad (\text{S86})$$

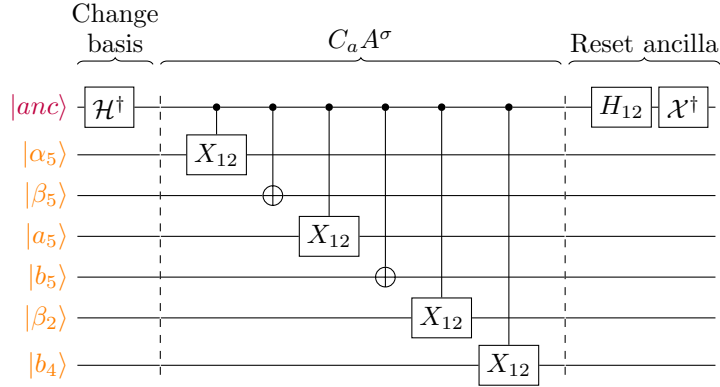
At this point, we could in principle end the protocol; projective measurement on the ancilla is equivalent to measuring the local charge at the fusion site. However, the ancilla is entangled with the system and cannot be reused until projective measurement.

We embrace a quixotic spirit by asking the question: can the ancilla be disentangled? Besides freeing the ancilla for other parts of the protocol, this also has the additional benefit of allowing the usual local vertex projectors ($A_v^{\mathbb{Z}_2}$ and $A_v^{\mathbb{Z}_3}$) to directly determine the charge violations.

Here, we make use of the fact that for the $|\tilde{1}\rangle_L$ input state of the $\mathcal{X}$ -basis measurement protocol, there is no weight in the vacuum fusion channel, implying our state simplifies to

$$|1\rangle_a |2-\rangle + |2\rangle_a |2+\rangle. \quad (\text{S87})$$

To disentangle the ancilla from the above state, we need to implement a special $C_a A^\sigma$ gate on the vertex of the fusion site: apply identity when the control is in $|1\rangle$ state, and apply A^σ when the control is in $|2\rangle$ state (since $A^\sigma |2-\rangle = |2+\rangle$). To decompose this further, the action of A^σ involves X gates on all edges of a vertex, and additional $\mathcal{C}$ gates on the oppositely aligned edges. The circuit implementation is as follows:



The definitions and physical implementations of the the qutrit-control qubit-target $C_2\text{NOT}$ gate and the qutrit-control qutrit-target CX_{12} gate are given in Eqs. (S51) and (S52), respectively.

Finally, after applying the $C_{a_{12}} A^\sigma$ gate, the ancilla is disentangled and is in the $|1\rangle + |2\rangle$ state. Therefore, a H_{12} gate (defined as qubit Hadamard on the $\{|1\rangle, |2\rangle\}$ subspace in Eq. (S47)) followed by a $\mathcal{X}^\dagger$ are applied to revert the ancilla back to the $|0\rangle$ state.

C. $\mathcal{Z}$ -basis measurement

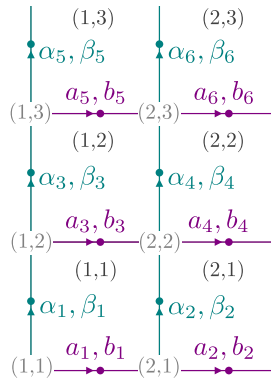
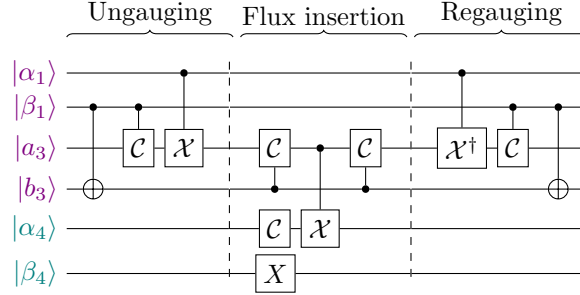


Figure 11. Labeling of vertices, edges, and plaquettes on a 3×2 periodic square lattice. Vertex and plaquette coordinates are (x, y) , indexed from 1. Edges are ordered left to right and bottom to top, alternating between horizontal and vertical.

1. Example of generalized ribbon operator

Here, we provide an example of how the generalized ribbon operator (see Section S1 D 3) is used to specify a logical qutrit in the absolute encoding IV A. It involves prepending the usual ribbon with a “direct string” that connects from the starting vertex of the ribbon to the origin. Operationally, this means that the ungauging (regauging) step of the ribbon operator starts (ends) at the origin.

As an example, we show a circuit implementation of step 1 of the $\mathcal{Z}$ -basis measurement protocol IV D, where we create the data qutrit in the $|0\rangle_L$ state at the (1,2) and (2,2) plaquettes:

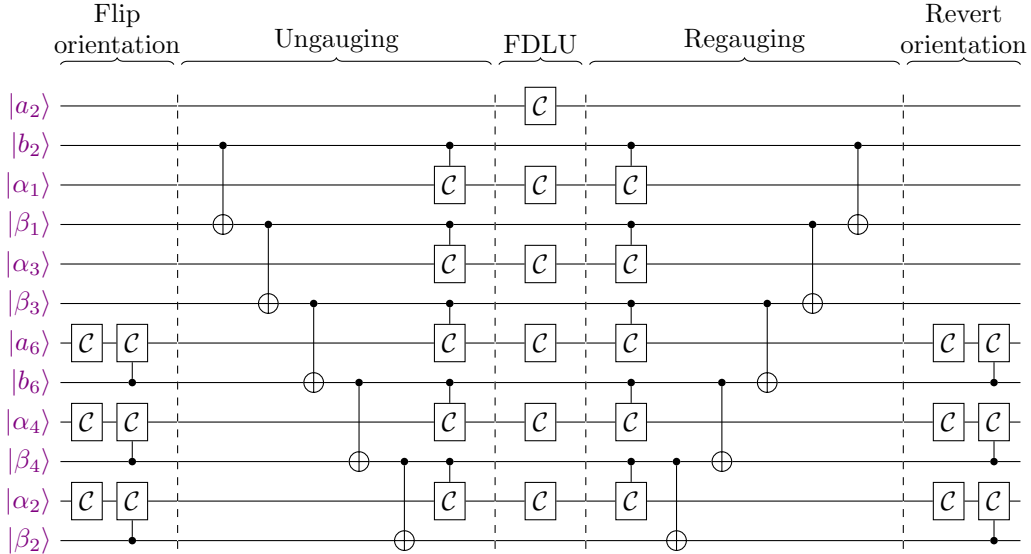


We can decompose the above circuit into direct string and the following triangle operators in two ways: 1) the direct string has support on the (α_1, β_1) and (α_3, β_3) edges, and the first triangle operator is the dual triangle acting on the (α_4, β_4) edge 2) the direct string has support on just the (α_1, β_1) edge, and the following is a ribbon consisting of two triangle operators (a direct triangle acting on the (α_3, β_3) edge and a dual triangle acting on the (α_4, β_4) edge). In either decomposition, the direct string is necessary, since there is no way to express the above circuit using only consecutive triangle operators.

Similar circuits are used for creating logical states in Section IV E and V.

2. [2] charge braiding

When we braid [2] charges around the two fluxes for the $\mathcal{Z}$ -basis comparison measurement, we can perform it in one shot, i.e. implementing a single closed ribbon operator without the need for coherent moving. An example circuit is given below, at the step where we compare the data qutrit with the reference $|0\rangle_L$ qutrit:



It is slightly modified from the [2] charge unitary ribbon given in Eq. S20. The orientation of some direct edges are flipped at first, such that all direct edges' orientation is aligned with respect to the direction of the ribbon operator. The orientation is reverted after the application of the unitary ribbon.

Because of the nontrivial braiding of the $[2]$ charge with a C_3 flux (e.g. when performing comparison measurement between the data qutrit and the reference $|1\rangle_L$ state), a closed loop ribbon operator is not “rotationally-invariant” along the loop. The start and end point of the ribbon operator determines the location of any remnant charge (in this case, a possible remnant $[-]$ charge) that breaks this rotational invariance.

Similar circuits are used for $\mathcal{Z}$ comparison measurements implemented in Section IV E and V.

D. $\mathcal{Z}$ bureau of standards

The circuits used in $\mathcal{Z}$ bureau of standards are very similar to that used in $\mathcal{Z}$ -basis measurement (see Section S10 C). The main difference is that the initial states are all in the $|0\rangle_L$ state (an example circuit creating such state is given in Section S10 B 2. Subsequent $[2]$ charge braiding circuits are adapted from the example given in Section S10 C 2.

1. Multi-verse interpretation of relative encoding

First, we outline an argument by Mochon [S16] that illustrates how we can interpret the relative encoding as the multi-verse of possible flux labels, such that the absolute encoding is one “branch” of this multi-verse. Consider the state $|\hat{0}\rangle_L = \frac{1}{\sqrt{3}}(|0\rangle + |1\rangle + |2\rangle)$. If we forget about the coherence between the $\mathcal{Z}$ eigenstates, it is a completely mixed state $\rho = \frac{1}{3}(|0\rangle\langle 0| + |1\rangle\langle 1| + |2\rangle\langle 2|)$. Equivalently, we can obtain this mixed state ρ by starting with a qutrit Bell state $|0\rangle|0\rangle + |1\rangle|1\rangle + |2\rangle|2\rangle$ and trace out the second qutrit. Then, we can use it as a $\mathcal{Z}$ -basis reference state. Without loss of generality, we can assert it to be the $|0\rangle_L$ in the relative encoding. Then, other $\mathcal{Z}$ states can be instantiated by $\mathcal{Z}$ -basis comparison measurement with the reference state.

For example, given a reference $|0\rangle_L$ in the relative encoding, if we instantiate another $|0\rangle_L$ state, this is actually the state $\psi \sim |0\rangle|0\rangle + |1\rangle|1\rangle + |2\rangle|2\rangle$ in the absolute encoding. What matters is that the logical qutrit values are the same for every “branch” of this state.

2. $[2]$ charge braiding once suffices for instantiating $\mathcal{Z}$ reference states

Next, we explain why we only perform $[2]$ charge braiding once rather than twice to instantiate the $\mathcal{Z}$ bureau of standards reference states.

For a full $\mathcal{Z}$ -basis measurement, we need to braid a $[2]$ charge pair twice to ensure that we obtain the state with the correct coherence in the orthogonal subspace after the measurement. For example, in Section V, two $[2]$ charge braiding are implemented to ensure that after projecting out of the $|2\rangle_L$ subspace, the resulting state $|0\rangle_L + \omega|1\rangle_L$ has the correct relative phase factor.

However, for the $\mathcal{Z}$ bureau of standards protocol (Section IV E), we do not require performing such $[2]$ charge braiding twice. This is because the coherence between the different “branches” of the relative encoding is not important; all that matters is that the absolute flux label of the states are distinct pair-wise for the three reference states. Hence, we do not need to repeat braiding with the $[2]$ charge twice to obtain the correct logical reference states for the $\mathcal{Z}$ bureau of standards.

E. Magic state creation

The circuits used in creating the magic state ($|0\rangle + \omega|1\rangle$) are very similar to that used in $\mathcal{Z}$ -basis measurement (see Section S10 C). The main difference is that the initial states are in the logical $|2\rangle_L$ and $|\tilde{1}\rangle_L$ state. To create the logical $|2\rangle_L$ state, it is a slight modification of the circuit given in S10 C 1: the circuit is conjugated with $\mathcal{X}^\dagger$ (see remark on C_2 ribbon in Section S1 D 2 for explanation). For the logical $|\tilde{1}\rangle_L$ state, an example circuit is given in S10 B 1. Subsequent $[2]$ charge braiding circuits are adapted from the example given in Section S10 C 2.

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